
Stream: Internet Research Task Force (IRTF)
RFC: [9380](#)
Category: Informational
Published: July 2023
ISSN: 2070-1721
Authors: A. Faz-Hernandez S. Scott N. Sullivan R. S. Wahby
Cloudflare, Inc. Cornell Tech Cloudflare, Inc. Stanford University

C. A. Wood
Cloudflare, Inc.

RFC 9380

Hashing to Elliptic Curves

Abstract

This document specifies a number of algorithms for encoding or hashing an arbitrary string to a point on an elliptic curve. This document is a product of the Crypto Forum Research Group (CFRG) in the IRTF.

Status of This Memo

This document is not an Internet Standards Track specification; it is published for informational purposes.

This document is a product of the Internet Research Task Force (IRTF). The IRTF publishes the results of Internet-related research and development activities. These results might not be suitable for deployment. This RFC represents the consensus of the Crypto Forum Research Group of the Internet Research Task Force (IRTF). Documents approved for publication by the IRSG are not candidates for any level of Internet Standard; see Section 2 of RFC 7841.

Information about the current status of this document, any errata, and how to provide feedback on it may be obtained at <https://www.rfc-editor.org/info/rfc9380>.

Copyright Notice

Copyright (c) 2023 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust's Legal Provisions Relating to IETF Documents (<https://trustee.ietf.org/license-info>) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document.

Table of Contents

1. Introduction
 - 1.1. Requirements Notation
2. Background
 - 2.1. Elliptic Curves
 - 2.2. Terminology
 - 2.2.1. Mappings
 - 2.2.2. Encodings
 - 2.2.3. Random Oracle Encodings
 - 2.2.4. Serialization
 - 2.2.5. Domain Separation
3. Encoding Byte Strings to Elliptic Curves
 - 3.1. Domain Separation Requirements
4. Utility Functions
 - 4.1. The `sgn0` Function
5. Hashing to a Finite Field
 - 5.1. Efficiency Considerations in Extension Fields
 - 5.2. `hash_to_field` Implementation
 - 5.3. `expand_message`
 - 5.3.1. `expand_message_xmd`
 - 5.3.2. `expand_message_xof`
 - 5.3.3. Using DSTs Longer Than 255 Bytes
 - 5.3.4. Defining Other `expand_message` Variants
6. Deterministic Mappings
 - 6.1. Choosing a Mapping Function
 - 6.2. Interface
 - 6.3. Notation
 - 6.4. Sign of the Resulting Point

- 6.5. Exceptional Cases
- 6.6. Mappings for Weierstrass Curves
 - 6.6.1. Shallue-van de Woestijne Method
 - 6.6.2. Simplified Shallue-van de Woestijne-Ulas Method
 - 6.6.3. Simplified SWU for $AB \neq 0$
- 6.7. Mappings for Montgomery Curves
 - 6.7.1. Elligator 2 Method
- 6.8. Mappings for Twisted Edwards Curves
 - 6.8.1. Rational Maps from Montgomery to Twisted Edwards Curves
 - 6.8.2. Elligator 2 Method
- 7. Clearing the Cofactor
- 8. Suites for Hashing
 - 8.1. Implementing a Hash-to-Curve Suite
 - 8.2. Suites for NIST P-256
 - 8.3. Suites for NIST P-384
 - 8.4. Suites for NIST P-521
 - 8.5. Suites for curve25519 and edwards25519
 - 8.6. Suites for curve448 and edwards448
 - 8.7. Suites for secp256k1
 - 8.8. Suites for BLS12-381
 - 8.8.1. BLS12-381 G1
 - 8.8.2. BLS12-381 G2
 - 8.9. Defining a New Hash-to-Curve Suite
 - 8.10. Suite ID Naming Conventions
- 9. IANA Considerations
- 10. Security Considerations
 - 10.1. Properties of Encodings
 - 10.2. Hashing Passwords
 - 10.3. Constant-Time Requirements
 - 10.4. `encode_to_curve`: Output Distribution and Indifferentiability

10.5. `hash_to_field` Security

10.6. `expand_message_xmd` Security

10.7. Domain Separation for `expand_message` Variants

10.8. Target Security Levels

11. References

11.1. Normative References

11.2. Informative References

Appendix A. Related Work

Appendix B. Hashing to `ristretto255`

Appendix C. Hashing to `decaf448`

Appendix D. Rational Maps

D.1. Generic Mapping from Montgomery to Twisted Edwards

D.2. Mapping from Weierstrass to Montgomery

Appendix E. Isogeny Maps for Suites

E.1. 3-Isogeny Map for `secp256k1`

E.2. 11-Isogeny Map for BLS12-381 G1

E.3. 3-Isogeny Map for BLS12-381 G2

Appendix F. Straight-Line Implementations of Deterministic Mappings

F.1. Shallue-van de Woestijne Method

F.2. Simplified SWU Method

F.2.1. `sqrt_ratio` Subroutine

F.3. Elligator 2 Method

Appendix G. Curve-Specific Optimized Sample Code

G.1. Interface and Projective Coordinate Systems

G.2. Elligator 2

G.2.1. `curve25519` ($q = 5 \pmod{8}$, $K = 1$)

G.2.2. `edwards25519`

G.2.3. `curve448` ($q = 3 \pmod{4}$, $K = 1$)

G.2.4. `edwards448`

G.2.5. Montgomery Curves with $q = 3 \pmod{4}$

G.2.6. Montgomery Curves with $q = 5 \pmod{8}$

G.3. Cofactor Clearing for BLS12-381 G2

Appendix H. Scripts for Parameter Generation

H.1. Finding Z for the Shallue-van de Woestijne Map

H.2. Finding Z for Simplified SWU

H.3. Finding Z for Elligator 2

Appendix I. sqrt and is_square Functions

I.1. sqrt for $q = 3 \pmod{4}$

I.2. sqrt for $q = 5 \pmod{8}$

I.3. sqrt for $q = 9 \pmod{16}$

I.4. Constant-Time Tonelli-Shanks Algorithm

I.5. is_square for $F = GF(p^2)$

Appendix J. Suite Test Vectors

J.1. NIST P-256

J.1.1. P256_XMD:SHA-256_SSWU_RO_

J.1.2. P256_XMD:SHA-256_SSWU_NU_

J.2. NIST P-384

J.2.1. P384_XMD:SHA-384_SSWU_RO_

J.2.2. P384_XMD:SHA-384_SSWU_NU_

J.3. NIST P-521

J.3.1. P521_XMD:SHA-512_SSWU_RO_

J.3.2. P521_XMD:SHA-512_SSWU_NU_

J.4. curve25519

J.4.1. curve25519_XMD:SHA-512_ELL2_RO_

J.4.2. curve25519_XMD:SHA-512_ELL2_NU_

J.5. edwards25519

J.5.1. edwards25519_XMD:SHA-512_ELL2_RO_

J.5.2. edwards25519_XMD:SHA-512_ELL2_NU_

J.6. curve448

J.6.1. curve448_XOF:SHAKE256_ELL2_RO_

J.6.2. curve448_XOF:SHAKE256_ELL2_NU_

J.7. edwards448

J.7.1. edwards448_XOF:SHAKE256_ELL2_RO_

J.7.2. edwards448_XOF:SHAKE256_ELL2_NU_

J.8. secp256k1

J.8.1. secp256k1_XMD:SHA-256_SSWU_RO_

J.8.2. secp256k1_XMD:SHA-256_SSWU_NU_

J.9. BLS12-381 G1

J.9.1. BLS12381G1_XMD:SHA-256_SSWU_RO_

J.9.2. BLS12381G1_XMD:SHA-256_SSWU_NU_

J.10. BLS12-381 G2

J.10.1. BLS12381G2_XMD:SHA-256_SSWU_RO_

J.10.2. BLS12381G2_XMD:SHA-256_SSWU_NU_

Appendix K. Expand Test Vectors

K.1. expand_message_xmd(SHA-256)

K.2. expand_message_xmd(SHA-256) (Long DST)

K.3. expand_message_xmd(SHA-512)

K.4. expand_message_xof(SHAKE128)

K.5. expand_message_xof(SHAKE128) (Long DST)

K.6. expand_message_xof(SHAKE256)

Acknowledgements

Contributors

Authors' Addresses

1. Introduction

Many cryptographic protocols require a procedure that encodes an arbitrary input, e.g., a password, to a point on an elliptic curve. This procedure is known as hashing to an elliptic curve, where the hashing procedure provides collision resistance and does not reveal the discrete logarithm of the output point. Prominent examples of cryptosystems that hash to elliptic curves include password-authenticated key exchanges [BM92] [J96] [BMP00] [p1363.2], Identity-Based Encryption [BF01], Boneh-Lynn-Shacham signatures [BLS01] [BLS-SIG], Verifiable Random Functions [MRV99] [VRF], and Oblivious Pseudorandom Functions [NR97] [OPRFs].

Unfortunately for implementors, the precise hash function that is suitable for a given protocol implemented using a given elliptic curve is often unclear from the protocol's description. Meanwhile, an incorrect choice of hash function can have disastrous consequences for security.

This document aims to bridge this gap by providing a comprehensive set of recommended algorithms for a range of curve types. Each algorithm conforms to a common interface: it takes as input an arbitrary-length byte string and produces as output a point on an elliptic curve. We provide implementation details for each algorithm, describe the security rationale behind each recommendation, and give guidance for elliptic curves that are not explicitly covered. We also present optimized implementations for internal functions used by these algorithms.

Readers wishing to quickly specify or implement a conforming hash function should consult [Section 8](#), which lists recommended hash-to-curve suites and describes both how to implement an existing suite and how to specify a new one.

This document does not specify probabilistic rejection sampling methods, sometimes referred to as "try-and-increment" or "hunt-and-peck," because the goal is to specify algorithms that can plausibly be computed in constant time. Use of these probabilistic rejection methods is **NOT RECOMMENDED**, because they have been a perennial cause of side-channel vulnerabilities. See Dragonblood [VR20] as one example of this problem in practice, and see [Appendix A](#) for an informal description of rejection sampling methods and the timing side-channels they introduce.

This document represents the consensus of the Crypto Forum Research Group (CFRG).

1.1. Requirements Notation

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

2. Background

2.1. Elliptic Curves

The following is a brief definition of elliptic curves, with an emphasis on important parameters and their relation to hashing to curves. For further reference on elliptic curves, consult [\[CFADLNV05\]](#) or [\[W08\]](#).

Let F be the finite field $GF(q)$ of prime characteristic $p > 3$. (This document does not consider elliptic curves over fields of characteristic 2 or 3.) In most cases, F is a prime field, so $q = p$. Otherwise, F is an extension field, so $q = p^m$ for an integer $m > 1$. This document writes elements of extension fields in a primitive element or polynomial basis, i.e., as a vector of m elements of $GF(p)$ written in ascending order by degree. The entries of this vector are indexed in ascending order starting from 1, i.e., $x = (x_1, x_2, \dots, x_m)$. For example, if $q = p^2$ and the primitive element basis is $(1, I)$, then $x = (a, b)$ corresponds to the element $a + b * I$, where $x_1 = a$ and $x_2 = b$. (Note that all choices of basis are isomorphic, but certain choices may result in a more efficient implementation; this document does not make any particular assumptions about choice of basis.)

An elliptic curve E is specified by an equation in two variables and a finite field F . An elliptic curve equation takes one of several standard forms, including (but not limited to) Weierstrass, Montgomery, and Edwards.

The curve E induces an algebraic group of order n , meaning that the group has n distinct elements. (This document uses additive notation for the elliptic curve group operation.) Elements of an elliptic curve group are points with affine coordinates (x, y) satisfying the curve equation, where x and y are elements of F . In addition, all elliptic curve groups have a distinguished element, the identity point, which acts as the identity element for the group operation. On certain curves (including Weierstrass and Montgomery curves), the identity point cannot be represented as an (x, y) coordinate pair.

For security reasons, cryptographic applications of elliptic curves generally require using a (sub)group of prime order. Let G be such a subgroup of the curve of prime order r , where $n = h * r$. In this equation, h is an integer called the cofactor. An algorithm that takes as input an arbitrary point on the curve E and produces as output a point in the subgroup G of E is said to "clear the cofactor." Such algorithms are discussed in [Section 7](#).

Certain hash-to-curve algorithms restrict the form of the curve equation, the characteristic of the field, or the parameters of the curve. For each algorithm presented, this document lists the relevant restrictions.

The table below summarizes quantities relevant to hashing to curves:

Symbol	Meaning	Relevance
F, q, p	A finite field F of characteristic p and $\#F = q = p^m$.	For prime fields, $q = p$; otherwise, $q = p^m$ and $m > 1$.
E	Elliptic curve.	E is specified by an equation and a field F .
n	Number of points on the elliptic curve E .	$n = h * r$, for h and r defined below.
G	A prime-order subgroup of the points on E .	G is a destination group to which byte strings are encoded.
r	Order of G .	r is a prime factor of n (usually, the largest such factor).
h	Cofactor, $h \geq 1$.	h is an integer satisfying $n = h * r$.

Table 1: Summary of Symbols and Their Definitions

2.2. Terminology

In this section, we define important terms used throughout the document.

2.2.1. Mappings

A mapping is a deterministic function from an element of the field F to a point on an elliptic curve E defined over F .

In general, the set of all points that a mapping can produce over all possible inputs may be only a subset of the points on an elliptic curve (i.e., the mapping may not be surjective). In addition, a mapping may output the same point for two or more distinct inputs (i.e., the mapping may not be injective). For example, consider a mapping from F to an elliptic curve having n points: if the number of elements of F is not equal to n , then this mapping cannot be bijective (i.e., both injective and surjective), since the mapping is defined to be deterministic.

Mappings may also be invertible, meaning that there is an efficient algorithm that, for any point P output by the mapping, outputs an x in F such that applying the mapping to x outputs P . Some of the mappings given in [Section 6](#) are invertible, but this document does not discuss inversion algorithms.

2.2.2. Encodings

Encodings are closely related to mappings. Like a mapping, an encoding is a function that outputs a point on an elliptic curve. In contrast to a mapping, however, the input to an encoding is an arbitrary-length byte string.

This document constructs deterministic encodings by composing a hash function H_f with a deterministic mapping. In particular, H_f takes as input an arbitrary string and outputs an element of F . The deterministic mapping takes that element as input and outputs a point on an

elliptic curve E defined over F . Since H_f takes arbitrary-length byte strings as inputs, it cannot be injective: the set of inputs is larger than the set of outputs, so there must be distinct inputs that give the same output (i.e., there must be collisions). Thus, any encoding built from H_f is also not injective.

Like mappings, encodings may be invertible, meaning that there is an efficient algorithm that, for any point P output by the encoding, outputs a string s such that applying the encoding to s outputs P . However, the instantiation of H_f used by all encodings specified in this document ([Section 5](#)) is not invertible; thus, those encodings are also not invertible.

In some applications of hashing to elliptic curves, it is important that encodings do not leak information through side channels. [\[VR20\]](#) is one example of this type of leakage leading to a security vulnerability. See [Section 10.3](#) for further discussion.

2.2.3. Random Oracle Encodings

A random-oracle encoding satisfies a strong property: it can be proved indistinguishable from a random oracle [\[MRH04\]](#) under a suitable assumption.

Both constructions described in [Section 3](#) are indistinguishable from random oracles [\[MRH04\]](#) when instantiated following the guidelines in this document. The constructions differ in their output distributions: one gives a uniformly random point on the curve, the other gives a point sampled from a nonuniform distribution.

A random-oracle encoding with a uniform output distribution is suitable for use in many cryptographic protocols proven secure in the random-oracle model. See [Section 10.1](#) for further discussion.

2.2.4. Serialization

A procedure related to encoding is the conversion of an elliptic curve point to a bit string. This is called serialization, and it is typically used for compactly storing or transmitting points. The inverse operation, deserialization, converts a bit string to an elliptic curve point. For example, [\[SEC1\]](#) and [\[p1363a\]](#) give standard methods for serialization and deserialization.

Deserialization is different from encoding in that only certain strings (namely, those output by the serialization procedure) can be deserialized. In contrast, this document is concerned with encodings from arbitrary strings to elliptic curve points. This document does not cover serialization or deserialization.

2.2.5. Domain Separation

Cryptographic protocols proven secure in the random-oracle model are often analyzed under the assumption that the random oracle only answers queries associated with that protocol (including queries made by adversaries) [\[BR93\]](#). In practice, this assumption does not hold if two protocols use the same function to instantiate the random oracle. Concretely, consider protocols P_1 and P_2 that query a random-oracle RO : if P_1 and P_2 both query RO on the same value x , the security analysis of one or both protocols may be invalidated.

A common way of addressing this issue is called domain separation, which allows a single random oracle to simulate multiple, independent oracles. This is effected by ensuring that each simulated oracle sees queries that are distinct from those seen by all other simulated oracles. For example, to simulate two oracles RO1 and RO2 given a single oracle RO, one might define

```
RO1(x) := RO("RO1" || x)
RO2(x) := RO("RO2" || x)
```

where `||` is the concatenation operator. In this example, "RO1" and "RO2" are called domain separation tags (DSTs); they ensure that queries to RO1 and RO2 cannot result in identical queries to RO, meaning that it is safe to treat RO1 and RO2 as independent oracles.

In general, domain separation requires defining a distinct injective encoding for each oracle being simulated. In the above example, "RO1" and "RO2" have the same length and thus satisfy this requirement when used as prefixes. The algorithms specified in this document take a different approach to ensuring injectivity; see Sections 5.3 and 10.7 for more details.

3. Encoding Byte Strings to Elliptic Curves

This section presents a general framework and interface for encoding byte strings to points on an elliptic curve. The constructions in this section rely on three basic functions:

- The function `hash_to_field` hashes arbitrary-length byte strings to a list of one or more elements of a finite field F ; its implementation is defined in [Section 5](#).

```
hash_to_field(msg, count)

Input:
- msg, a byte string containing the message to hash.
- count, the number of elements of  $F$  to output.

Output:
-  $(u_0, \dots, u_{(count - 1)})$ , a list of field elements.

Steps: defined in Section 5.
```

- The function `map_to_curve` calculates a point on the elliptic curve E from an element of the finite field F over which E is defined. [Section 6](#) describes mappings for a range of curve families.

```
map_to_curve(u)

Input:  $u$ , an element of field  $F$ .
Output:  $Q$ , a point on the elliptic curve  $E$ .
Steps: defined in Section 6.
```

- The function `clear_cofactor` sends any point on the curve E to the subgroup G of E . [Section 7](#) describes methods to perform this operation.

```
clear_cofactor(Q)
```

Input: Q , a point on the elliptic curve E .

Output: P , a point in G .

Steps: defined in [Section 7](#).

The two encodings ([Section 2.2.2](#)) defined in this section have the same interface and are both random-oracle encodings ([Section 2.2.3](#)). Both are implemented as a composition of the three basic functions above. The difference between the two is that their outputs are sampled from different distributions:

- `encode_to_curve` is a nonuniform encoding from byte strings to points in G . That is, the distribution of its output is not uniformly random in G : the set of possible outputs of `encode_to_curve` is only a fraction of the points in G , and some points in this set are more likely to be output than others. [Section 10.4](#) gives a more precise definition of `encode_to_curve`'s output distribution.

```
encode_to_curve(msg)
```

Input: `msg`, an arbitrary-length byte string.

Output: P , a point in G .

Steps:

1. $u = \text{hash_to_field}(\text{msg}, 1)$
2. $Q = \text{map_to_curve}(u[0])$
3. $P = \text{clear_cofactor}(Q)$
4. return P

- `hash_to_curve` is a uniform encoding from byte strings to points in G . That is, the distribution of its output is statistically close to uniform in G .

This function is suitable for most applications requiring a random oracle returning points in G , when instantiated with any of the `map_to_curve` functions described in [Section 6](#). See [Section 10.1](#) for further discussion.

```
hash_to_curve(msg)
```

Input: `msg`, an arbitrary-length byte string.

Output: P , a point in G .

Steps:

1. $u = \text{hash_to_field}(\text{msg}, 2)$
2. $Q_0 = \text{map_to_curve}(u[0])$
3. $Q_1 = \text{map_to_curve}(u[1])$
4. $R = Q_0 + Q_1$ # Point addition
5. $P = \text{clear_cofactor}(R)$
6. return P

Each hash-to-curve suite in [Section 8](#) instantiates one of these encoding functions for a specific elliptic curve.

3.1. Domain Separation Requirements

All uses of the encoding functions defined in this document **MUST** include domain separation ([Section 2.2.5](#)) to avoid interfering with other uses of similar functionality.

Applications that instantiate multiple, independent instances of either `hash_to_curve` or `encode_to_curve` **MUST** enforce domain separation between those instances. This requirement applies in both the case of multiple instances targeting the same curve and the case of multiple instances targeting different curves. (This is because the internal `hash_to_field` primitive ([Section 5](#)) requires domain separation to guarantee independent outputs.)

Domain separation is enforced with a domain separation tag (DST), which is a byte string constructed according to the following requirements:

1. Tags **MUST** be supplied as the DST parameter to `hash_to_field`, as described in [Section 5](#).
2. Tags **MUST** have nonzero length. A minimum length of 16 bytes is **RECOMMENDED** to reduce the chance of collisions with other applications.
3. Tags **SHOULD** begin with a fixed identification string that is unique to the application.
4. Tags **SHOULD** include a version number.
5. For applications that define multiple ciphersuites, each ciphersuite's tag **MUST** be different. For this purpose, it is **RECOMMENDED** to include a ciphersuite identifier in each tag.
6. For applications that use multiple encodings, to either the same curve or different curves, each encoding **MUST** use a different tag. For this purpose, it is **RECOMMENDED** to include the encoding's Suite ID ([Section 8](#)) in the domain separation tag. For independent encodings based on the same suite, each tag **SHOULD** also include a distinct identifier, e.g., "ENC1" and "ENC2".

As an example, consider a fictional application named Quux that defines several different ciphersuites, each for a different curve. A reasonable choice of tag is "QUUX-V<xx>-CS<yy>-<suiteID>", where <xx> and <yy> are two-digit numbers indicating the version and ciphersuite, respectively, and <suiteID> is the Suite ID of the encoding used in ciphersuite <yy>.

As another example, consider a fictional application named Baz that requires two independent random oracles to the same curve. Reasonable choices of tags for these oracles are "BAZ-V<xx>-CS<yy>-<suiteID>-ENC1" and "BAZ-V<xx>-CS<yy>-<suiteID>-ENC2", respectively, where <xx>, <yy>, and <suiteID> are as described above.

The example tags given above are assumed to be ASCII-encoded byte strings without null termination, which is the **RECOMMENDED** format. Other encodings can be used, but in all cases the encoding as a sequence of bytes **MUST** be specified unambiguously.

4. Utility Functions

Algorithms in this document use the utility functions described below, plus standard arithmetic operations (addition, multiplication, modular reduction, etc.) and elliptic curve point operations (point addition and scalar multiplication).

For security, implementations of these functions **SHOULD** be constant time: in brief, this means that execution time and memory access patterns **SHOULD NOT** depend on the values of secret inputs, intermediate values, or outputs. For such constant-time implementations, all arithmetic, comparisons, and assignments **MUST** also be implemented in constant time. [Section 10.3](#) briefly discusses constant-time security issues.

Guidance on implementing low-level operations (in constant time or otherwise) is beyond the scope of this document; readers should consult standard reference material [[MOV96](#)] [[CFADLNV05](#)].

- `CMOV(a, b, c)`: If `c` is `False`, `CMOV` returns `a`; otherwise, it returns `b`. For constant-time implementations, this operation must run in a time that is independent of the value of `c`.
- `AND`, `OR`, `NOT`, and `XOR` are standard bitwise logical operators. For constant-time implementations, short-circuit operators **MUST** be avoided.
- `is_square(x)`: This function returns `True` whenever the value `x` is a square in the field `F`. By Euler's criterion, this function can be calculated in constant time as

```
is_square(x) := { True,  if x^((q - 1) / 2) is 0 or 1 in F;  
                { False, otherwise.
```

In certain extension fields, `is_square` can be computed in constant time more quickly than by the above exponentiation. [[AR13](#)] and [[S85](#)] describe optimized methods for extension fields. [Appendix I.5](#) gives an optimized straight-line method for $GF(p^2)$.

- `sqrt(x)`: The `sqrt` operation is a multi-valued function, i.e., there exist two roots of `x` in the field `F` whenever `x` is square (except when `x = 0`). To maintain compatibility across implementations while allowing implementors leeway for optimizations, this document does not require `sqrt()` to return a particular value. Instead, as explained in [Section 6.4](#), any function that calls `sqrt` also specifies how to determine the correct root.

The preferred way of computing square roots is to fix a deterministic algorithm particular to `F`. We give several algorithms in [Appendix I](#).

- `sgn0(x)`: This function returns either 0 or 1 indicating the "sign" of `x`, where `sgn0(x) == 1` just when `x` is "negative". (In other words, this function always considers 0 to be positive.) [Section 4.1](#) defines this function and discusses its implementation.
- `inv0(x)`: This function returns the multiplicative inverse of `x` in `F`, extended to all of `F` by fixing `inv0(0) == 0`. A straightforward way to implement `inv0` in constant time is to compute

$$\text{inv0}(x) := x^{(q - 2)}.$$

Notice that on input 0, the output is 0 as required. Certain fields may allow faster inversion methods; detailed discussion of such methods is beyond the scope of this document.

- I2OSP and OS2IP: These functions are used to convert a byte string to and from a non-negative integer as described in [RFC8017]. (Note that these functions operate on byte strings in big-endian byte order.)
- `a || b`: denotes the concatenation of byte strings `a` and `b`. For example, `"ABC" || "DEF" == "ABCDEF"`.
- `substr(str, sbegin, slen)`: For a byte string `str`, this function returns the `slen`-byte substring starting at position `sbegin`; positions are zero indexed. For example, `substr("ABCDEFGH", 2, 3) == "CDE"`.
- `len(str)`: For a byte string `str`, this function returns the length of `str` in bytes. For example, `len("ABC") == 3`.
- `strxor(str1, str2)`: For byte strings `str1` and `str2`, `strxor(str1, str2)` returns the bitwise XOR of the two strings. For example, `strxor("abc", "XYZ") == "9;9"` (the strings in this example are ASCII literals, but `strxor` is defined for arbitrary byte strings). In this document, `strxor` is only applied to inputs of equal length.

4.1. The `sgn0` Function

This section defines a generic `sgn0` implementation that applies to any field $F = GF(p^m)$. It also gives simplified implementations for the cases $F = GF(p)$ and $F = GF(p^2)$.

The definition of the `sgn0` function for extension fields relies on the polynomial basis or vector representation of field elements, and iterates over the entire vector representation of the input element. As a result, `sgn0` depends on the primitive polynomial used to define the polynomial basis; see [Section 8](#) for more information about this basis, and see [Section 2.1](#) for a discussion of representing elements of extension fields as vectors.

```
sgn0(x)
```

Parameters:

- F , a finite field of characteristic p and order $q = p^m$.
- p , the characteristic of F (see immediately above).
- m , the extension degree of F , $m \geq 1$ (see immediately above).

Input: x , an element of F .

Output: 0 or 1.

Steps:

1. $sign = 0$
2. $zero = 1$
3. for i in $(1, 2, \dots, m)$:
4. $sign_i = x_i \bmod 2$
5. $zero_i = x_i == 0$
6. $sign = sign \text{ OR } (zero \text{ AND } sign_i)$ # Avoid short-circuit logic ops
7. $zero = zero \text{ AND } zero_i$
8. return $sign$

When $m == 1$, $sgn0$ can be significantly simplified:

```
sgn0_m_eq_1(x)
```

Input: x , an element of $GF(p)$.

Output: 0 or 1.

Steps:

1. return $x \bmod 2$

The case $m == 2$ is only slightly more complicated:

```
sgn0_m_eq_2(x)
```

Input: x , an element of $GF(p^2)$.

Output: 0 or 1.

Steps:

1. $sign_0 = x_0 \bmod 2$
2. $zero_0 = x_0 == 0$
3. $sign_1 = x_1 \bmod 2$
4. $s = sign_0 \text{ OR } (zero_0 \text{ AND } sign_1)$ # Avoid short-circuit logic ops
5. return s

5. Hashing to a Finite Field

The `hash_to_field` function hashes a byte string `msg` of arbitrary length into one or more elements of a field F . This function works in two steps: it first hashes the input byte string to produce a uniformly random byte string, and then interprets this byte string as one or more elements of F .

For the first step, `hash_to_field` calls an auxiliary function `expand_message`. This document defines two variants of `expand_message`: one appropriate for hash functions like SHA-2 [FIPS180-4] or SHA-3 [FIPS202], and another appropriate for extendable-output functions such as SHAKE128 [FIPS202]. Security considerations for each `expand_message` variant are discussed below (Sections 5.3.1 and 5.3.2).

Implementors **MUST NOT** use rejection sampling to generate a uniformly random element of F , to ensure that the `hash_to_field` function is amenable to constant-time implementation. The reason is that rejection sampling procedures are difficult to implement in constant time, and later well-meaning "optimizations" may silently render an implementation non-constant-time. This means that any `hash_to_field` function based on rejection sampling would be incompatible with constant-time implementation.

The `hash_to_field` function is also suitable for securely hashing to scalars. For example, when hashing to the scalar field for an elliptic curve (sub)group with prime order r , it suffices to instantiate `hash_to_field` with target field $GF(r)$.

The `hash_to_field` function is designed to be indiffereniable from a random oracle [MRH04] when `expand_message` (Section 5.3) is modeled as a random oracle (see Section 10.5 for details about its indiffereniability). Ensuring indiffereniability requires care; to see why, consider a prime p that is close to $3/4 * 2^{256}$. Reducing a random 256-bit integer modulo this p yields a value that is in the range $[0, p/3]$ with probability roughly $1/2$, meaning that this value is statistically far from uniform in $[0, p-1]$.

To control bias, `hash_to_field` instead uses random integers whose length is at least $\text{ceil}(\log_2(p)) + k$ bits, where k is the target security level for the suite in bits. Reducing such integers mod p gives bias at most 2^{-k} for any p ; this bias is appropriate when targeting k -bit security. For each such integer, `hash_to_field` uses `expand_message` to obtain L uniform bytes, where

$$L = \text{ceil}((\text{ceil}(\log_2(p)) + k) / 8)$$

These uniform bytes are then interpreted as an integer via OS2IP. For example, for a 255-bit prime p , and $k = 128$ -bit security, $L = \text{ceil}((255 + 128) / 8) = 48$ bytes.

Note that k is an upper bound on the security level for the corresponding curve. See Section 10.8 for more details and Section 8.9 for guidelines on choosing k for a given curve.

5.1. Efficiency Considerations in Extension Fields

The `hash_to_field` function described in this section is inefficient for certain extension fields. Specifically, when hashing to an element of the extension field $GF(p^m)$, `hash_to_field` requires expanding `msg` into $m * L$ bytes (for L as defined above). For extension fields where $\log_2(p)$ is significantly smaller than the security level k , this approach is inefficient: it requires

`expand_message` to output roughly $m * \log_2(p) + m * k$ bits, whereas $m * \log_2(p) + k$ bytes suffices to generate an element of $GF(p^m)$ with bias at most 2^{-k} . In such cases, applications **MAY** use an alternative `hash_to_field` function, provided it meets the following security requirements:

- The function **MUST** output one or more field elements that are uniformly random except with bias at most 2^{-k} .
- The function **MUST NOT** use rejection sampling.
- The function **SHOULD** be amenable to straight-line implementations.

For example, Pornin [P20] describes a method for hashing to $GF(9767^{19})$ that meets these requirements while using fewer output bits from `expand_message` than `hash_to_field` would for that field.

5.2. `hash_to_field` Implementation

The following procedure implements `hash_to_field`.

The `expand_message` parameter to this function **MUST** conform to the requirements given in [Section 5.3](#). [Section 3.1](#) discusses the **REQUIRED** method for constructing DST, the domain separation tag. Note that `hash_to_field` may fail (ABORT) if `expand_message` fails.

```
hash_to_field(msg, count)
```

Parameters:

- DST, a domain separation tag (see [Section 3.1](#)).
- F, a finite field of characteristic p and order $q = p^m$.
- p , the characteristic of F (see immediately above).
- m , the extension degree of F, $m \geq 1$ (see immediately above).
- $L = \lceil (\lceil \log_2(p) \rceil + k) / 8 \rceil$, where k is the security parameter of the suite (e.g., $k = 128$).
- `expand_message`, a function that expands a byte string and domain separation tag into a uniformly random byte string (see [Section 5.3](#)).

Input:

- `msg`, a byte string containing the message to hash.
- `count`, the number of elements of F to output.

Output:

- $(u_0, \dots, u_{(count - 1)})$, a list of field elements.

Steps:

1. `len_in_bytes = count * m * L`
2. `uniform_bytes = expand_message(msg, DST, len_in_bytes)`
3. `for i in (0, ..., count - 1):`
4. `for j in (0, ..., m - 1):`
5. `elm_offset = L * (j + i * m)`
6. `tv = substr(uniform_bytes, elm_offset, L)`
7. `e_j = OS2IP(tv) mod p`
8. `u_i = (e_0, ..., e_{(m - 1)})`
9. `return (u_0, ..., u_{(count - 1)})`

5.3. expand_message

expand_message is a function that generates a uniformly random byte string. It takes three arguments:

1. msg, a byte string containing the message to hash,
2. DST, a byte string that acts as a domain separation tag, and
3. len_in_bytes, the number of bytes to be generated.

This document defines the following two variants of expand_message:

- expand_message_xmd ([Section 5.3.1](#)) is appropriate for use with a wide range of hash functions, including SHA-2 [[FIPS180-4](#)], SHA-3 [[FIPS202](#)], BLAKE2 [[RFC7693](#)], and others.
- expand_message_xof ([Section 5.3.2](#)) is appropriate for use with extendable-output functions (XOFs), including functions in the SHAKE [[FIPS202](#)] or BLAKE2X [[BLAKE2X](#)] families.

These variants should suffice for the vast majority of use cases, but other variants are possible; [Section 5.3.4](#) discusses requirements.

5.3.1. expand_message_xmd

The expand_message_xmd function produces a uniformly random byte string using a cryptographic hash function H that outputs b bits. For security, H **MUST** meet the following requirements:

- The number of bits output by H **MUST** be $b \geq 2 * k$, where k is the target security level in bits, and b **MUST** be divisible by 8. The first requirement ensures k-bit collision resistance; the second ensures uniformity of expand_message_xmd's output.
- H **MAY** be a Merkle-Damgaard hash function like SHA-2. In this case, security holds when the underlying compression function is modeled as a random oracle [[CDMP05](#)]. (See [Section 10.6](#) for discussion.)
- H **MAY** be a sponge-based hash function like SHA-3 or BLAKE2. In this case, security holds when the inner function is modeled as a random transformation or as a random permutation [[BDPV08](#)].
- Otherwise, H **MUST** be a hash function that has been proved indistinguishable from a random oracle [[MRH04](#)] under a reasonable cryptographic assumption.

SHA-2 [[FIPS180-4](#)] and SHA-3 [[FIPS202](#)] are typical and **RECOMMENDED** choices. As an example, for the 128-bit security level, $b \geq 256$ bits and either SHA-256 or SHA3-256 would be an appropriate choice.

The hash function H is assumed to work by repeatedly ingesting fixed-length blocks of data. The length in bits of these blocks is called the input block size (s). As examples, $s = 1024$ for SHA-512 [[FIPS180-4](#)] and $s = 576$ for SHA3-512 [[FIPS202](#)]. For correctness, H requires $b \leq s$.

The following procedure implements expand_message_xmd.

```
expand_message_xmd(msg, DST, len_in_bytes)
```

Parameters:

- H, a hash function (see requirements above).
- b_in_bytes, $b / 8$ for b the output size of H in bits.
For example, for $b = 256$, b_in_bytes = 32.
- s_in_bytes, the input block size of H, measured in bytes (see discussion above). For example, for SHA-256, s_in_bytes = 64.

Input:

- msg, a byte string.
- DST, a byte string of at most 255 bytes.
See below for information on using longer DSTs.
- len_in_bytes, the length of the requested output in bytes,
not greater than the lesser of $(255 * b_in_bytes)$ or $2^{16}-1$.

Output:

- uniform_bytes, a byte string.

Steps:

1. ell = ceil(len_in_bytes / b_in_bytes)
2. ABORT if ell > 255 or len_in_bytes > 65535 or len(DST) > 255
3. DST_prime = DST || I2OSP(len(DST), 1)
4. Z_pad = I2OSP(0, s_in_bytes)
5. l_i_b_str = I2OSP(len_in_bytes, 2)
6. msg_prime = Z_pad || msg || l_i_b_str || I2OSP(0, 1) || DST_prime
7. b_0 = H(msg_prime)
8. b_1 = H(b_0 || I2OSP(1, 1) || DST_prime)
9. for i in (2, ..., ell):
10. b_i = H(strxor(b_0, b_(i - 1)) || I2OSP(i, 1) || DST_prime)
11. uniform_bytes = b_1 || ... || b_ell
12. return substr(uniform_bytes, 0, len_in_bytes)

Note that the string Z_pad (step 6) is prefixed to msg before computing b_0 (step 7). This is necessary for security when H is a Merkle-Damgaard hash, e.g., SHA-2 (see [Section 10.6](#)). Hashing this additional data means that the cost of computing b_0 is higher than the cost of simply computing H(msg). In most settings, this overhead is negligible, because the cost of evaluating H is much less than the other costs involved in hashing to a curve.

It is possible, however, to entirely avoid this overhead by taking advantage of the fact that Z_pad depends only on H, and not on the arguments to expand_message_xmd. To do so, first precompute and save the internal state of H after ingesting Z_pad. Then, when computing b_0, initialize H using the saved state. Further details are implementation dependent and are beyond the scope of this document.

5.3.2. expand_message_xof

The expand_message_xof function produces a uniformly random byte string using an extendable-output function (XOF) H. For security, H **MUST** meet the following criteria:

- The collision resistance of H **MUST** be at least k bits.

- **H MUST** be an XOF that has been proved indiffereniable from a random oracle under a reasonable cryptographic assumption.

The SHAKE XOF family [FIPS202] is a typical and **RECOMMENDED** choice. As an example, for 128-bit security, SHAKE128 would be an appropriate choice.

The following procedure implements `expand_message_xof`.

```
expand_message_xof(msg, DST, len_in_bytes)
```

Parameters:

- $H(m, d)$, an extendable-output function that processes input message m and returns d bytes.

Input:

- `msg`, a byte string.
- `DST`, a byte string of at most 255 bytes.
See below for information on using longer DSTs.
- `len_in_bytes`, the length of the requested output in bytes.

Output:

- `uniform_bytes`, a byte string.

Steps:

1. ABORT if `len_in_bytes > 65535` or `len(DST) > 255`
2. `DST_prime = DST || I2OSP(len(DST), 1)`
3. `msg_prime = msg || I2OSP(len_in_bytes, 2) || DST_prime`
4. `uniform_bytes = H(msg_prime, len_in_bytes)`
5. return `uniform_bytes`

5.3.3. Using DSTs Longer Than 255 Bytes

The `expand_message` variants defined in this section accept domain separation tags of at most 255 bytes. If applications require a domain separation tag longer than 255 bytes, e.g., because of requirements imposed by an invoking protocol, implementors **MUST** compute a short domain separation tag by hashing, as follows:

- For `expand_message_xmd` using hash function H , DST is computed as

```
DST = H("H2C-OVERSIZE-DST-" || a_very_long_DST)
```

- For `expand_message_xof` using extendable-output function H , DST is computed as

```
DST = H("H2C-OVERSIZE-DST-" || a_very_long_DST, ceil(2 * k / 8))
```

Here, `a_very_long_DST` is the DST whose length is greater than 255 bytes, "H2C-OVERSIZE-DST-" is a 17-byte ASCII string literal, and k is the target security level in bits.

5.3.4. Defining Other `expand_message` Variants

When defining a new `expand_message` variant, the most important consideration is that `hash_to_field` models `expand_message` as a random oracle. Thus, implementors **SHOULD** prove indistinguishability from a random oracle under an appropriate assumption about the underlying cryptographic primitives; see [Section 10.5](#) for more information.

In addition, `expand_message` variants:

- **MUST** give collision resistance commensurate with the security level of the target elliptic curve.
- **MUST** be built on primitives designed for use in applications requiring cryptographic randomness. As examples, a secure stream cipher is an appropriate primitive, whereas a Mersenne twister pseudorandom number generator [MT98] is not.
- **MUST NOT** use rejection sampling.
- **MUST** give independent values for distinct (`msg`, `DST`, `length`) inputs. Meeting this requirement is subtle. As a simplified example, hashing `msg || DST` does not work, because in this case distinct (`msg`, `DST`) pairs whose concatenations are equal will return the same output (e.g., ("AB", "CDEF") and ("ABC", "DEF")). The variants defined in this document use a suffix-free encoding of `DST` to avoid this issue.
- **MUST** use the domain separation tag `DST` to ensure that invocations of cryptographic primitives inside of `expand_message` are domain-separated from invocations outside of `expand_message`. For example, if the `expand_message` variant uses a hash function `H`, an encoding of `DST` **MUST** be added either as a prefix or a suffix of the input to each invocation of `H`. Adding `DST` as a suffix is the **RECOMMENDED** approach.
- **SHOULD** read `msg` exactly once, for efficiency when `msg` is long.

In addition, each `expand_message` variant **MUST** specify a unique `EXP_TAG` that identifies that variant in a Suite ID. See [Section 8.10](#) for more information.

6. Deterministic Mappings

The mappings in this section are suitable for implementing either nonuniform or uniform encodings using the constructions in [Section 3](#). Certain mappings restrict the form of the curve or its parameters. For each mapping presented, this document lists the relevant restrictions.

Note that mappings in this section are not interchangeable: different mappings will almost certainly output different points when evaluated on the same input.

6.1. Choosing a Mapping Function

This section gives brief guidelines on choosing a mapping function for a given elliptic curve. Note that the suites given in [Section 8](#) are recommended mappings for the respective curves.

If the target elliptic curve is a Montgomery curve (Section 6.7), the Elligator 2 method (Section 6.7.1) is recommended. Similarly, if the target elliptic curve is a twisted Edwards curve (Section 6.8), the twisted Edwards Elligator 2 method (Section 6.8.2) is recommended.

The remaining cases are Weierstrass curves. For curves supported by the Simplified Shallue-van de Woestijne-Ulas (SWU) method (Section 6.6.2), that mapping is the recommended one. Otherwise, the Simplified SWU method for $AB \neq 0$ (Section 6.6.3) is recommended if the goal is best performance, while the Shallue-van de Woestijne method (Section 6.6.1) is recommended if the goal is simplicity of implementation. (The reason for this distinction is that the Simplified SWU method for $AB \neq 0$ requires implementing an isogeny map in addition to the mapping function, while the Shallue-van de Woestijne method does not.)

The Shallue-van de Woestijne method (Section 6.6.1) works with any curve and may be used in cases where a generic mapping is required. Note, however, that this mapping is almost always more computationally expensive than the curve-specific recommendations above.

6.2. Interface

The generic interface shared by all mappings in this section is as follows:

```
(x, y) = map_to_curve(u)
```

The input u and outputs x and y are elements of the field F . The affine coordinates (x, y) specify a point on an elliptic curve defined over F . Note, however, that the point (x, y) is not a uniformly random point.

6.3. Notation

As a rough guide, the following conventions are used in pseudocode:

- All arithmetic operations are performed over a field F , unless explicitly stated otherwise.
- u : the input to the mapping function. This is an element of F produced by the `hash_to_field` function.
- (x, y) , (s, t) , (v, w) : the affine coordinates of the point output by the mapping. Indexed variables (e.g., x_1, y_2, \dots) are used for candidate values.
- tv_1, tv_2, \dots : reusable temporary variables.
- c_1, c_2, \dots : constant values, which can be computed in advance.

6.4. Sign of the Resulting Point

In general, elliptic curves have equations of the form $y^2 = g(x)$. The mappings in this section first identify an x such that $g(x)$ is square, then take a square root to find y . Since there are two square roots when $g(x) \neq 0$, this may result in an ambiguity regarding the sign of y .

When necessary, the mappings in this section resolve this ambiguity by specifying the sign of the y-coordinate in terms of the input to the mapping function. Two main reasons support this approach: first, this covers elliptic curves over any field in a uniform way, and second, it gives implementors leeway in optimizing square-root implementations.

6.5. Exceptional Cases

Mappings may have exceptional cases, i.e., inputs u on which the mapping is undefined. These cases must be handled carefully, especially for constant-time implementations.

For each mapping in this section, we discuss the exceptional cases and show how to handle them in constant time. Note that all implementations **SHOULD** use `inv0` (Section 4) to compute multiplicative inverses, to avoid exceptional cases that result from attempting to compute the inverse of 0.

6.6. Mappings for Weierstrass Curves

The mappings in this section apply to a target curve E defined by the equation

$$y^2 = g(x) = x^3 + A * x + B$$

where $4 * A^3 + 27 * B^2 \neq 0$.

6.6.1. Shallue-van de Woestijne Method

Shallue and van de Woestijne [SW06] describe a mapping that applies to essentially any elliptic curve. (Note, however, that this mapping is more expensive to evaluate than the other mappings in this document.)

The parameterization given below is for Weierstrass curves; its derivation is detailed in [W19]. This parameterization also works for Montgomery curves (Section 6.7) and twisted Edwards curves (Section 6.8) via the rational maps given in Appendix D: first, evaluate the Shallue-van de Woestijne mapping to an equivalent Weierstrass curve, then map that point to the target Montgomery or twisted Edwards curve using the corresponding rational map.

Preconditions: A Weierstrass curve $y^2 = x^3 + A * x + B$.

Constants:

- A and B , the parameter of the Weierstrass curve.
- Z , a non-zero element of F meeting the below criteria. Appendix H.1 gives a Sage script [SAGE] that outputs the **RECOMMENDED** Z .
 1. $g(Z) \neq 0$ in F .
 2. $-(3 * Z^2 + 4 * A) / (4 * g(Z)) \neq 0$ in F .
 3. $-(3 * Z^2 + 4 * A) / (4 * g(Z))$ is square in F .
 4. At least one of $g(Z)$ and $g(-Z / 2)$ is square in F .

Sign of y : Inputs u and $-u$ give the same x -coordinate for many values of u . Thus, we set $\text{sgn0}(y) == \text{sgn0}(u)$.

Exceptions: The exceptional cases for u occur when $(1 + u^2 * g(Z)) * (1 - u^2 * g(Z)) == 0$. The restrictions on Z given above ensure that implementations that use inv0 to invert this product are exception free.

Operations:

```

1. tv1 = u^2 * g(Z)
2. tv2 = 1 + tv1
3. tv1 = 1 - tv1
4. tv3 = inv0(tv1 * tv2)
5. tv4 = sqrt(-g(Z) * (3 * Z^2 + 4 * A)) # can be precomputed
6. If sgn0(tv4) == 1, set tv4 = -tv4 # sgn0(tv4) MUST equal 0
7. tv5 = u * tv1 * tv3 * tv4
8. tv6 = -4 * g(Z) / (3 * Z^2 + 4 * A) # can be precomputed
9. x1 = -Z / 2 - tv5
10. x2 = -Z / 2 + tv5
11. x3 = Z + tv6 * (tv2^2 * tv3)^2
12. If is_square(g(x1)), set x = x1 and y = sqrt(g(x1))
13. Else If is_square(g(x2)), set x = x2 and y = sqrt(g(x2))
14. Else set x = x3 and y = sqrt(g(x3))
15. If sgn0(u) != sgn0(y), set y = -y
16. return (x, y)

```

[Appendix F.1](#) gives an example straight-line implementation of this mapping.

6.6.2. Simplified Shallue-van de Woestijne-Ulas Method

The function `map_to_curve_simple_swu(u)` implements a simplification of the Shallue-van de Woestijne-Ulas mapping [U07] described by Brier et al. [BCIMRT10], which they call the "simplified SWU" map. Wahby and Boneh [WB19] generalize and optimize this mapping.

Preconditions: A Weierstrass curve $y^2 = x^3 + A * x + B$ where $A \neq 0$ and $B \neq 0$.

Constants:

- A and B , the parameters of the Weierstrass curve.
- Z , an element of F meeting the below criteria. [Appendix H.2](#) gives a Sage script [SAGE] that outputs the **RECOMMENDED** Z . The criteria are as follows:
 1. Z is non-square in F ,
 2. $Z \neq -1$ in F ,
 3. the polynomial $g(x) - Z$ is irreducible over F , and
 4. $g(B / (Z * A))$ is square in F .

Sign of y : Inputs u and $-u$ give the same x -coordinate. Thus, we set $\text{sgn0}(y) == \text{sgn0}(u)$.

Exceptions: The exceptional cases are values of u such that $Z^2 * u^4 + Z * u^2 == 0$. This includes $u == 0$ and may include other values that depend on Z . Implementations must detect this case and set $x_1 = B / (Z * A)$, which guarantees that $g(x_1)$ is square by the condition on Z given above.

Operations:

```

1. tv1 = inv0(Z^2 * u^4 + Z * u^2)
2. x1 = (-B / A) * (1 + tv1)
3. If tv1 == 0, set x1 = B / (Z * A)
4. gx1 = x1^3 + A * x1 + B
5. x2 = Z * u^2 * x1
6. gx2 = x2^3 + A * x2 + B
7. If is_square(gx1), set x = x1 and y = sqrt(gx1)
8. Else set x = x2 and y = sqrt(gx2)
9. If sgn0(u) != sgn0(y), set y = -y
10. return (x, y)

```

[Appendix F.2](#) gives a general and optimized straight-line implementation of this mapping. For more information on optimizing this mapping, see Section 4 of [\[WB19\]](#) or the example code found at [\[hash2curve-repo\]](#).

6.6.3. Simplified SWU for $AB == 0$

Wahby and Boneh [\[WB19\]](#) show how to adapt the Simplified SWU mapping to Weierstrass curves having $A == 0$ or $B == 0$, which the mapping of [Section 6.6.2](#) does not support. (The case $A == B == 0$ is excluded because $y^2 = x^3$ is not an elliptic curve.)

This method applies to curves like secp256k1 [\[SEC2\]](#) and to pairing-friendly curves in the Barreto-Lynn-Scott family [\[BLS03\]](#), Barreto-Naehrig family [\[BN05\]](#), and other families.

This method requires finding another elliptic curve E' given by the equation

$$y'^2 = g'(x') = x'^3 + A' * x' + B'$$

that is isogenous to E and has $A' != 0$ and $B' != 0$. (See [\[WB19\]](#), Appendix A, for one way of finding E' using [\[SAGE\]](#).) This isogeny defines a map $\text{iso_map}(x', y')$ given by a pair of rational functions. iso_map takes as input a point on E' and produces as output a point on E .

Once E' and iso_map are identified, this mapping works as follows: on input u , first apply the Simplified SWU mapping to get a point on E' , then apply the isogeny map to that point to get a point on E .

Note that iso_map is a group homomorphism, meaning that point addition commutes with iso_map . Thus, when using this mapping in the hash_to_curve construction discussed in [Section 3](#), one can effect a small optimization by first mapping u_0 and u_1 to E' , adding the resulting points on E' , and then applying iso_map to the sum. This gives the same result while requiring only one evaluation of iso_map .

Preconditions: An elliptic curve E' with $A' \neq 0$ and $B' \neq 0$ that is isogenous to the target curve E with isogeny map `iso_map` from E' to E .

Helper functions:

- `map_to_curve_simple_swu` is the mapping of [Section 6.6.2](#) to E'
- `iso_map` is the isogeny map from E' to E

Sign of y : For this map, the sign is determined by `map_to_curve_simple_swu`. No further sign adjustments are necessary.

Exceptions: `map_to_curve_simple_swu` handles its exceptional cases. Exceptional cases of `iso_map` are inputs that cause the denominator of either rational function to evaluate to zero; such cases **MUST** return the identity point on E .

Operations:

```

1. (x', y') = map_to_curve_simple_swu(u)    # (x', y') is on E'
2.  (x, y) = iso_map(x', y')              # (x, y) is on E
3. return (x, y)

```

See [\[hash2curve-repo\]](#) or Section 4.3 of [\[WB19\]](#) for details on implementing the isogeny map.

6.7. Mappings for Montgomery Curves

The mapping defined in this section applies to a target curve M defined by the equation

$$K * t^2 = s^3 + J * s^2 + s$$

6.7.1. Elligator 2 Method

Bernstein, Hamburg, Krasnova, and Lange give a mapping that applies to any curve with a point of order 2 [\[BHKL13\]](#), which they call Elligator 2.

Preconditions: A Montgomery curve $K * t^2 = s^3 + J * s^2 + s$ where $J \neq 0$, $K \neq 0$, and $(J^2 - 4) / K^2$ is non-zero and non-square in F .

Constants:

- J and K , the parameters of the elliptic curve.
- Z , a non-square element of F . [Appendix H.3](#) gives a Sage script [\[SAGE\]](#) that outputs the **RECOMMENDED** Z .

Sign of t : This mapping fixes the sign of t as specified in [\[BHKL13\]](#). No additional adjustment is required.

Exceptions: The exceptional case is $Z * u^2 == -1$, i.e., $1 + Z * u^2 == 0$. Implementations must detect this case and set $x_1 = -(J / K)$. Note that this can only happen when $q = 3 \pmod{4}$.

Operations:

```

1.  $x_1 = -(J / K) * \text{inv}\theta(1 + Z * u^2)$ 
2. If  $x_1 == 0$ , set  $x_1 = -(J / K)$ 
3.  $gx_1 = x_1^3 + (J / K) * x_1^2 + x_1 / K^2$ 
4.  $x_2 = -x_1 - (J / K)$ 
5.  $gx_2 = x_2^3 + (J / K) * x_2^2 + x_2 / K^2$ 
6. If  $\text{is\_square}(gx_1)$ , set  $x = x_1$ ,  $y = \text{sqrt}(gx_1)$  with  $\text{sgn}\theta(y) == 1$ .
7. Else set  $x = x_2$ ,  $y = \text{sqrt}(gx_2)$  with  $\text{sgn}\theta(y) == 0$ .
8.  $s = x * K$ 
9.  $t = y * K$ 
10. return (s, t)

```

[Appendix F.3](#) gives an example straight-line implementation of this mapping. [Appendix G.2](#) gives optimized straight-line procedures that apply to specific classes of curves and base fields.

6.8. Mappings for Twisted Edwards Curves

Twisted Edwards curves (a class of curves that includes Edwards curves) are given by the equation

$$a * v^2 + w^2 = 1 + d * v^2 * w^2$$

with $a \neq 0$, $d \neq 0$, and $a \neq d$ [\[BBJLP08\]](#).

These curves are closely related to Montgomery curves ([Section 6.7](#)): every twisted Edwards curve is birationally equivalent to a Montgomery curve ([\[BBJLP08\]](#), Theorem 3.2). This equivalence yields an efficient way of hashing to a twisted Edwards curve: first, hash to an equivalent Montgomery curve, then transform the result into a point on the twisted Edwards curve via a rational map. This method of hashing to a twisted Edwards curve thus requires identifying a corresponding Montgomery curve and rational map. We describe how to identify such a curve and map immediately below.

6.8.1. Rational Maps from Montgomery to Twisted Edwards Curves

There are two ways to select a Montgomery curve and rational map for use when hashing to a given twisted Edwards curve. The selected Montgomery curve and rational map **MUST** be specified as part of the hash-to-curve suite for a given twisted Edwards curve; see [Section 8](#).

1. When hashing to a standardized twisted Edwards curve for which a corresponding Montgomery form and rational map are also standardized, the standard Montgomery form and rational map **SHOULD** be used to ensure compatibility with existing software.

In certain cases, e.g., `edwards25519` [\[RFC7748\]](#), the sign of the rational map from the twisted Edwards curve to its corresponding Montgomery curve is not given explicitly. In this case, the sign **MUST** be fixed such that applying the rational map to the twisted Edwards curve's base point yields the Montgomery curve's base point with correct sign. (For `edwards25519`, see [\[RFC7748\]](#) and [\[Err4730\]](#).)

When defining new twisted Edwards curves, a Montgomery equivalent and rational map **SHOULD** also be specified, and the sign of the rational map **SHOULD** be stated explicitly.

2. When hashing to a twisted Edwards curve that does not have a standardized Montgomery form or rational map, the map given in [Appendix D](#) **SHOULD** be used.

6.8.2. Elligator 2 Method

Preconditions: A twisted Edwards curve E and an equivalent Montgomery curve M meeting the requirements in [Section 6.8.1](#).

Helper functions:

- `map_to_curve_elligator2` is the mapping of [Section 6.7.1](#) to the curve M .
- `rational_map` is a function that takes a point (s, t) on M and returns a point (v, w) on E . This rational map should be chosen as defined in [Section 6.8.1](#).

Sign of t (and v): For this map, the sign is determined by `map_to_curve_elligator2`. No further sign adjustments are required.

Exceptions: The exceptions for the Elligator 2 mapping are as given in [Section 6.7.1](#). The exceptions for the rational map are as given in [Section 6.8.1](#). No other exceptions are possible.

The following procedure implements the Elligator 2 mapping for a twisted Edwards curve. (Note that the output point is denoted (v, w) because it is a point on the target twisted Edwards curve.)

```
map_to_curve_elligator2_edwards(u)
```

```
Input: u, an element of F.
```

```
Output: (v, w), a point on E.
```

```
1. (s, t) = map_to_curve_elligator2(u)      # (s, t) is on M
2. (v, w) = rational_map(s, t)             # (v, w) is on E
3. return (v, w)
```

7. Clearing the Cofactor

The mappings of [Section 6](#) always output a point on the elliptic curve, i.e., a point in a group of order $h * r$ ([Section 2.1](#)). Obtaining a point in G may require a final operation commonly called "clearing the cofactor," which takes as input any point on the curve and produces as output a point in the prime-order (sub)group G ([Section 2.1](#)).

The cofactor can always be cleared via scalar multiplication by h . For elliptic curves where $h = 1$, i.e., the curves with a prime number of points, no operation is required. This applies, for example, to the NIST curves P-256, P-384, and P-521 [[FIPS186-4](#)].

In some cases, it is possible to clear the cofactor via a faster method than scalar multiplication by h . These methods are equivalent to (but usually faster than) multiplication by some scalar h_{eff} whose value is determined by the method and the curve. Examples of fast cofactor clearing methods include the following:

- For certain pairing-friendly curves having subgroup G_2 over an extension field, Scott et al. [SBCDK09] describe a method for fast cofactor clearing that exploits an efficiently computable endomorphism. Fuentes-Castaneda et al. [FKR11] propose an alternative method that is sometimes more efficient. Budroni and Pintore [BP17] give concrete instantiations of these methods for Barreto-Lynn-Scott pairing-friendly curves [BLS03]. This method is described for the specific case of BLS12-381 in [Appendix G.3](#).
- Wahby and Boneh ([WB19], Section 5) describe a trick due to Scott for fast cofactor clearing on any elliptic curve for which the prime factorization of h and the structure of the elliptic curve group meet certain conditions.

The `clear_cofactor` function is parameterized by a scalar h_{eff} . Specifically,

```
clear_cofactor(P) := h_eff * P
```

where $*$ represents scalar multiplication. When a curve does not support a fast cofactor clearing method, $h_{\text{eff}} = h$ and the cofactor **MUST** be cleared via scalar multiplication.

When a curve admits a fast cofactor clearing method, `clear_cofactor` **MAY** be evaluated either via that method or via scalar multiplication by the equivalent h_{eff} ; these two methods give the same result. Note that in this case scalar multiplication by the cofactor h does not generally give the same result as the fast method and **MUST NOT** be used.

8. Suites for Hashing

This section lists recommended suites for hashing to standard elliptic curves.

A hash-to-curve suite fully specifies the procedure for hashing byte strings to points on a specific elliptic curve group. [Section 8.1](#) describes how to implement a suite. Applications that require hashing to an elliptic curve should use either an existing suite or a new suite specified as described in [Section 8.9](#).

All applications using a hash-to-curve suite **MUST** choose a domain separation tag (DST) in accordance with the guidelines in [Section 3.1](#). In addition, applications whose security requires a random oracle that returns uniformly random points on the target curve **MUST** use a suite whose encoding type is `hash_to_curve`; see [Section 3](#) and immediately below for more information.

A hash-to-curve suite comprises the following parameters:

- Suite ID, a short name used to refer to a given suite. [Section 8.10](#) discusses the naming conventions for Suite IDs.

- encoding type, either uniform (`hash_to_curve`) or nonuniform (`encode_to_curve`). See [Section 3](#) for definitions of these encoding types.
- E , the target elliptic curve over a field F .
- p , the characteristic of the field F .
- m , the extension degree of the field F . If $m > 1$, the suite **MUST** also specify the polynomial basis used to represent extension field elements.
- k , the target security level of the suite in bits. (See [Section 10.8](#) for discussion.)
- L , the length parameter for `hash_to_field` ([Section 5](#)).
- `expand_message`, one of the variants specified in [Section 5.3](#) plus any parameters required for the specified variant (for example, H , the underlying hash function).
- f , a mapping function from [Section 6](#).
- h_{eff} , the scalar parameter for `clear_cofactor` ([Section 7](#)).

In addition to the above parameters, the mapping f may require additional parameters Z , M , `rational_map`, E' , or `iso_map`. When applicable, these **MUST** be specified.

The table below lists suites **RECOMMENDED** for some elliptic curves. The corresponding parameters are given in the following subsections. Applications instantiating cryptographic protocols whose security analysis relies on a random oracle that outputs points with a uniform distribution **MUST NOT** use a nonuniform encoding. Moreover, applications that use a nonuniform encoding **SHOULD** carefully analyze the security implications of nonuniformity. When the required encoding is not clear, applications **SHOULD** use a uniform encoding for security.

E	Suites	Section
NIST P-256	P256_XMD:SHA-256_SSWU_RO_ P256_XMD:SHA-256_SSWU_NU_	8.2
NIST P-384	P384_XMD:SHA-384_SSWU_RO_ P384_XMD:SHA-384_SSWU_NU_	8.3
NIST P-521	P521_XMD:SHA-512_SSWU_RO_ P521_XMD:SHA-512_SSWU_NU_	8.4
curve25519	curve25519_XMD:SHA-512_ELL2_RO_ curve25519_XMD:SHA-512_ELL2_NU_	8.5
edwards25519	edwards25519_XMD:SHA-512_ELL2_RO_ edwards25519_XMD:SHA-512_ELL2_NU_	8.5
curve448	curve448_XOF:SHAKE256_ELL2_RO_ curve448_XOF:SHAKE256_ELL2_NU_	8.6
edwards448	edwards448_XOF:SHAKE256_ELL2_RO_ edwards448_XOF:SHAKE256_ELL2_NU_	8.6
secp256k1	secp256k1_XMD:SHA-256_SSWU_RO_ secp256k1_XMD:SHA-256_SSWU_NU_	8.7

E	Suites	Section
BLS12-381 G1	BLS12381G1_XMD:SHA-256_SSWU_RO_ BLS12381G1_XMD:SHA-256_SSWU_NU_	8.8
BLS12-381 G2	BLS12381G2_XMD:SHA-256_SSWU_RO_ BLS12381G2_XMD:SHA-256_SSWU_NU_	8.8

Table 2: Suites for hashing to elliptic curves.

8.1. Implementing a Hash-to-Curve Suite

A hash-to-curve suite requires the following functions. Note that some of these require utility functions from [Section 4](#).

1. Base field arithmetic operations for the target elliptic curve, e.g., addition, multiplication, and square root.
2. Elliptic curve point operations for the target curve, e.g., point addition and scalar multiplication.
3. The `hash_to_field` function; see [Section 5](#). This includes the `expand_message` variant ([Section 5.3](#)) and any constituent hash function or XOF.
4. The suite-specified mapping function; see the corresponding subsection of [Section 6](#).
5. A cofactor clearing function; see [Section 7](#). This may be implemented as scalar multiplication by `h_eff` or as a faster equivalent method.
6. The desired encoding function; see [Section 3](#). This is either `hash_to_curve` or `encode_to_curve`.

8.2. Suites for NIST P-256

This section defines ciphersuites for the NIST P-256 elliptic curve [[FIPS186-4](#)].

`P256_XMD:SHA-256_SSWU_RO_` is defined as follows:

- encoding type: `hash_to_curve` ([Section 3](#))
- E: $y^2 = x^3 + A * x + B$, where
 - $A = -3$
 - $B = 0x5ac635d8aa3a93e7b3ebbd55769886bc651d06b0cc53b0f63bce3c3e27d2604b$
- p: $2^{256} - 2^{224} + 2^{192} + 2^{96} - 1$
- m: 1
- k: 128
- `expand_message`: `expand_message_xmd` ([Section 5.3.1](#))
- H: SHA-256
- L: 48
- f: Simplified SWU method ([Section 6.6.2](#))

- Z: -10
- h_eff: 1

P256_XMD:SHA-256_SSWU_NU_ is identical to P256_XMD:SHA-256_SSWU_RO_, except that the encoding type is `encode_to_curve` (Section 3).

An optimized example implementation of the Simplified SWU mapping to P-256 is given in Appendix F.2.

8.3. Suites for NIST P-384

This section defines ciphersuites for the NIST P-384 elliptic curve [FIPS186-4].

P384_XMD:SHA-384_SSWU_RO_ is defined as follows:

- encoding type: `hash_to_curve` (Section 3)
- E: $y^2 = x^3 + A * x + B$, where
 - A = -3
 - B =
0xb3312fa7e23ee7e4988e056be3f82d19181d9c6efe8141120314088f5013875ac656398d8a2ed19d2a85c8edc
- p: $2^{384} - 2^{128} - 2^{96} + 2^{32} - 1$
- m: 1
- k: 192
- expand_message: `expand_message_xmd` (Section 5.3.1)
- H: SHA-384
- L: 72
- f: Simplified SWU method (Section 6.6.2)
- Z: -12
- h_eff: 1

P384_XMD:SHA-384_SSWU_NU_ is identical to P384_XMD:SHA-384_SSWU_RO_, except that the encoding type is `encode_to_curve` (Section 3).

An optimized example implementation of the Simplified SWU mapping to P-384 is given in Appendix F.2.

8.4. Suites for NIST P-521

This section defines ciphersuites for the NIST P-521 elliptic curve [FIPS186-4].

P521_XMD:SHA-512_SSWU_RO_ is defined as follows:

- encoding type: `hash_to_curve` (Section 3)
- E: $y^2 = x^3 + A * x + B$, where
 - A = -3

- B =
0x51953eb9618e1c9a1f929a21a0b68540eea2da725b99b315f3b8b489918ef109e156193951ec7e937b1652c0
- p: $2^{521} - 1$
- m: 1
- k: 256
- expand_message: expand_message_xmd ([Section 5.3.1](#))
- H: SHA-512
- L: 98
- f: Simplified SWU method ([Section 6.6.2](#))
- Z: -4
- h_eff: 1

P521_XMD:SHA-512_SSWU_NU_ is identical to P521_XMD:SHA-512_SSWU_RO_, except that the encoding type is encode_to_curve ([Section 3](#)).

An optimized example implementation of the Simplified SWU mapping to P-521 is given in [Appendix F.2](#).

8.5. Suites for curve25519 and edwards25519

This section defines ciphersuites for curve25519 and edwards25519 [[RFC7748](#)]. Note that these ciphersuites **MUST NOT** be used when hashing to ristretto255 [[ristretto255-decaf448](#)]. See [Appendix B](#) for information on how to hash to that group.

curve25519_XMD:SHA-512_ELL2_RO_ is defined as follows:

- encoding type: hash_to_curve ([Section 3](#))
- E: $K * t^2 = s^3 + J * s^2 + s$, where
 - J = 486662
 - K = 1
- p: $2^{255} - 19$
- m: 1
- k: 128
- expand_message: expand_message_xmd ([Section 5.3.1](#))
- H: SHA-512
- L: 48
- f: Elligator 2 method ([Section 6.7.1](#))
- Z: 2
- h_eff: 8

edwards25519_XMD:SHA-512_ELL2_RO_ is identical to curve25519_XMD:SHA-512_ELL2_RO_, except for the following parameters:

- E: $a * v^2 + w^2 = 1 + d * v^2 * w^2$, where
 - $a = -1$
 - $d = 0x52036cee2b6ffe738cc740797779e89800700a4d4141d8ab75eb4dca135978a3$
- f: Twisted Edwards Elligator 2 method ([Section 6.8.2](#))
- M: curve25519, defined in [[RFC7748](#)], [Section 4.1](#)
- rational_map: the birational maps defined in [[RFC7748](#)], [Section 4.1](#)

curve25519_XMD:SHA-512_ELL2_NU_ is identical to curve25519_XMD:SHA-512_ELL2_RO_, except that the encoding type is `encode_to_curve` ([Section 3](#)).

edwards25519_XMD:SHA-512_ELL2_NU_ is identical to edwards25519_XMD:SHA-512_ELL2_RO_, except that the encoding type is `encode_to_curve` ([Section 3](#)).

Optimized example implementations of the above mappings are given in [Appendix G.2.1](#) and [Appendix G.2.2](#).

8.6. Suites for curve448 and edwards448

This section defines ciphersuites for curve448 and edwards448 [[RFC7748](#)]. Note that these ciphersuites **MUST NOT** be used when hashing to decaf448 [[ristretto255-decaf448](#)]. See [Appendix C](#) for information on how to hash to that group.

curve448_XOF:SHAKE256_ELL2_RO_ is defined as follows:

- encoding type: `hash_to_curve` ([Section 3](#))
- E: $K * t^2 = s^3 + J * s^2 + s$, where
 - $J = 156326$
 - $K = 1$
- p: $2^{448} - 2^{224} - 1$
- m: 1
- k: 224
- expand_message: `expand_message_xof` ([Section 5.3.2](#))
- H: SHAKE256
- L: 84
- f: Elligator 2 method ([Section 6.7.1](#))
- Z: -1
- h_eff: 4

edwards448_XOF:SHAKE256_ELL2_RO_ is identical to curve448_XOF:SHAKE256_ELL2_RO_, except for the following parameters:

- E: $a * v^2 + w^2 = 1 + d * v^2 * w^2$, where
 - $a = 1$
 - $d = -39081$
- f: Twisted Edwards Elligator 2 method ([Section 6.8.2](#))
- M: curve448, defined in [[RFC7748](#)], [Section 4.2](#)
- rational_map: the 4-isogeny map defined in [[RFC7748](#)], [Section 4.2](#)

curve448_XOF:SHAKE256_ELL2_NU_ is identical to curve448_XOF:SHAKE256_ELL2_RO_, except that the encoding type is `encode_to_curve` ([Section 3](#)).

edwards448_XOF:SHAKE256_ELL2_NU_ is identical to edwards448_XOF:SHAKE256_ELL2_RO_, except that the encoding type is `encode_to_curve` ([Section 3](#)).

Optimized example implementations of the above mappings are given in [Appendix G.2.3](#) and [Appendix G.2.4](#).

8.7. Suites for secp256k1

This section defines ciphersuites for the secp256k1 elliptic curve [[SEC2](#)].

secp256k1_XMD:SHA-256_SSWU_RO_ is defined as follows:

- encoding type: `hash_to_curve` ([Section 3](#))
- E: $y^2 = x^3 + 7$
- p: $2^{256} - 2^{32} - 2^9 - 2^8 - 2^7 - 2^6 - 2^4 - 1$
- m: 1
- k: 128
- expand_message: `expand_message_xmd` ([Section 5.3.1](#))
- H: SHA-256
- L: 48
- f: Simplified SWU for $AB == 0$ ([Section 6.6.3](#))
- Z: -11
- E': $y'^2 = x'^3 + A' * x' + B'$, where
 - A': `0x3f8731abdd661adca08a5558f0f5d272e953d363cb6f0e5d405447c01a444533`
 - B': 1771
- iso_map: the 3-isogeny map from E' to E given in [Appendix E.1](#)
- h_eff: 1

secp256k1_XMD:SHA-256_SSWU_NU_ is identical to secp256k1_XMD:SHA-256_SSWU_RO_, except that the encoding type is `encode_to_curve` ([Section 3](#)).

An optimized example implementation of the Simplified SWU mapping to the curve E' isogenous to secp256k1 is given in [Appendix F.2](#).

8.8. Suites for BLS12-381

This section defines ciphersuites for groups G_1 and G_2 of the BLS12-381 elliptic curve [\[BLS12-381\]](#).

8.8.1. BLS12-381 G_1

`BLS12381G1_XMD:SHA-256_SSWU_RO_` is defined as follows:

- encoding type: `hash_to_curve` ([Section 3](#))
- E : $y^2 = x^3 + 4$
- p :
0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabfffeb153ffffb9fefffffffaaa
- m : 1
- k : 128
- `expand_message`: `expand_message_xmd` ([Section 5.3.1](#))
- H : SHA-256
- L : 64
- f : Simplified SWU for $AB == 0$ ([Section 6.6.3](#))
- Z : 11
- E' : $y'^2 = x'^3 + A' * x' + B'$, where
 - A' =
0x144698a3b8e9433d693a02c96d4982b0ea985383ee66a8d8e8981aefd881ac98936f8da0e0f97f5cf428082d5
 - B' =
0x12e2908d11688030018b12e8753eee3b2016c1f0f24f4070a0b9c14fcef35ef55a23215a316ceaa5d1cc48e98e
- `iso_map`: the 11-isogeny map from E' to E given in [Appendix E.2](#)
- `h_eff`: 0xd20100000010001

`BLS12381G1_XMD:SHA-256_SSWU_NU_` is identical to `BLS12381G1_XMD:SHA-256_SSWU_RO_`, except that the encoding type is `encode_to_curve` ([Section 3](#)).

Note that the `h_eff` values for these suites are chosen for compatibility with the fast cofactor clearing method described by Scott ([\[WB19\]](#), Section 5).

An optimized example implementation of the Simplified SWU mapping to the curve E' isogenous to BLS12-381 G_1 is given in [Appendix F.2](#).

8.8.2. BLS12-381 G_2

`BLS12381G2_XMD:SHA-256_SSWU_RO_` is defined as follows:

- encoding type: `hash_to_curve` ([Section 3](#))

- E: $y^2 = x^3 + 4 * (1 + I)$
- base field F is $GF(p^m)$, where
 - p:

0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabfffeb153ffffb9fefffffaa
 - m: 2
 - (1, I) is the basis for F, where $I^2 + 1 = 0$ in F
- k: 128
- expand_message: expand_message_xmd ([Section 5.3.1](#))
- H: SHA-256
- L: 64
- f: Simplified SWU for $AB = 0$ ([Section 6.6.3](#))
- Z: $-(2 + I)$
- E': $y'^2 = x'^3 + A' * x' + B'$, where
 - $A' = 240 * I$
 - $B' = 1012 * (1 + I)$
- iso_map: the isogeny map from E' to E given in [Appendix E.3](#)
- h_eff:

0xbc69f08f2ee75b3584c6a0ea91b352888e2a8e9145ad7689986ff031508ffe1329c2f178731db956d82bf015d121

BLS12381G2_XMD:SHA-256_SSWU_NU_ is identical to BLS12381G2_XMD:SHA-256_SSWU_RO_, except that the encoding type is encode_to_curve ([Section 3](#)).

Note that the h_eff values for these suites are chosen for compatibility with the fast cofactor clearing method described by Budroni and Pintore ([BP17], Section 4.1) and are summarized in [Appendix G.3](#).

An optimized example implementation of the Simplified SWU mapping to the curve E' isogenous to BLS12-381 G2 is given in [Appendix F.2](#).

8.9. Defining a New Hash-to-Curve Suite

For elliptic curves not listed elsewhere in [Section 8](#), a new hash-to-curve suite can be defined by the following:

1. E, F, p, and m are determined by the elliptic curve and its base field.
2. k is an upper bound on the target security level of the suite ([Section 10.8](#)). A reasonable choice of k is $\text{ceil}(\log_2(r) / 2)$, where r is the order of the subgroup G of the curve E ([Section 2.1](#)).
3. Choose encoding type, either hash_to_curve or encode_to_curve ([Section 3](#)).
4. Compute L as described in [Section 5](#).
5. Choose an expand_message variant from [Section 5.3](#) plus any underlying cryptographic primitives (e.g., a hash function H).

6. Choose a mapping following the guidelines in [Section 6.1](#), and select any required parameters for that mapping.
7. Choose `h_eff` to be either the cofactor of `E` or, if a fast cofactor clearing method is to be used, a value appropriate to that method as discussed in [Section 7](#).
8. Construct a Suite ID following the guidelines in [Section 8.10](#).

8.10. Suite ID Naming Conventions

Suite IDs **MUST** be constructed as follows:

```
CURVE_ID || "_" || HASH_ID || "_" || MAP_ID || "_" || ENC_VAR || "_"
```

The fields `CURVE_ID`, `HASH_ID`, `MAP_ID`, and `ENC_VAR` are ASCII-encoded strings of at most 64 characters each. Fields **MUST** contain only ASCII characters between 0x21 and 0x7E (inclusive), except that underscore (i.e., 0x5F) is not allowed.

As indicated above, each field (including the last) is followed by an underscore ("`_`", ASCII 0x5F). This helps to ensure that Suite IDs are prefix free. Suite IDs **MUST** include the final underscore and **MUST NOT** include any characters after the final underscore.

Suite ID fields **MUST** be chosen as follows:

- `CURVE_ID`: a human-readable representation of the target elliptic curve.
- `HASH_ID`: a human-readable representation of the `expand_message` function and any underlying hash primitives used in `hash_to_field` ([Section 5](#)). This field **MUST** be constructed as follows:

```
EXP_TAG || ":" || HASH_NAME
```

`EXP_TAG` indicates the `expand_message` variant:

- "XMD" for `expand_message_xmd` ([Section 5.3.1](#)).
- "XOF" for `expand_message_xof` ([Section 5.3.2](#)).

`HASH_NAME` is a human-readable name for the underlying hash primitive. As examples:

1. For `expand_message_xof` ([Section 5.3.2](#)) with SHAKE128, `HASH_ID` is "XOF:SHAKE128".
2. For `expand_message_xmd` ([Section 5.3.1](#)) with SHA3-256, `HASH_ID` is "XMD:SHA3-256".

Suites that use an alternative `hash_to_field` function that meets the requirements in [Section 5.1](#) **MUST** indicate this by appending a tag identifying that function to the `HASH_ID` field, separated by a colon ("`:`", ASCII 0x3A).

- `MAP_ID`: a human-readable representation of the `map_to_curve` function as defined in [Section 6](#). These are defined as follows:
 - "SVDW" for Shallue and van de Woestijne ([Section 6.6.1](#)).

- "SSWU" for Simplified SWU (Sections 6.6.2 and 6.6.3).
- "ELL2" for Elligator 2 (Sections 6.7.1 and 6.8.2).
- ENC_VAR: a string indicating the encoding type and other information. The first two characters of this string indicate whether the suite represents a hash_to_curve or an encode_to_curve operation (Section 3), as follows:
 - If ENC_VAR begins with "RO", the suite uses hash_to_curve.
 - If ENC_VAR begins with "NU", the suite uses encode_to_curve.
 - ENC_VAR **MUST NOT** begin with any other string.

ENC_VAR **MAY** also be used to encode other information used to identify variants, for example, a version number. The **RECOMMENDED** way to do so is to add one or more subfields separated by colons. For example, "RO:V02" is an appropriate ENC_VAR value for the second version of a uniform encoding suite, while "RO:V02:FOO01:BAR17" might be used to indicate a variant of that suite.

9. IANA Considerations

This document has no IANA actions.

10. Security Considerations

This section contains additional security considerations about the hash-to-curve mechanisms described in this document.

10.1. Properties of Encodings

Each encoding type (Section 3) accepts an arbitrary byte string and maps it to a point on the curve sampled from a distribution that depends on the encoding type. It is important to note that using a nonuniform encoding or directly evaluating one of the mappings of Section 6 produces an output that is easily distinguished from a uniformly random point. Applications that use a nonuniform encoding **SHOULD** carefully analyze the security implications of nonuniformity. When the required encoding is not clear, applications **SHOULD** use a uniform encoding.

Both encodings given in Section 3 can output the identity element of the group G . The probability that either encoding function outputs the identity element is roughly $1/r$ for a random input, which is negligible for cryptographically useful elliptic curves. Further, it is computationally infeasible to find an input to either encoding function whose corresponding output is the identity element. (Both of these properties hold when the encoding functions are instantiated with a hash_to_field function that follows all guidelines in Section 5.) Protocols that use these encoding functions **SHOULD NOT** add a special case to detect and "fix" the identity element.

When the hash_to_curve function (Section 3) is instantiated with a hash_to_field function that is indistinguishable from a random oracle (Section 5), the resulting function is indistinguishable from a random oracle ([MRH04] [BCIMRT10] [FFSTV13] [LBB19] [H20]). In many cases, such a function

can be safely used in cryptographic protocols whose security analysis assumes a random oracle that outputs uniformly random points on an elliptic curve. As Ristenpart et al. discuss in [RSS11], however, not all security proofs that rely on random oracles continue to hold when those oracles are replaced by indifferentiable functionalities. This limitation should be considered when analyzing the security of protocols relying on the `hash_to_curve` function.

10.2. Hashing Passwords

When hashing passwords using any function described in this document, an adversary who learns the output of the hash function (or potentially any intermediate value, e.g., the output of `hash_to_field`) may be able to carry out a dictionary attack. To mitigate such attacks, it is recommended to first execute a more costly key derivation function (e.g., PBKDF2 [RFC8018], `scrypt` [RFC7914], or Argon2 [RFC9106]) on the password, then hash the output of that function to the target elliptic curve. For collision resistance, the hash underlying the key derivation function should be chosen according to the guidelines listed in Section 5.3.1.

10.3. Constant-Time Requirements

Constant-time implementations of all functions in this document are **STRONGLY RECOMMENDED** for all uses, to avoid leaking information via side channels. It is especially important to use a constant-time implementation when inputs to an encoding are secret values; in such cases, constant-time implementations are **REQUIRED** for security against timing attacks (e.g., [VR20]). When constant-time implementations are required, all basic operations and utility functions must be implemented in constant time, as discussed in Section 4. In some applications (e.g., embedded systems), leakage through other side channels (e.g., power or electromagnetic side channels) may be pertinent. Defending against such leakage is outside the scope of this document, because the nature of the leakage and the appropriate defense depend on the application.

10.4. `encode_to_curve`: Output Distribution and Indifferentiability

The `encode_to_curve` function (Section 3) returns points sampled from a distribution that is statistically far from uniform. This distribution is bounded roughly as follows: first, it includes at least one eighth of the points in G , and second, the probability of points in the distribution varies by at most a factor of four. These bounds hold when `encode_to_curve` is instantiated with any of the `map_to_curve` functions in Section 6.

The bounds above are derived from several works in the literature. Specifically:

- Shallue and van de Woestijne [SW06] and Fouque and Tibouchi [FT12] derive bounds on the Shallue-van de Woestijne mapping (Section 6.6.1).
- Fouque and Tibouchi [FT10] and Tibouchi [T14] derive bounds for the Simplified SWU mapping (Sections 6.6.2 and 6.6.3).
- Bernstein et al. [BHKL13] derive bounds for the Elligator 2 mapping (Sections 6.7.1 and 6.8.2).

Indifferentiability of `encode_to_curve` follows from an argument similar to the one given by Brier et al. [BCIMRT10]; we briefly sketch this argument as follows. Consider an ideal random oracle $H_c()$ that samples from the distribution induced by the `map_to_curve` function called by `encode_to_curve`, and assume for simplicity that the target elliptic curve has cofactor 1 (a similar argument applies for non-unity cofactors). Indifferentiability holds just if it is possible to efficiently simulate the "inner" random oracle in `encode_to_curve`, namely, `hash_to_field`. The simulator works as follows: on a fresh query `msg`, the simulator queries $H_c(msg)$ and receives a point P in the image of `map_to_curve` (if `msg` is the same as a prior query, the simulator just returns the value it gave in response to that query). The simulator then computes the possible preimages of P under `map_to_curve`, i.e., elements u of F such that `map_to_curve(u) == P` (Tibouchi [T14] shows that this can be done efficiently for the Shallue-van de Woestijne and Simplified SWU maps, and Bernstein et al. show the same for Elligator 2). The simulator selects one such preimage at random and returns this value as the simulated output of the "inner" random oracle. By hypothesis, $H_c()$ samples from the distribution induced by `map_to_curve` on a uniformly random input element of F , so this value is uniformly random and induces the correct point P when passed through `map_to_curve`.

10.5. `hash_to_field` Security

The `hash_to_field` function, defined in Section 5, is indifferentiable from a random oracle [MRH04] when `expand_message` (Section 5.3) is modeled as a random oracle. Since indifferentiability proofs are composable, this also holds when `expand_message` is proved indifferentiable from a random oracle relative to an underlying primitive that is modeled as a random oracle. When following the guidelines in Section 5.3, both variants of `expand_message` defined in that section meet this requirement (see also Section 10.6).

We very briefly sketch the indifferentiability argument for `hash_to_field`. Notice that each integer mod p that `hash_to_field` returns (i.e., each element of the vector representation of F) is a member of an equivalence class of roughly 2^k integers of length $\log_2(p) + k$ bits, all of which are equal modulo p . For each integer mod p that `hash_to_field` returns, the simulator samples one member of this equivalence class at random and outputs the byte string returned by `I2OSP`. (Notice that this is essentially the inverse of the `hash_to_field` procedure.)

10.6. `expand_message_xmd` Security

The `expand_message_xmd` function, defined in Section 5.3.1, is indifferentiable from a random oracle [MRH04] when one of the following holds:

1. H is indifferentiable from a random oracle,
2. H is a sponge-based hash function whose inner function is modeled as a random transformation or random permutation [BDPV08], or
3. H is a Merkle-Damgaard hash function whose compression function is modeled as a random oracle [CDMP05].

For cases (1) and (2), the indifferentiability of `expand_message_xmd` follows directly from the indifferentiability of H .

For case (3), i.e., where H is a Merkle-Damgaard hash function, indistinguishability follows from [CDMP05], Theorem 5. In particular, `expand_message_xmd` computes b_0 by prefixing the message with one block of zeros plus auxiliary information (length, counter, and DST). Then, each of the output blocks b_i , $i \geq 1$ in `expand_message_xmd` is the result of invoking H on a unique, prefix-free encoding of b_0 . This is true, first because the length of the input to all such invocations is equal and fixed by the choice of H and DST, and second because each such input has a unique suffix (because of the inclusion of the counter byte $I2OSP(i, 1)$).

The essential difference between the construction discussed in [CDMP05] and `expand_message_xmd` is that the latter hashes a counter appended to `strxor(b_0, b_{(i-1)})` (`{#hashtofield-expand-xmd}`, step 10) rather than to b_0 . This approach increases the Hamming distance between inputs to different invocations of H , which reduces the likelihood that nonidealities in H affect the distribution of the b_i values.

We note that `expand_message_xmd` can be used to instantiate a general-purpose indistinguishable functionality with variable-length output based on any hash function meeting one of the above criteria. Applications that use `expand_message_xmd` outside of `hash_to_field` should ensure domain separation by picking a distinct value for DST.

10.7. Domain Separation for `expand_message` Variants

As discussed in Section 2.2.5, the purpose of domain separation is to ensure that security analyses of cryptographic protocols that query multiple independent random oracles remain valid even if all of these random oracles are instantiated based on one underlying function H .

The `expand_message` variants in this document (Section 5.3) ensure domain separation by appending a suffix-free-encoded domain separation tag `DST_prime` to all strings hashed by H , an underlying hash or extendable-output function. (Other `expand_message` variants that follow the guidelines in Section 5.3.4 are expected to behave similarly, but these should be analyzed on a case-by-case basis.) For security, applications that use the same function H outside of `expand_message` should enforce domain separation between those uses of H and `expand_message`, and they should separate all of these from uses of H in other applications.

This section suggests four methods for enforcing domain separation from `expand_message` variants, explains how each method achieves domain separation, and lists the situations in which each is appropriate. These methods share a high-level structure: the application designer fixes a tag `DST_ext` distinct from `DST_prime` and augments calls to H with `DST_ext`. Each method augments calls to H differently, and each may impose additional requirements on `DST_ext`.

These methods can be used to instantiate multiple domain-separated functions (e.g., H_1 and H_2) by selecting distinct `DST_ext` values for each (e.g., `DST_ext1`, `DST_ext2`).

1. (Suffix-only domain separation.) This method is useful when domain-separating invocations of H from `expand_message_xmd` or `expand_message_xof`. It is not appropriate for domain-separating `expand_message` from HMAC-H [RFC2104]; for that purpose, see method 4.

To instantiate a suffix-only domain-separated function H_{so} , compute

$$\text{Hso}(\text{msg}) = \text{H}(\text{msg} \parallel \text{DST_ext})$$

DST_ext should be suffix-free encoded (e.g., by appending one byte encoding the length of DST_ext) to make it infeasible to find distinct (msg, DST_ext) pairs that hash to the same value.

This method ensures domain separation because all distinct invocations of H have distinct suffixes, since DST_ext is distinct from DST_prime.

2. (Prefix-suffix domain separation.) This method can be used in the same cases as the suffix-only method.

To instantiate a prefix-suffix domain-separated function Hps, compute

$$\text{Hps}(\text{msg}) = \text{H}(\text{DST_ext} \parallel \text{msg} \parallel \text{I2OSP}(0, 1))$$

DST_ext should be prefix-free encoded (e.g., by adding a one-byte prefix that encodes the length of DST_ext) to make it infeasible to find distinct (msg, DST_ext) pairs that hash to the same value.

This method ensures domain separation because appending the byte I2OSP(0, 1) ensures that inputs to H inside Hps are distinct from those inside expand_message. Specifically, the final byte of DST_prime encodes the length of DST, which is required to be nonzero (Section 3.1, requirement 2), and DST_prime is always appended to invocations of H inside expand_message.

3. (Prefix-only domain separation.) This method is only useful for domain-separating invocations of H from expand_message_xmd. It does not give domain separation for expand_message_xof or HMAC-H.

To instantiate a prefix-only domain-separated function Hpo, compute

$$\text{Hpo}(\text{msg}) = \text{H}(\text{DST_ext} \parallel \text{msg})$$

In order for this method to give domain separation, DST_ext should be at least b bits long, where b is the number of bits output by the hash function H. In addition, at least one of the first b bits must be nonzero. Finally, DST_ext should be prefix-free encoded (e.g., by adding a one-byte prefix that encodes the length of DST_ext) to make it infeasible to find distinct (msg, DST_ext) pairs that hash to the same value.

This method ensures domain separation as follows. First, since DST_ext contains at least one nonzero bit among its first b bits, it is guaranteed to be distinct from the value Z_pad (Section 5.3.1, step 4), which ensures that all inputs to H are distinct from the input used to generate b_0 in expand_message_xmd. Second, since DST_ext is at least b bits long, it is almost certainly distinct from the values b_0 and strxor(b_0, b_(i - 1)), and therefore all inputs to H are distinct from the inputs used to generate b_i, i >= 1, with high probability.

4. (XMD-HMAC domain separation.) This method is useful for domain-separating invocations of H inside HMAC-H (i.e., HMAC [RFC2104] instantiated with hash function H) from `expand_message_xmd`. It also applies to HKDF-H (i.e., HKDF [RFC5869] instantiated with hash function H), as discussed below.

Specifically, this method applies when HMAC-H is used with a non-secret key to instantiate a random oracle based on a hash function H (note that `expand_message_xmd` can also be used for this purpose; see Section 10.6). When using HMAC-H with a high-entropy secret key, domain separation is not necessary; see discussion below.

To choose a non-secret HMAC key `DST_key` that ensures domain separation from `expand_message_xmd`, compute

```
DST_key_preimage = "DERIVE-HMAC-KEY-" || DST_ext || I2OSP(0, 1)
DST_key = H(DST_key_preimage)
```

Then, to instantiate the random oracle `Hro` using HMAC-H, compute

```
Hro(msg) = HMAC-H(DST_key, msg)
```

The trailing zero byte in `DST_key_preimage` ensures that this value is distinct from inputs to H inside `expand_message_xmd` (because all such inputs have suffix `DST_prime`, which cannot end with a zero byte as discussed above). This ensures domain separation because, with overwhelming probability, all inputs to H inside of HMAC-H using key `DST_key` have prefixes that are distinct from the values `Z_pad`, `b_0`, and `strxor(b_0, b_(i - 1))` inside of `expand_message_xmd`.

For uses of HMAC-H that instantiate a private random oracle by fixing a high-entropy secret key, domain separation from `expand_message_xmd` is not necessary. This is because, similarly to the case above, all inputs to H inside HMAC-H using this secret key almost certainly have distinct prefixes from all inputs to H inside `expand_message_xmd`.

Finally, this method can be used with HKDF-H [RFC5869] by fixing the salt input to HKDF-Extract to `DST_key`, computed as above. This ensures domain separation for HKDF-Extract by the same argument as for HMAC-H using `DST_key`. Moreover, assuming that the input keying material (IKM) supplied to HKDF-Extract has sufficiently high entropy (say, commensurate with the security parameter), the HKDF-Expand step is domain-separated by the same argument as for HMAC-H with a high-entropy secret key (since a pseudorandom key is exactly that).

10.8. Target Security Levels

Each ciphersuite specifies a target security level (in bits) for the underlying curve. This parameter ensures the corresponding `hash_to_field` instantiation is conservative and correct. We stress that this parameter is only an upper bound on the security level of the curve and is neither a guarantee nor endorsement of its suitability for a given application. Mathematical and cryptographic advancements may reduce the effective security level for any curve.

11. References

11.1. Normative References

- [Err4730] RFC Errata, "Erratum ID 4730", RFC 7748, July 2016, <<https://www.rfc-editor.org/errata/eid4730>>.
- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", BCP 14, RFC 2119, DOI 10.17487/RFC2119, March 1997, <<https://www.rfc-editor.org/info/rfc2119>>.
- [RFC7748] Langley, A., Hamburg, M., and S. Turner, "Elliptic Curves for Security", RFC 7748, DOI 10.17487/RFC7748, January 2016, <<https://www.rfc-editor.org/info/rfc7748>>.
- [RFC8017] Moriarty, K., Ed., Kaliski, B., Jonsson, J., and A. Rusch, "PKCS #1: RSA Cryptography Specifications Version 2.2", RFC 8017, DOI 10.17487/RFC8017, November 2016, <<https://www.rfc-editor.org/info/rfc8017>>.
- [RFC8174] Leiba, B., "Ambiguity of Uppercase vs Lowercase in RFC 2119 Key Words", BCP 14, RFC 8174, DOI 10.17487/RFC8174, May 2017, <<https://www.rfc-editor.org/info/rfc8174>>.

11.2. Informative References

- [AFQTZ14] Aranha, D. F., Fouque, P.-A., Qian, C., Tibouchi, M., and J. C. Zapalowicz, "Binary Elligator Squared", In Selected Areas in Cryptography - SAC 2014, pages 20-37, DOI 10.1007/978-3-319-13051-4_2, November 2014, <https://doi.org/10.1007/978-3-319-13051-4_2>.
- [AR13] Adj, G. and F. Rodríguez-Henríquez, "Square Root Computation over Even Extension Fields", In IEEE Transactions on Computers. vol 63 issue 11, pages 2829-2841, DOI 10.1109/TC.2013.145, November 2014, <<https://doi.org/10.1109/TC.2013.145>>.
- [BBJLP08] Bernstein, D. J., Birkner, P., Joye, M., Lange, T., and C. Peters, "Twisted Edwards Curves", In AFRICACRYPT 2008, pages 389-405, DOI 10.1007/978-3-540-68164-9_26, June 2008, <https://doi.org/10.1007/978-3-540-68164-9_26>.
- [BCIMRT10] Brier, E., Coron, J.-S., Icart, T., Madore, D., Randriam, H., and M. Tibouchi, "Efficient Indifferentiable Hashing into Ordinary Elliptic Curves", In Advances in Cryptology - CRYPTO 2010, pages 237-254, DOI 10.1007/978-3-642-14623-7_13, August 2010, <https://doi.org/10.1007/978-3-642-14623-7_13>.

-
- [BDPV08]** Bertoni, G., Daemen, J., Peeters, M., and G. Van Assche, "On the Indifferentiability of the Sponge Construction", In Advances in Cryptology - EUROCRYPT 2008, pages 181-197, DOI 10.1007/978-3-540-78967-3_11, April 2008, <https://doi.org/10.1007/978-3-540-78967-3_11>.
- [BF01]** Boneh, D. and M. Franklin, "Identity-Based Encryption from the Weil Pairing", In Advances in Cryptology - CRYPTO 2001, pages 213-229, DOI 10.1007/3-540-44647-8_13, August 2001, <https://doi.org/10.1007/3-540-44647-8_13>.
- [BHKL13]** Bernstein, D. J., Hamburg, M., Krasnova, A., and T. Lange, "Elligator: elliptic-curve points indistinguishable from uniform random strings", In Proceedings of the 2013 ACM SIGSAC Conference on Computer and Communications Security, pages 967-980, DOI 10.1145/2508859.2516734, November 2013, <<https://doi.org/10.1145/2508859.2516734>>.
- [BLAKE2X]** Aumasson, J.-P., Neves, S., Wilcox-O'Hearn, Z., and C. Winnerlein, "BLAKE2X", December 2016, <<https://blake2.net/blake2x.pdf>>.
- [BLMP19]** Bernstein, D. J., Lange, T., Martindale, C., and L. Panny, "Quantum Circuits for the CSIDH: Optimizing Quantum Evaluation of Isogenies", In Advances in Cryptology - EUROCRYPT 2019, pages 409-441, DOI 10.1007/978-3-030-17656-3, May 2019, <https://doi.org/10.1007/978-3-030-17656-3_15>.
- [BLS-SIG]** Boneh, D., Gorbunov, S., Wahby, R. S., Wee, H., Wood, C. A., and Z. Zhang, "BLS Signatures", Work in Progress, Internet-Draft, draft-irtf-cfrg-bls-signature-05, 16 June 2022, <<https://datatracker.ietf.org/doc/html/draft-irtf-cfrg-bls-signature-05>>.
- [BLS01]** Boneh, D., Lynn, B., and H. Shacham, "Short Signatures from the Weil Pairing", In Journal of Cryptology, vol 17, pages 297-319, DOI 10.1007/s00145-004-0314-9, July 2004, <<https://doi.org/10.1007/s00145-004-0314-9>>.
- [BLS03]** Barreto, P. S. L. M., Lynn, B., and M. Scott, "Constructing Elliptic Curves with Prescribed Embedding Degrees", In Security in Communication Networks, pages 257-267, DOI 10.1007/3-540-36413-7_19, September 2002, <https://doi.org/10.1007/3-540-36413-7_19>.
- [BLS12-381]** Bowe, S., "BLS12-381: New zk-SNARK Elliptic Curve Construction", March 2017, <<https://electriccoin.co/blog/new-snark-curve/>>.
- [BM92]** Bellare, S. M. and M. Merritt, "Encrypted key exchange: password-based protocols secure against dictionary attacks", In IEEE Symposium on Security and Privacy - Oakland 1992, pages 72-84, DOI 10.1109/RISP.1992.213269, May 1992, <<https://doi.org/10.1109/RISP.1992.213269>>.
- [BMP00]** Boyko, V., MacKenzie, P., and S. Patel, "Provably Secure Password-Authenticated Key Exchange Using Diffie-Hellman", In Advances in Cryptology - EUROCRYPT 2000, pages 156-171, DOI 10.1007/3-540-45539-6_12, May 2000, <https://doi.org/10.1007/3-540-45539-6_12>.

-
- [BN05] Barreto, P. S. L. M. and M. Naehrig, "Pairing-Friendly Elliptic Curves of Prime Order", In Selected Areas in Cryptography 2005, pages 319-331, DOI 10.1007/11693383_22, 2006, <https://doi.org/10.1007/11693383_22>.
- [BP17] Budroni, A. and F. Pintore, "Efficient hash maps to \mathbb{G}_2 on BLS curves", Cryptology ePrint Archive, Paper 2017/419, May 2017, <<https://eprint.iacr.org/2017/419>>.
- [BR93] Bellare, M. and P. Rogaway, "Random oracles are practical: a paradigm for designing efficient protocols", In Proceedings of the 1993 ACM Conference on Computer and Communications Security, pages 62-73, DOI 10.1145/168588.168596, December 1993, <<https://doi.org/10.1145/168588.168596>>.
- [C93] Cohen, H., "A Course in Computational Algebraic Number Theory", Springer-Verlag, ISBN 9783642081422, DOI 10.1007/978-3-662-02945-9, 1993, <<https://doi.org/10.1007/978-3-662-02945-9>>.
- [CDMP05] Coron, J.-S., Dodis, Y., Malinaud, C., and P. Puniya, "Merkle-Damgård Revisited: How to Construct a Hash Function", In Advances in Cryptology -- CRYPTO 2005, pages 430-448, DOI 10.1007/11535218_26, August 2005, <https://doi.org/10.1007/11535218_26>.
- [CFADLNV05] Cohen, H., Frey, G., Avanzi, R., Doche, C., Lange, T., Nguyen, K., and F. Vercauteren, "Handbook of Elliptic and Hyperelliptic Curve Cryptography", Chapman and Hall / CRC, ISBN 9781584885184, 2005, <<https://www.crcpress.com/9781584885184>>.
- [CK11] Couveignes, J.-M. and J.-G. Kammerer, "The geometry of flex tangents to a cubic curve and its parameterizations", In Journal of Symbolic Computation, vol 47 issue 3, pages 266-281, DOI 10.1016/j.jsc.2011.11.003, March 2012, <<https://doi.org/10.1016/j.jsc.2011.11.003>>.
- [F11] Farashahi, R. R., "Hashing into Hessian Curves", In AFRICACRYPT 2011, pages 278-289, DOI 10.1007/978-3-642-21969-6_17, July 2011, <https://doi.org/10.1007/978-3-642-21969-6_17>.
- [FFSTV13] Farashahi, R. R., Fouque, P.-A., Shparlinski, I. E., Tibouchi, M., and J. F. Voloch, "Indifferentiable deterministic hashing to elliptic and hyperelliptic curves", In Mathematics of Computation. vol 82, pages 491-512, DOI 10.1090/S0025-5718-2012-02606-8, 2013, <<https://doi.org/10.1090/S0025-5718-2012-02606-8>>.
- [FIPS180-4] National Institute of Standards and Technology (NIST), "Secure Hash Standard (SHS)", FIPS 180-4, DOI 10.6028/NIST.FIPS.180-4, August 2015, <<https://nvlpubs.nist.gov/nistpubs/FIPS/NIST.FIPS.180-4.pdf>>.
- [FIPS186-4] National Institute of Standards and Technology (NIST), "Digital Signature Standard (DSS)", FIPS 186-4, DOI 10.6028/NIST.FIPS.186-4, July 2013, <<https://nvlpubs.nist.gov/nistpubs/FIPS/NIST.FIPS.186-4.pdf>>.
-

-
- [FIPS202]** National Institute of Standards and Technology (NIST), "SHA-3 Standard: Permutation-Based Hash and Extendable-Output Functions", FIPS 202, DOI 10.6028/NIST.FIPS.202, August 2015, <<https://nvlpubs.nist.gov/nistpubs/FIPS/NIST.FIPS.202.pdf>>.
- [FJT13]** Fouque, P.-A., Joux, A., and M. Tibouchi, "Injective Encodings to Elliptic Curves", In ACISP 2013, pages 203-218, DOI 10.1007/978-3-642-39059-3_14, 2013, <https://doi.org/10.1007/978-3-642-39059-3_14>.
- [FKR11]** Fuentes-Castañeda, L., Knapp, E., and F. Rodriguez-Henriquez, "Faster Hashing to G²", In Selected Areas in Cryptography, pages 412-430, DOI 10.1007/978-3-642-28496-0_25, August 2011, <https://doi.org/10.1007/978-3-642-28496-0_25>.
- [FSV09]** Farashahi, R. R., Shparlinski, I. E., and J. F. Voloch, "On hashing into elliptic curves", In Journal of Mathematical Cryptology, vol 3 no 4, pages 353-360, DOI 10.1515/JMC.2009.022, March 2009, <<https://doi.org/10.1515/JMC.2009.022>>.
- [FT10]** Fouque, P.-A. and M. Tibouchi, "Estimating the Size of the Image of Deterministic Hash Functions to Elliptic Curves", In Progress in Cryptology - LATINCRYPT 2010, pages 81-91, DOI 10.1007/978-3-642-14712-8_5, August 2010, <https://doi.org/10.1007/978-3-642-14712-8_5>.
- [FT12]** Fouque, P.-A. and M. Tibouchi, "Indifferentiable Hashing to Barreto--Naehrig Curves", In Progress in Cryptology - LATINCRYPT 2012, pages 1-17, DOI 10.1007/978-3-642-33481-8_1, 2012, <https://doi.org/10.1007/978-3-642-33481-8_1>.
- [H20]** Hamburg, M., "Indifferentiable hashing from Elligator 2", Cryptology ePrint Archive, Paper 2020/1513, 2020, <<https://eprint.iacr.org/2020/1513>>.
- [hash2curve-repo]** "Hashing to Elliptic Curves", commit 664b135, June 2022, <<https://github.com/cfrg/draft-irtf-cfrg-hash-to-curve>>.
- [Icart09]** Icart, T., "How to Hash into Elliptic Curves", In Advances in Cryptology - CRYPTO 2009, pages 303-316, DOI 10.1007/978-3-642-03356-8_18, August 2009, <https://doi.org/10.1007/978-3-642-03356-8_18>.
- [J96]** Jablon, D. P., "Strong password-only authenticated key exchange", In SIGCOMM Computer Communication Review, vol 26 issue 5, pages 5-26, DOI 10.1145/242896.242897, October 1996, <<https://doi.org/10.1145/242896.242897>>.
- [jubjub-fq]** "zkcrypto/jubjub - fq.rs", 2019, <<https://github.com/zkcrypto/jubjub/pull/18>>.
- [KLR10]** Kammerer, J.-G., Lercier, R., and G. Renault, "Encoding Points on Hyperelliptic Curves over Finite Fields in Deterministic Polynomial Time", In Pairing-Based Cryptography - Pairing 2010, pages 278-297, DOI 10.1007/978-3-642-17455-1_18, 2010, <https://doi.org/10.1007/978-3-642-17455-1_18>.

-
- [L13]** Langley, A., "Implementing Elligator for Curve25519", December 2013, <<https://www.imperialviolet.org/2013/12/25/elligator.html>>.
- [LBB19]** Lipp, B., Blanchet, B., and K. Bhargavan, "A Mechanised Cryptographic Proof of the WireGuard Virtual Private Network Protocol", In INRIA Research Report 9269, April 2019, <<https://hal.inria.fr/hal-02100345/>>.
- [MOV96]** Menezes, A. J., van Oorschot, P. C., and S. A. Vanstone, "Handbook of Applied Cryptography", CRC Press, ISBN 9780849385230, October 1996, <<http://cacr.uwaterloo.ca/hac/>>.
- [MRH04]** Maurer, U., Renner, R., and C. Holenstein, "Indifferentiability, Impossibility Results on Reductions, and Applications to the Random Oracle Methodology", In TCC 2004: Theory of Cryptography, pages 21-39, DOI 10.1007/978-3-540-24638-1_2, February 2004, <https://doi.org/10.1007/978-3-540-24638-1_2>.
- [MRV99]** Micali, S., Rabin, M., and S. Vadhan, "Verifiable random functions", 40th Annual Symposium on Foundations of Computer Science (Cat. No.99CB37039), pages 120-130, DOI 10.1109/SFFCS.1999.814584, October 1999, <<https://doi.org/10.1109/SFFCS.1999.814584>>.
- [MT98]** Matsumoto, M. and T. Nishimura, "Mersenne twister: A 623-dimensionally equidistributed uniform pseudo-random number generator", In ACM Transactions on Modeling and Computer Simulation (TOMACS), vol 8 issue 1, pages 3-30, DOI 10.1145/272991.272995, January 1998, <<https://doi.org/10.1145/272991.272995>>.
- [NR97]** Naor, M. and O. Reingold, "Number-theoretic constructions of efficient pseudo-random functions", In Proceedings 38th Annual Symposium on Foundations of Computer Science, pages 458-467, DOI 10.1109/SFCS.1997.646134, October 1997, <<https://doi.org/10.1109/SFCS.1997.646134>>.
- [OPRFs]** Davidson, A., Faz-Hernandez, A., Sullivan, N., and C. A. Wood, "Oblivious Pseudorandom Functions (OPRFs) using Prime-Order Groups", Work in Progress, Internet-Draft, draft-irtf-cfrg-voprf-21, 21 February 2023, <<https://datatracker.ietf.org/doc/html/draft-irtf-cfrg-voprf-21>>.
- [p1363.2]** IEEE, "IEEE Standard Specification for Password-Based Public-Key Cryptography Techniques", IEEE 1363.2-2008, September 2008, <https://standards.ieee.org/standard/1363_2-2008.html>.
- [p1363a]** IEEE, "IEEE Standard Specifications for Public-Key Cryptography - Amendment 1: Additional Techniques", IEEE 1363a-2004, March 2004, <<https://standards.ieee.org/standard/1363a-2004.html>>.
- [P20]** Pornin, T., "Efficient Elliptic Curve Operations On Microcontrollers With Finite Field Extensions", Cryptology ePrint Archive, Paper 2020/009, 2020, <<https://eprint.iacr.org/2020/009>>.

-
- [RCB16] Renes, J., Costello, C., and L. Batina, "Complete Addition Formulas for Prime Order Elliptic Curves", In *Advances in Cryptology - EUROCRYPT 2016*, pages 403-428, DOI 10.1007/978-3-662-49890-3_16, April 2016, <https://doi.org/10.1007/978-3-662-49890-3_16>.
- [RFC2104] Krawczyk, H., Bellare, M., and R. Canetti, "HMAC: Keyed-Hashing for Message Authentication", RFC 2104, DOI 10.17487/RFC2104, February 1997, <<https://www.rfc-editor.org/info/rfc2104>>.
- [RFC5869] Krawczyk, H. and P. Eronen, "HMAC-based Extract-and-Expand Key Derivation Function (HKDF)", RFC 5869, DOI 10.17487/RFC5869, May 2010, <<https://www.rfc-editor.org/info/rfc5869>>.
- [RFC7693] Saarinen, M., Ed. and J. Aumasson, "The BLAKE2 Cryptographic Hash and Message Authentication Code (MAC)", RFC 7693, DOI 10.17487/RFC7693, November 2015, <<https://www.rfc-editor.org/info/rfc7693>>.
- [RFC7914] Percival, C. and S. Josefsson, "The scrypt Password-Based Key Derivation Function", RFC 7914, DOI 10.17487/RFC7914, August 2016, <<https://www.rfc-editor.org/info/rfc7914>>.
- [RFC8018] Moriarty, K., Ed., Kaliski, B., and A. Rusch, "PKCS #5: Password-Based Cryptography Specification Version 2.1", RFC 8018, DOI 10.17487/RFC8018, January 2017, <<https://www.rfc-editor.org/info/rfc8018>>.
- [RFC9106] Biryukov, A., Dinu, D., Khovratovich, D., and S. Josefsson, "Argon2 Memory-Hard Function for Password Hashing and Proof-of-Work Applications", RFC 9106, DOI 10.17487/RFC9106, September 2021, <<https://www.rfc-editor.org/info/rfc9106>>.
- [ristretto255-decaf448] de Valence, H., Grigg, J., Hamburg, M., Lovecruft, I., Tankersley, G., and F. Valsorda, "The ristretto255 and decaf448 Groups", Work in Progress, Internet-Draft, draft-irtf-cfrg-ristretto255-decaf448-07, 3 April 2023, <<https://datatracker.ietf.org/doc/html/draft-irtf-cfrg-ristretto255-decaf448-07>>.
- [RSS11] Ristenpart, T., Shacham, H., and T. Shrimpton, "Careful with Composition: Limitations of the Indifferentiability Framework", In *Advances in Cryptology - EUROCRYPT 2011*, pages 487-506, DOI 10.1007/978-3-642-20465-4_27, May 2011, <https://doi.org/10.1007/978-3-642-20465-4_27>.
- [S05] Skałba, M., "Points on elliptic curves over finite fields", In *Acta Arithmetica*, vol 117 no 3, pages 293-301, DOI 10.4064/aa117-3-7, 2005, <<https://doi.org/10.4064/aa117-3-7>>.
- [S85] Schoof, R., "Elliptic curves over finite fields and the computation of square roots mod p ", In *Mathematics of Computation*, vol 44 issue 170, pages 483-494, DOI 10.1090/S0025-5718-1985-0777280-6, April 1985, <<https://doi.org/10.1090/S0025-5718-1985-0777280-6>>.
- [SAGE] The Sage Developers, "SageMath, the Sage Mathematics Software System", <<https://www.sagemath.org>>.
-

- [SBCDK09]** Scott, M., Benger, N., Charlemagne, M., Dominguez Perez, L. J., and E. J. Kachisa, "Fast Hashing to G2 on Pairing-Friendly Curves", In Pairing-Based Cryptography - Pairing 2009, pages 102-113, DOI 10.1007/978-3-642-03298-1_8, August 2009, <https://doi.org/10.1007/978-3-642-03298-1_8>.
- [SEC1]** Standards for Efficient Cryptography Group (SECG), "SEC 1: Elliptic Curve Cryptography", May 2009, <<http://www.secg.org/sec1-v2.pdf>>.
- [SEC2]** Standards for Efficient Cryptography Group (SECG), "SEC 2: Recommended Elliptic Curve Domain Parameters", January 2010, <<http://www.secg.org/sec2-v2.pdf>>.
- [SS04]** Schinzel, A. and M. Skalba, "On equations $y^2 = x^n + k$ in a finite field", In Bulletin Polish Academy of Sciences. Mathematics, vol 52 no 3, pages 223-226, DOI 10.4064/ba52-3-1, 2004, <<https://doi.org/10.4064/ba52-3-1>>.
- [SW06]** Shallue, A. and C. E. van de Woestijne, "Construction of Rational Points on Elliptic Curves over Finite Fields", In Algorithmic Number Theory - ANTS 2006, pages 510-524, DOI 10.1007/11792086_36, July 2006, <https://doi.org/10.1007/11792086_36>.
- [T14]** Tibouchi, M., "Elligator Squared: Uniform Points on Elliptic Curves of Prime Order as Uniform Random Strings", In Financial Cryptography and Data Security - FC 2014, pages 139-156, DOI 10.1007/978-3-662-45472-5_10, November 2014, <https://doi.org/10.1007/978-3-662-45472-5_10>.
- [TK17]** Tibouchi, M. and T. Kim, "Improved elliptic curve hashing and point representation", In Designs, Codes, and Cryptography, vol 82, pages 161-177, DOI 10.1007/s10623-016-0288-2, January 2017, <<https://doi.org/10.1007/s10623-016-0288-2>>.
- [U07]** Ulas, M., "Rational Points on Certain Hyperelliptic Curves over Finite Fields", In Bulletin Polish Academy of Science. Mathematics, vol 55 no 2, pages 97-104, DOI 10.4064/ba55-2-1, July 2007, <<https://doi.org/10.4064/ba55-2-1>>.
- [VR20]** Vanhoef, M. and E. Ronen, "Dragonblood: Analyzing the Dragonfly Handshake of WPA3 and EAP-pwd", In IEEE Symposium on Security & Privacy (SP), May 2020, <<https://eprint.iacr.org/2019/383>>.
- [VRF]** Goldberg, S., Reyzin, L., Papadopoulos, D., and J. Včelák, "Verifiable Random Functions (VRFs)", Work in Progress, Internet-Draft, draft-irtf-cfrg-vrf-15, 9 August 2022, <<https://datatracker.ietf.org/doc/html/draft-irtf-cfrg-vrf-15>>.
- [W08]** Washington, L. C., "Elliptic Curves: Number Theory and Cryptography, Second Edition", Chapman and Hall / CRC, ISBN 9781420071467, April 2008, <<https://www.crcpress.com/9781420071467>>.
- [W19]** Wahby, R. S., "An explicit, generic parameterization for the Shallue--van de Woestijne map", commit e2a625f, March 2020, <https://github.com/cfrg/draft-irtf-cfrg-hash-to-curve/blob/draft-irtf-cfrg-hash-to-curve-14/doc/svdw_params.pdf>.

- [WB19] Wahby, R. S. and D. Boneh, "Fast and simple constant-time hashing to the BLS12-381 elliptic curve", In IACR Transactions on Cryptographic Hardware and Embedded Systems, vol 2019 issue 4, Cryptology ePrint Archive, Paper 2019/403, DOI 10.13154/tches.v2019.i4.154-179, August 2019, <<https://eprint.iacr.org/2019/403>>.

Appendix A. Related Work

The problem of mapping arbitrary bit strings to elliptic curve points has been the subject of both practical and theoretical research. This section briefly describes the background and research results that underlie the recommendations in this document. This section is provided for informational purposes only.

A naive but generally insecure method of mapping a string msg to a point on an elliptic curve E having n points is to first fix a point P that generates the elliptic curve group, and a hash function H_n from bit strings to integers less than n ; then compute $H_n(msg) * P$, where the $*$ operator represents scalar multiplication. The reason this approach is insecure is that the resulting point has a known discrete log relationship to P . Thus, except in cases where this method is specified by the protocol, it must not be used; doing so risks catastrophic security failures.

Boneh et al. [BLS01] describe an encoding method they call `MapToGroup`, which works roughly as follows: first, use the input string to initialize a pseudorandom number generator, then use the generator to produce a value x in F . If x is the x -coordinate of a point on the elliptic curve, output that point. Otherwise, generate a new value x in F and try again. Since a random value x in F has probability about $1/2$ of corresponding to a point on the curve, the expected number of tries is just two. However, the running time of this method, which is generally referred to as a probabilistic try-and-increment algorithm, depends on the input string. As such, it is not safe to use in protocols sensitive to timing side channels, as was exemplified by the Dragonblood attack [VR20].

Schinzel and Skalba [SS04] introduce a method of constructing elliptic curve points deterministically, for a restricted class of curves and a very small number of points. Skalba [S05] generalizes this construction to more curves and more points on those curves. Shallue and van de Woestijne [SW06] further generalize and simplify Skalba's construction, yielding concretely efficient maps to a constant fraction of the points on almost any curve. Fouque and Tibouchi [FT12] give a parameterization of this mapping for Barreto-Naehrig pairing-friendly curves [BN05].

Ulas [U07] describes a simpler version of the Shallue-van de Woestijne map, and Brier et al. [BCIMRT10] give a further simplification, which the authors call the "Simplified SWU" map. That simplified map applies only to fields of characteristic $p = 3 \pmod{4}$; Wahby and Boneh [WB19] generalize to fields of any characteristic and give further optimizations.

Boneh and Franklin give a deterministic algorithm mapping to certain supersingular curves over fields of characteristic $p = 2 \pmod{3}$ [BF01]. Icart gives another deterministic algorithm that maps to any curve over a field of characteristic $p = 2 \pmod{3}$ [Icart09]. Several extensions and generalizations follow this work, including [FSV09], [FT10], [KLR10], [F11], and [CK11].

Following the work of Farashahi [F11], Fouque et al. [FJT13] describe a mapping to curves over fields of characteristic $p = 3 \pmod{4}$ having a number of points divisible by 4. Bernstein et al. [BHKL13] optimize this mapping and describe a related mapping that they call "Elligator 2," which applies to any curve over a field of odd characteristic having a point of order 2. This includes Curve25519 and Curve448, both of which are CFRG-recommended curves [RFC7748]. Bernstein et al. [BLMP19] extend the Elligator 2 map to a class of supersingular curves over fields of characteristic $p = 3 \pmod{4}$.

An important caveat regarding all of the above deterministic mapping functions is that none of them map to the entire curve, but rather to some fraction of the points. This means that they cannot be used directly to construct a random oracle that outputs points on the curve.

Brier et al. [BCIMRT10] give two solutions to this problem. The first, which Brier et al. prove applies to Icart's method, computes $f(H_0(\text{msg})) + f(H_1(\text{msg}))$ for two distinct hash functions H_0 and H_1 from bit strings to F and a mapping f from F to the elliptic curve E . The second, which applies to essentially all deterministic mappings but is more costly, computes $f(H_0(\text{msg})) + H_2(\text{msg}) * P$, where P is a generator of the elliptic curve group, H_2 is a hash from bit strings to integers modulo r , and r is the order of the elliptic curve group.

Farashahi et al. [FFSTV13] improve the analysis of the first method, showing that it applies to essentially all deterministic mappings. Tibouchi and Kim [TK17] further refine the analysis and describe additional optimizations.

Complementary to the problem of mapping from bit strings to elliptic curve points, Bernstein et al. [BHKL13] study the problem of mapping from elliptic curve points to uniformly random bit strings, giving solutions for a class of curves that includes Montgomery and twisted Edwards curves. Tibouchi [T14] and Aranha et al. [AFQTZ14] generalize these results. This document does not deal with this complementary problem.

Appendix B. Hashing to ristretto255

ristretto255 [ristretto255-decaf448] provides a prime-order group based on curve25519 [RFC7748]. This section describes `hash_to_ristretto255`, which implements a random-oracle encoding to this group that has a uniform output distribution (Section 2.2.3) and the same security properties and interface as the `hash_to_curve` function (Section 3).

The ristretto255 API defines a one-way map ([ristretto255-decaf448], Section 4.3.4); this section refers to that map as `ristretto255_map`.

The `hash_to_ristretto255` function **MUST** be instantiated with an `expand_message` function that conforms to the requirements given in Section 5.3. In addition, it **MUST** use a domain separation tag constructed as described in Section 3.1, and all domain separation recommendations given in Section 10.7 apply when implementing protocols that use `hash_to_ristretto255`.

```
hash_to_ristretto255(msg)
```

Parameters:

- DST, a domain separation tag (see discussion above).
- `expand_message`, a function that expands a byte string and domain separation tag into a uniformly random byte string (see discussion above).
- `ristretto255_map`, the one-way map from the `ristretto255` API.

Input: `msg`, an arbitrary-length byte string.

Output: `P`, an element of the `ristretto255` group.

Steps:

1. `uniform_bytes = expand_message(msg, DST, 64)`
2. `P = ristretto255_map(uniform_bytes)`
3. return `P`

Since `hash_to_ristretto255` is not a hash-to-curve suite, it does not have a Suite ID. If a similar identifier is needed, it **MUST** be constructed following the guidelines in [Section 8.10](#), with the following parameters:

- CURVE_ID: "ristretto255"
- HASH_ID: as described in [Section 8.10](#)
- MAP_ID: "R255MAP"
- ENC_VAR: "RO"

For example, if `expand_message` is `expand_message_xmd` using SHA-512, the **REQUIRED** identifier is:

```
ristretto255_XMD:SHA-512_R255MAP_RO_
```

Appendix C. Hashing to decaf448

Similar to `ristretto255`, `decaf448` [[ristretto255-decaf448](#)] provides a prime-order group based on `curve448` [[RFC7748](#)]. This section describes `hash_to_decaf448`, which implements a random-oracle encoding to this group that has a uniform output distribution ([Section 2.2.3](#)) and the same security properties and interface as the `hash_to_curve` function ([Section 3](#)).

The `decaf448` API defines a one-way map ([[ristretto255-decaf448](#)], [Section 5.3.4](#)); this section refers to that map as `decaf448_map`.

The `hash_to_decaf448` function **MUST** be instantiated with an `expand_message` function that conforms to the requirements given in [Section 5.3](#). In addition, it **MUST** use a domain separation tag constructed as described in [Section 3.1](#), and all domain separation recommendations given in [Section 10.7](#) apply when implementing protocols that use `hash_to_decaf448`.

```
hash_to_decaf448(msg)
```

Parameters:

- DST, a domain separation tag (see discussion above).
- `expand_message`, a function that expands a byte string and domain separation tag into a uniformly random byte string (see discussion above).
- `decaf448_map`, the one-way map from the decaf448 API.

Input: `msg`, an arbitrary-length byte string.

Output: `P`, an element of the decaf448 group.

Steps:

1. `uniform_bytes = expand_message(msg, DST, 112)`
2. `P = decaf448_map(uniform_bytes)`
3. return `P`

Since `hash_to_decaf448` is not a hash-to-curve suite, it does not have a Suite ID. If a similar identifier is needed, it **MUST** be constructed following the guidelines in [Section 8.10](#), with the following parameters:

- CURVE_ID: "decaf448"
- HASH_ID: as described in [Section 8.10](#)
- MAP_ID: "D448MAP"
- ENC_VAR: "RO"

For example, if `expand_message` is `expand_message_xof` using SHAKE256, the **REQUIRED** identifier is:

```
decaf448_XOF : SHAKE256_D448MAP_RO_
```

Appendix D. Rational Maps

This section gives rational maps that can be used when hashing to twisted Edwards or Montgomery curves.

Given a twisted Edwards curve, [Appendix D.1](#) shows how to derive a corresponding Montgomery curve and how to map from that curve to the twisted Edwards curve. This mapping may be used when hashing to twisted Edwards curves as described in [Section 6.8](#).

Given a Montgomery curve, [Appendix D.2](#) shows how to derive a corresponding Weierstrass curve and how to map from that curve to the Montgomery curve. This mapping can be used to hash to Montgomery or twisted Edwards curves via the Shallue-van de Woestijne method ([Section 6.6.1](#)) or Simplified SWU method ([Section 6.6.2](#)), as follows:

- For Montgomery curves, first map to the Weierstrass curve, then convert to Montgomery coordinates via the mapping.

- For twisted Edwards curves, compose the mapping from Weierstrass to Montgomery with the mapping from Montgomery to twisted Edwards ([Appendix D.1](#)) to obtain a Weierstrass curve and a mapping to the target twisted Edwards curve. Map to this Weierstrass curve, then convert to Edwards coordinates via the mapping.

D.1. Generic Mapping from Montgomery to Twisted Edwards

This section gives a generic birational map between twisted Edwards and Montgomery curves.

The map in this section is a simplified version of the map given in [\[BBJLP08\]](#), Theorem 3.2. Specifically, this section's map handles exceptional cases in a simplified way that is geared towards hashing to a twisted Edwards curve's prime-order subgroup.

The twisted Edwards curve

$$a * v^2 + w^2 = 1 + d * v^2 * w^2$$

is birationally equivalent to the Montgomery curve

$$K * t^2 = s^3 + J * s^2 + s$$

which has the form required by the Elligator 2 mapping of [Section 6.7.1](#). The coefficients of the Montgomery curve are

- $J = 2 * (a + d) / (a - d)$
- $K = 4 / (a - d)$

The rational map from the point (s, t) on the above Montgomery curve to the point (v, w) on the twisted Edwards curve is given by

- $v = s / t$
- $w = (s - 1) / (s + 1)$

This mapping is undefined when $t == 0$ or $s == -1$, i.e., when the denominator of either of the above rational functions is zero. Implementations **MUST** detect exceptional cases and return the value $(v, w) = (0, 1)$, which is the identity point on all twisted Edwards curves.

The following straight-line implementation of the above rational map handles the exceptional cases.

```
monty_to_edw_generic(s, t)
```

Input: (s, t) , a point on the curve $K * t^2 = s^3 + J * s^2 + s$.
Output: (v, w) , a point on an equivalent twisted Edwards curve.

```
1. tv1 = s + 1
2. tv2 = tv1 * t      # (s + 1) * t
3. tv2 = inv0(tv2)   # 1 / ((s + 1) * t)
4. v = tv2 * tv1     # 1 / t
5. v = v * s         # s / t
6. w = tv2 * t       # 1 / (s + 1)
7. tv1 = s - 1
8. w = w * tv1       # (s - 1) / (s + 1)
9. e = tv2 == 0
10. w = CMOV(w, 1, e) # handle exceptional case
11. return (v, w)
```

For completeness, we also give the inverse relations. (Note that this map is not required when hashing to twisted Edwards curves.) The coefficients of the twisted Edwards curve corresponding to the above Montgomery curve are

- $a = (J + 2) / K$
- $d = (J - 2) / K$

The rational map from the point (v, w) on the twisted Edwards curve to the point (s, t) on the Montgomery curve is given by

- $s = (1 + w) / (1 - w)$
- $t = (1 + w) / (v * (1 - w))$

The mapping is undefined when $v == 0$ or $w == 1$. When the goal is to map into the prime-order subgroup of the Montgomery curve, it suffices to return the identity point on the Montgomery curve in the exceptional cases.

D.2. Mapping from Weierstrass to Montgomery

The rational map from the point (s, t) on the Montgomery curve

$$K * t^2 = s^3 + J * s^2 + s$$

to the point (x, y) on the equivalent Weierstrass curve

$$y^2 = x^3 + A * x + B$$

is given by

- $A = (3 - J^2) / (3 * K^2)$

- $k_{(2,1)} = 0xedadc6f64383dc1df7c4b2d51b54225406d36b641f5e41bbc52a56612a8c6d14$

The constants used to compute y_{num} are as follows:

- $k_{(3,0)} = 0x4bda12f684bda12f684bda12f684bda12f684bda12f684bda12f684b8e38e23c$
- $k_{(3,1)} = 0xc75e0c32d5cb7c0fa9d0a54b12a0a6d5647ab046d686da6fdffc90fc201d71a3$
- $k_{(3,2)} = 0x29a6194691f91a73715209ef6512e576722830a201be2018a765e85a9ecee931$
- $k_{(3,3)} = 0x2f684bda12f684bda12f684bda12f684bda12f684bda12f684bda12f38e38d84$

The constants used to compute y_{den} are as follows:

- $k_{(4,0)} = 0xff93b$
- $k_{(4,1)} = 0x7a06534bb8bdb49fd5e9e6632722c2989467c1bfc8e8d978dfb425d2685c2573$
- $k_{(4,2)} = 0x6484aa716545ca2cf3a70c3fa8fe337e0a3d21162f0d6299a7bf8192bfd2a76f$

E.2. 11-Isogeny Map for BLS12-381 G1

The 11-isogeny map from (x', y') on E' to (x, y) on E is given by the following rational functions:

- $x = x_{num} / x_{den}$, where
 - $x_{num} = k_{(1,11)} * x'^{11} + k_{(1,10)} * x'^{10} + k_{(1,9)} * x'^9 + \dots + k_{(1,0)}$
 - $x_{den} = x'^{10} + k_{(2,9)} * x'^9 + k_{(2,8)} * x'^8 + \dots + k_{(2,0)}$
- $y = y' * y_{num} / y_{den}$, where
 - $y_{num} = k_{(3,15)} * x'^{15} + k_{(3,14)} * x'^{14} + k_{(3,13)} * x'^{13} + \dots + k_{(3,0)}$
 - $y_{den} = x'^{15} + k_{(4,14)} * x'^{14} + k_{(4,13)} * x'^{13} + \dots + k_{(4,0)}$

The constants used to compute x_{num} are as follows:

- $k_{(1,0)} = 0x11a05f2b1e833340b809101dd99815856b303e88a2d7005ff2627b56cdb4e2c85610c2d5f2e62d6eaeac166273$
- $k_{(1,1)} = 0x17294ed3e943ab2f0588bab22147a81c7c17e75b2f6a8417f565e33c70d1e86b4838f2a6f318c356e834eef1b3c$
- $k_{(1,2)} = 0xd54005db97678ec1d1048c5d10a9a1bce032473295983e56878e501ec68e25c958c3e3d2a09729fe0179f9dac9$
- $k_{(1,3)} = 0x1778e7166fcc6db74e0609d307e55412d7f5e4656a8dbf25f1b33289f1b330835336e25ce3107193c5b388641d9$
- $k_{(1,4)} = 0xe99726a3199f4436642b4b3e4118e5499db995a1257fb3f086eeb65982fac18985a286f301e77c451154ce9ac88$
- $k_{(1,5)} = 0x1630c3250d7313ff01d1201bf7a74ab5db3cb17dd952799b9ed3ab9097e68f90a0870d2dcae73d19cd13c1c66f$
- $k_{(1,6)} = 0xd6ed6553fe44d296a3726c38ae652bfb11586264f0f8ce19008e218f9c86b2a8da25128c1052ecadd7f225a139$
- $k_{(1,7)} = 0x17b81e7701abdbe2e8743884d1117e53356de5ab275b4db1a682c62ef0f2753339b7c8f8c8f475af9ccb5618e3f$

- $k_{(1,8)} =$
0x80d3cf1f9a78fc47b90b33563be990dc43b756ce79f5574a2c596c928c5d1de4fa295f296b74e956d71986a8497
- $k_{(1,9)} =$
0x169b1f8e1bcfa7c42e0c37515d138f22dd2ecb803a0c5c99676314baf4bb1b7fa3190b2edc0327797f241067be3
- $k_{(1,10)} =$
0x10321da079ce07e272d8ec09d2565b0dfa7dcccde6787f96d50af36003b14866f69b771f8c285decca67df3f160
- $k_{(1,11)} =$
0x6e08c248e260e70bd1e962381edee3d31d79d7e22c837bc23c0bf1bc24c6b68c24b1b80b64d391fa9c8ba2e8ba

The constants used to compute x_{den} are as follows:

- $k_{(2,0)} =$
0x8ca8d548cff19ae18b2e62f4bd3fa6f01d5ef4ba35b48ba9c9588617fc8ac62b558d681be343df8993cf9fa40d21
- $k_{(2,1)} =$
0x12561a5deb559c4348b4711298e536367041e8ca0cf0800c0126c2588c48bf5713daa8846cb026e9e5c8276ec82
- $k_{(2,2)} =$
0xb2962fe57a3225e8137e629bff2991f6f89416f5a718cd1fca64e00b11aceacd6a3d0967c94fedcfcc239ba5cb83e
- $k_{(2,3)} =$
0x3425581a58ae2fec83aafef7c40eb545b08243f16b1655154cca8abc28d6fd04976d5243eecf5c4130de8938dc62
- $k_{(2,4)} =$
0x13a8e162022914a80a6f1d5f43e7a07dffdfc759a12062bb8d6b44e833b306da9bd29ba81f35781d539d395b35
- $k_{(2,5)} =$
0xe7355f8e4e667b955390f7f0506c6e9395735e9ce9cad4d0a43bcef24b8982f7400d24bc4228f11c02df9a29f630
- $k_{(2,6)} =$
0x772caacf16936190f3e0c63e0596721570f5799af53a1894e2e073062aede9cea73b3538f0de06cec2574496ee84
- $k_{(2,7)} =$
0x14a7ac2a9d64a8b230b3f5b074cf01996e7f63c21bca68a81996e1cdf9822c580fa5b9489d11e2d311f7d99bbdc
- $k_{(2,8)} =$
0xa10ecf6ada54f825e920b3dafc7a3cce07f8d1d7161366b74100da67f39883503826692abba43704776ec3a79a1
- $k_{(2,9)} =$
0x95fc13ab9e92ad4476d6e3eb3a56680f682b4ee96f7d03776df533978f31c1593174e4b4b7865002d6384d168e

The constants used to compute y_{num} are as follows:

- $k_{(3,0)} =$
0x90d97c81ba24ee0259d1f094980dcfa11ad138e48a869522b52af6c956543d3cd0c7aee9b3ba3c2be984571970
- $k_{(3,1)} =$
0x134996a104ee5811d51036d776fb46831223e96c254f383d0f906343eb67ad34d6c56711962fa8bfe097e75a2e9
- $k_{(3,2)} =$
0xcc786baa966e66f4a384c86a3b49942552e2d658a31ce2c344be4b91400da7d26d521628b00523b8dfe240c72c
- $k_{(3,3)} =$
0x1f86376e8981c217898751ad8746757d42aa7b90eeb791c09e4a3ec03251cf9de405aba9ec61deca6355c77b0e
- $k_{(3,4)} =$
0x8cc03fdefe0ff135caf4fe2a21529c4195536f3ce50b879833fd221351adc2ee7f8dc099040a841b6daecf2e8fec

- $k_{(3,5)} =$
0x16603fca40634b6a2211e11db8f0a6a074a7d0d4afadb7bd76505c3d3ad5544e203f6326c95a807299b23ab136
- $k_{(3,6)} =$
0x4ab0b9bcfac1bbcb2c977d027796b3ce75bb8ca2be184cb5231413c4d634f3747a87ac2460f415ec961f8855fe9
- $k_{(3,7)} =$
0x987c8d5333ab86fde9926bd2ca6c674170a05bfe3bdd81ffd038da6c26c842642f64550fedfe935a15e4ca31870f
- $k_{(3,8)} =$
0x9fc4018bd96684be88c9e221e4da1bb8f3abd16679dc26c1e8b6e6a1f20cabe69d65201c78607a360370e577bd
- $k_{(3,9)} =$
0xe1bba7a1186bdb5223abde7ada14a23c42a0ca7915af6fe06985e7ed1e4d43b9b3f7055dd4eba6f2bafaaebca7
- $k_{(3,10)} =$
0x19713e47937cd1be0dfd0b8f1d43fb93cd2fcbcb6caf493fd1183e416389e61031bf3a5cce3fbafce813711ad011
- $k_{(3,11)} =$
0x18b46a908f36f6deb918c143fed2edcc523559b8aaf0c2462e6bfe7f911f643249d9cdf41b44d606ce07c8a4d007
- $k_{(3,12)} =$
0xb182cac101b9399d155096004f53f447aa7b12a3426b08ec02710e807b4633f06c851c1919211f20d4c04f00b97
- $k_{(3,13)} =$
0x245a394ad1eca9b72fc00ae7be315dc757b3b080d4c158013e6632d3c40659cc6cf90ad1c232a6442d9d3f5db9
- $k_{(3,14)} =$
0x5c129645e44cf1102a159f748c4a3fc5e673d81d7e86568d9ab0f5d396a7ce46ba1049b6579afb7866b1e715475
- $k_{(3,15)} =$
0x15e6be4e990f03ce4ea50b3b42df2eb5cb181d8f84965a3957add4fa95af01b2b665027efec01c7704b456be69c

The constants used to compute y_{den} are as follows:

- $k_{(4,0)} =$
0x16112c4c3a9c98b252181140fad0eae9601a6de578980be6eec3232b5be72e7a07f3688ef60c206d01479253b0f
- $k_{(4,1)} =$
0x1962d75c2381201e1a0cbd6c43c348b885c84ff731c4d59ca4a10356f453e01f78a4260763529e3532f6102c2e49
- $k_{(4,2)} =$
0x58df3306640da276faaae7d6e8eb15778c4855551ae7f310c35a5dd279cd2eca6757cd636f96f891e2538b53dbf
- $k_{(4,3)} =$
0x16b7d288798e5395f20d23bf89edb4d1d115c5dbddbcd30e123da489e726af41727364f2c28297ada8d26d984
- $k_{(4,4)} =$
0xbe0e079545f43e4b00cc912f8228ddcc6d19c9f0f69bbb0542eda0fc9dec916a20b15dc0fd2ededda39142311a5
- $k_{(4,5)} =$
0x8d9e5297186db2d9fb266eaac783182b70152c65550d881c5ecd87b6f0f5a6449f38db9dfa9cce202c6477faaf9
- $k_{(4,6)} =$
0x166007c08a99db2fc3ba8734ace9824b5eecdafa8d0cf8ef5dd365bc400a0051d5fa9c01a58b1fb93d1a1399126
- $k_{(4,7)} =$
0x16a3ef08be3ea7ea03bcddfabb6ff6ee5a4375efa1f4fd7feb34fd206357132b920f5b00801dee460ee415a1581
- $k_{(4,8)} =$
0x1866c8ed336c61231a1be54fd1d74cc4f9fb0ce4c6af5920abc5750c4bf39b4852cfe2f7bb9248836b233d9d555

- $k_{(4,9)} =$
0x167a55cda70a6e1cea820597d94a84903216f763e13d87bb5308592e7ea7d4fbc7385ea3d529b35e346ef48bb8
- $k_{(4,10)} =$
0x4d2f259eea405bd48f010a01ad2911d9c6dd039bb61a6290e591b36e636a5c871a5c29f4f83060400f8b49cba8
- $k_{(4,11)} =$
0xaccbb67481d033ff5852c1e48c50c477f94ff8aefce42d28c0f9a88cea7913516f968986f7ebbea9684b529e25610
- $k_{(4,12)} =$
0xad6b9514c767fe3c3613144b45f1496543346d98adf02267d5ceef9a00d9b8693000763e3b90ac11e99b138573
- $k_{(4,13)} =$
0x2660400eb2e4f3b628bdd0d53cd76f2bf565b94e72927c1cb748df27942480e420517bd8714cc80d1fadc1326e
- $k_{(4,14)} =$
0xe0fa1d816ddc03e6b24255e0d7819c171c40f65e273b853324efcd6356caa205ca2f570f13497804415473a1d63

E.3. 3-Isogeny Map for BLS12-381 G2

The 3-isogeny map from (x', y') on E' to (x, y) on E is given by the following rational functions:

- $x = x_num / x_den$, where
 - $x_num = k_{(1,3)} * x'^3 + k_{(1,2)} * x'^2 + k_{(1,1)} * x' + k_{(1,0)}$
 - $x_den = x'^2 + k_{(2,1)} * x' + k_{(2,0)}$
- $y = y' * y_num / y_den$, where
 - $y_num = k_{(3,3)} * x'^3 + k_{(3,2)} * x'^2 + k_{(3,1)} * x' + k_{(3,0)}$
 - $y_den = x'^3 + k_{(4,2)} * x'^2 + k_{(4,1)} * x' + k_{(4,0)}$

The constants used to compute x_num are as follows:

- $k_{(1,0)} =$
0x5c759507e8e333ebb5b7a9a47d7ed8532c52d39fd3a042a88b58423c50ae15d5c2638e343d9c71c6238aaaaaa
+
0x5c759507e8e333ebb5b7a9a47d7ed8532c52d39fd3a042a88b58423c50ae15d5c2638e343d9c71c6238aaaaaa
* I
- $k_{(1,1)} =$
0x11560bf17baa99bc32126fcd787c88f984f87adf7ae0c7f9a208c6b4f20a4181472aaa9cb8d555526a9ffffffc71
* I
- $k_{(1,2)} =$
0x11560bf17baa99bc32126fcd787c88f984f87adf7ae0c7f9a208c6b4f20a4181472aaa9cb8d555526a9ffffffc71
+
0x8ab05f8bdd54cde190937e76bc3e447cc27c3d6fbd7063fcd104635a790520c0a395554e5c6aaaa9354ffffffe3
* I
- $k_{(1,3)} =$
0x171d6541fa38ccaed6dea691f5fb614cb14b4e7f4e810aa22d6108f142b85757098e38d0f671c7188e2aaaaaaa

The constants used to compute x_{den} are as follows:

- $k_{(2,0)} =$
 $0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabfffeb153ffffb9feffffffffffaa63$
 $* I$
- $k_{(2,1)} = 0xc +$
 $0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabfffeb153ffffb9feffffffffffaa9f$
 $* I$

The constants used to compute y_{num} are as follows:

- $k_{(3,0)} =$
 $0x1530477c7ab4113b59a4c18b076d11930f7da5d4a07f649bf54439d87d27e500fc8c25ebf8c92f6812cfc71c71c6$
 $+$
 $0x1530477c7ab4113b59a4c18b076d11930f7da5d4a07f649bf54439d87d27e500fc8c25ebf8c92f6812cfc71c71c6$
 $* I$
- $k_{(3,1)} =$
 $0x5c759507e8e333ebb5b7a9a47d7ed8532c52d39fd3a042a88b58423c50ae15d5c2638e343d9c71c6238aaaaaa$
 $* I$
- $k_{(3,2)} =$
 $0x11560bf17baa99bc32126fcd787c88f984f87adf7ae0c7f9a208c6b4f20a4181472aaa9cb8d555526a9ffffffffffc71$
 $+$
 $0x8ab05f8bdd54cde190937e76bc3e447cc27c3d6fbd7063fcd104635a790520c0a395554e5c6aaaa9354ffffffffffe3$
 $* I$
- $k_{(3,3)} =$
 $0x124c9ad43b6cf79bfbf7043de3811ad0761b0f37a1e26286b0e977c69aa274524e79097a56dc4bd9e1b371c71c$

The constants used to compute y_{den} are as follows:

- $k_{(4,0)} =$
 $0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabfffeb153ffffb9feffffffffffa8fb$
 $+$
 $0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabfffeb153ffffb9feffffffffffa8fb$
 $* I$
- $k_{(4,1)} =$
 $0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabfffeb153ffffb9feffffffffffa9d3$
 $* I$
- $k_{(4,2)} = 0x12 +$
 $0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabfffeb153ffffb9feffffffffffaa99$
 $* I$

Appendix F. Straight-Line Implementations of Deterministic Mappings

This section gives straight-line implementations of the mappings of [Section 6](#). These implementations are generic, i.e., they are defined for any curve and field. [Appendix G](#) gives further implementations that are optimized for specific classes of curves and fields.

F.1. Shallue-van de Woestijne Method

This section gives a straight-line implementation of the Shallue-van de Woestijne method for any Weierstrass curve of the form given in [Section 6.6](#). See [Section 6.6.1](#) for information on the constants used in this mapping.

Note that the constant c_3 below **MUST** be chosen such that $\text{sgn}_0(c_3) = 0$. In other words, if the square-root computation returns a value c_x such that $\text{sgn}_0(c_x) = 1$, set $c_3 = -c_x$; otherwise, set $c_3 = c_x$.

```
map_to_curve_svdw(u)
```

```
Input: u, an element of F.
```

```
Output: (x, y), a point on E.
```

```
Constants:
```

```
1. c1 = g(Z)
2. c2 = -Z / 2
3. c3 = sqrt(-g(Z) * (3 * Z^2 + 4 * A)) # sgn0(c3) MUST equal 0
4. c4 = -4 * g(Z) / (3 * Z^2 + 4 * A)
```

```
Steps:
```

```
1. tv1 = u^2
2. tv1 = tv1 * c1
3. tv2 = 1 + tv1
4. tv1 = 1 - tv1
5. tv3 = tv1 * tv2
6. tv3 = inv0(tv3)
7. tv4 = u * tv1
8. tv4 = tv4 * tv3
9. tv4 = tv4 * c3
10. x1 = c2 - tv4
11. gx1 = x1^2
12. gx1 = gx1 + A
13. gx1 = gx1 * x1
14. gx1 = gx1 + B
15. e1 = is_square(gx1)
16. x2 = c2 + tv4
17. gx2 = x2^2
18. gx2 = gx2 + A
19. gx2 = gx2 * x2
20. gx2 = gx2 + B
21. e2 = is_square(gx2) AND NOT e1 # Avoid short-circuit logic ops
22. x3 = tv2^2
23. x3 = x3 * tv3
24. x3 = x3^2
25. x3 = x3 * c4
26. x3 = x3 + Z
27. x = CMOV(x3, x1, e1) # x = x1 if gx1 is square, else x = x3
28. x = CMOV(x, x2, e2) # x = x2 if gx2 is square and gx1 is not
29. gx = x^2
30. gx = gx + A
31. gx = gx * x
32. gx = gx + B
33. y = sqrt(gx)
34. e3 = sgn0(u) == sgn0(y)
35. y = CMOV(-y, y, e3) # Select correct sign of y
36. return (x, y)
```

F.2. Simplified SWU Method

This section gives a straight-line implementation of the Simplified SWU method for any Weierstrass curve of the form given in [Section 6.6](#). See [Section 6.6.2](#) for information on the constants used in this mapping.

This optimized, straight-line procedure applies to any base field. The `sqrt_ratio` subroutine is defined in [Appendix F.2.1](#).

```
map_to_curve_simple_swu(u)
```

Input: u , an element of F .

Output: (x, y) , a point on E .

Steps:

```
1. tv1 = u^2
2. tv1 = Z * tv1
3. tv2 = tv1^2
4. tv2 = tv2 + tv1
5. tv3 = tv2 + 1
6. tv3 = B * tv3
7. tv4 = CMOV(Z, -tv2, tv2 != 0)
8. tv4 = A * tv4
9. tv2 = tv3^2
10. tv6 = tv4^2
11. tv5 = A * tv6
12. tv2 = tv2 + tv5
13. tv2 = tv2 * tv3
14. tv6 = tv6 * tv4
15. tv5 = B * tv6
16. tv2 = tv2 + tv5
17. x = tv1 * tv3
18. (is_gx1_square, y1) = sqrt_ratio(tv2, tv6)
19. y = tv1 * u
20. y = y * y1
21. x = CMOV(x, tv3, is_gx1_square)
22. y = CMOV(y, y1, is_gx1_square)
23. e1 = sgn0(u) == sgn0(y)
24. y = CMOV(-y, y, e1)
25. x = x / tv4
26. return (x, y)
```

F.2.1. `sqrt_ratio` Subroutine

This section defines three variants of the `sqrt_ratio` subroutine used by the above procedure. The first variant can be used with any field; the others are optimized versions for specific fields.

The routines given in this section depend on the constant Z from the Simplified SWU map. For correctness, `sqrt_ratio` and `map_to_curve_simple_swu` **MUST** use the same value for Z .

F.2.1.1. sqrt_ratio for Any Field

```
sqrt_ratio(u, v)
```

Parameters:

- F, a finite field of characteristic p and order $q = p^m$.
- Z, the constant from the Simplified SWU map.

Input: u and v, elements of F, where $v \neq 0$.

Output: (b, y), where

- b = True and $y = \sqrt{u / v}$ if (u / v) is square in F, and
- b = False and $y = \sqrt{Z * (u / v)}$ otherwise.

Constants:

1. c_1 , the largest integer such that 2^{c_1} divides $q - 1$.
2. $c_2 = (q - 1) / (2^{c_1})$ # Integer arithmetic
3. $c_3 = (c_2 - 1) / 2$ # Integer arithmetic
4. $c_4 = 2^{c_1} - 1$ # Integer arithmetic
5. $c_5 = 2^{(c_1 - 1)}$ # Integer arithmetic
6. $c_6 = Z^{c_2}$
7. $c_7 = Z^{((c_2 + 1) / 2)}$

Procedure:

1. $tv_1 = c_6$
2. $tv_2 = v^{c_4}$
3. $tv_3 = tv_2^2$
4. $tv_3 = tv_3 * v$
5. $tv_5 = u * tv_3$
6. $tv_5 = tv_5^{c_3}$
7. $tv_5 = tv_5 * tv_2$
8. $tv_2 = tv_5 * v$
9. $tv_3 = tv_5 * u$
10. $tv_4 = tv_3 * tv_2$
11. $tv_5 = tv_4^{c_5}$
12. $isQR = tv_5 == 1$
13. $tv_2 = tv_3 * c_7$
14. $tv_5 = tv_4 * tv_1$
15. $tv_3 = CMOV(tv_2, tv_3, isQR)$
16. $tv_4 = CMOV(tv_5, tv_4, isQR)$
17. for i in $(c_1, c_1 - 1, \dots, 2)$:
 18. $tv_5 = i - 2$
 19. $tv_5 = 2^{tv_5}$
 20. $tv_5 = tv_4^{tv_5}$
 21. $e_1 = tv_5 == 1$
 22. $tv_2 = tv_3 * tv_1$
 23. $tv_1 = tv_1 * tv_1$
 24. $tv_5 = tv_4 * tv_1$
 25. $tv_3 = CMOV(tv_2, tv_3, e_1)$
 26. $tv_4 = CMOV(tv_5, tv_4, e_1)$
27. return (isQR, tv_3)

F.2.1.2. Optimized `sqrt_ratio` for $q = 3 \bmod 4$

```
sqrt_ratio_3mod4(u, v)
```

Parameters:

- F , a finite field of characteristic p and order $q = p^m$, where $q = 3 \bmod 4$.
- Z , the constant from the Simplified SWU map.

Input: u and v , elements of F , where $v \neq 0$.

Output: (b, y) , where

- $b = \text{True}$ and $y = \text{sqrt}(u / v)$ if (u / v) is square in F , and
- $b = \text{False}$ and $y = \text{sqrt}(Z * (u / v))$ otherwise.

Constants:

1. $c1 = (q - 3) / 4$ # Integer arithmetic
2. $c2 = \text{sqrt}(-Z)$

Procedure:

1. $tv1 = v^2$
2. $tv2 = u * v$
3. $tv1 = tv1 * tv2$
4. $y1 = tv1^{c1}$
5. $y1 = y1 * tv2$
6. $y2 = y1 * c2$
7. $tv3 = y1^2$
8. $tv3 = tv3 * v$
9. $isQR = tv3 == u$
10. $y = \text{CMOV}(y2, y1, isQR)$
11. return $(isQR, y)$

F.2.1.3. Optimized `sqrt_ratio` for $q = 5 \bmod 8$

```
sqrt_ratio_5mod8(u, v)
```

Parameters:

- F , a finite field of characteristic p and order $q = p^m$, where $q = 5 \bmod 8$.
- Z , the constant from the Simplified SWU map.

Input: u and v , elements of F , where $v \neq 0$.

Output: (b, y) , where

- $b = \text{True}$ and $y = \text{sqrt}(u / v)$ if (u / v) is square in F , and
- $b = \text{False}$ and $y = \text{sqrt}(Z * (u / v))$ otherwise.

Constants:

1. $c1 = (q - 5) / 8$
2. $c2 = \text{sqrt}(-1)$
3. $c3 = \text{sqrt}(Z / c2)$

Steps:

1. $tv1 = v^2$
2. $tv2 = tv1 * v$
3. $tv1 = tv1^2$
4. $tv2 = tv2 * u$
5. $tv1 = tv1 * tv2$
6. $y1 = tv1^{c1}$
7. $y1 = y1 * tv2$
8. $tv1 = y1 * c2$
9. $tv2 = tv1^2$
10. $tv2 = tv2 * v$
11. $e1 = tv2 == u$
12. $y1 = \text{CMOV}(y1, tv1, e1)$
13. $tv2 = y1^2$
14. $tv2 = tv2 * v$
15. $isQR = tv2 == u$
16. $y2 = y1 * c3$
17. $tv1 = y2 * c2$
18. $tv2 = tv1^2$
19. $tv2 = tv2 * v$
20. $tv3 = Z * u$
21. $e2 = tv2 == tv3$
22. $y2 = \text{CMOV}(y2, tv1, e2)$
23. $y = \text{CMOV}(y2, y1, isQR)$
24. return $(isQR, y)$

F.3. Elligator 2 Method

This section gives a straight-line implementation of the Elligator 2 method for any Montgomery curve of the form given in [Section 6.7](#). See [Section 6.7.1](#) for information on the constants used in this mapping.

[Appendix G.2](#) gives optimized straight-line procedures that apply to specific classes of curves and base fields, including `curve25519` and `curve448` [[RFC7748](#)].

```

map_to_curve_elligator2(u)

Input: u, an element of F.
Output: (s, t), a point on M.

Constants:
1.  c1 = J / K
2.  c2 = 1 / K^2

Steps:
1.  tv1 = u^2
2.  tv1 = Z * tv1           # Z * u^2
3.  e1 = tv1 == -1         # exceptional case: Z * u^2 == -1
4.  tv1 = CMOV(tv1, 0, e1) # if tv1 == -1, set tv1 = 0
5.  x1 = tv1 + 1
6.  x1 = inv0(x1)
7.  x1 = -c1 * x1          # x1 = -(J / K) / (1 + Z * u^2)
8.  gx1 = x1 + c1
9.  gx1 = gx1 * x1
10. gx1 = gx1 + c2
11. gx1 = gx1 * x1         # gx1 = x1^3 + (J / K) * x1^2 + x1 / K^2
12. x2 = -x1 - c1
13. gx2 = tv1 * gx1
14. e2 = is_square(gx1)    # If is_square(gx1)
15.  x = CMOV(x2, x1, e2)   # then x = x1, else x = x2
16. y2 = CMOV(gx2, gx1, e2) # then y2 = gx1, else y2 = gx2
17.  y = sqrt(y2)
18. e3 = sgn0(y) == 1
19.  y = CMOV(y, -y, e2 XOR e3) # fix sign of y
20.  s = x * K
21.  t = y * K
22. return (s, t)

```

Appendix G. Curve-Specific Optimized Sample Code

This section gives sample implementations optimized for some of the elliptic curves listed in [Section 8](#). Sample Sage code [[SAGE](#)] for each algorithm can also be found in [[hash2curve-repo](#)].

G.1. Interface and Projective Coordinate Systems

The sample code in this section uses a different interface than the mappings of [Section 6](#). Specifically, each mapping function in this section has the following signature:

```
(xn, xd, yn, yd) = map_to_curve(u)
```

The resulting affine point (x, y) is given by $(xn / xd, yn / yd)$.

The reason for this modified interface is that it enables further optimizations when working with points in a projective coordinate system. This is desirable, for example, when the resulting point will be immediately multiplied by a scalar, since most scalar multiplication algorithms operate on projective points.

Projective coordinates are also useful when implementing random-oracle encodings ([Section 3](#)). One reason is that, in general, point addition is faster using projective coordinates. Another reason is that, for Weierstrass curves, projective coordinates allow using complete addition formulas [[RCB16](#)]. This is especially convenient when implementing a constant-time encoding, because it eliminates the need for a special case when $Q_0 == Q_1$, which incomplete addition formulas usually do not handle.

The following are two commonly used projective coordinate systems and the corresponding conversions:

- A point (X, Y, Z) in homogeneous projective coordinates corresponds to the affine point $(x, y) = (X/Z, Y/Z)$; the inverse conversion is given by $(X, Y, Z) = (x, y, 1)$. To convert (x_n, x_d, y_n, y_d) to homogeneous projective coordinates, compute $(X, Y, Z) = (x_n * y_d, y_n * x_d, x_d * y_d)$.
- A point (X', Y', Z') in Jacobian projective coordinates corresponds to the affine point $(x, y) = (X'/Z'^2, Y'/Z'^3)$; the inverse conversion is given by $(X', Y', Z') = (x, y, 1)$. To convert (x_n, x_d, y_n, y_d) to Jacobian projective coordinates, compute $(X', Y', Z') = (x_n * x_d * y_d^2, y_n * y_d^2 * x_d^3, x_d * y_d)$.

G.2. Elligator 2

G.2.1. curve25519 ($q = 5 \pmod{8}$, $K = 1$)

The following is a straight-line implementation of Elligator 2 for curve25519 [[RFC7748](#)] as specified in [Section 8.5](#).

This implementation can also be used for any Montgomery curve with $K = 1$ over $\text{GF}(q)$ where $q = 5 \pmod{8}$.

```
map_to_curve_elligator2_curve25519(u)
```

Input: u , an element of F .

Output: (x_n, x_d, y_n, y_d) such that $(x_n / x_d, y_n / y_d)$ is a point on curve25519.

Constants:

```
1. c1 = (q + 3) / 8      # Integer arithmetic
2. c2 = 2^c1
3. c3 = sqrt(-1)
4. c4 = (q - 5) / 8     # Integer arithmetic
```

Steps:

```
1. tv1 = u^2
2. tv1 = 2 * tv1
3.  xd = tv1 + 1      # Nonzero: -1 is square (mod p), tv1 is not
4.  x1n = -J          # x1 = x1n / xd = -J / (1 + 2 * u^2)
5.  tv2 = xd^2
6.  gxd = tv2 * xd   # gxd = xd^3
7.  gx1 = J * tv1    # x1n + J * xd
8.  gx1 = gx1 * x1n  # x1n^2 + J * x1n * xd
9.  gx1 = gx1 + tv2  # x1n^2 + J * x1n * xd + xd^2
10. gx1 = gx1 * x1n # x1n^3 + J * x1n^2 * xd + x1n * xd^2
11. tv3 = gxd^2
```



```

12. tv2 = tv3^2           # gxd^4
13. tv3 = tv3 * gxd      # gxd^3
14. tv3 = tv3 * gx1      # gx1 * gxd^3
15. tv2 = tv2 * tv3      # gx1 * gxd^7
16. y11 = tv2^c4         # (gx1 * gxd^7)^((p - 5) / 8)
17. y11 = y11 * tv3      # gx1 * gxd^3 * (gx1 * gxd^7)^((p - 5) / 8)
18. y12 = y11 * c3
19. tv2 = y11^2
20. tv2 = tv2 * gxd
21. e1 = tv2 == gx1
22. y1 = CM0V(y12, y11, e1) # If g(x1) is square, this is its sqrt
23. x2n = x1n * tv1      # x2 = x2n / xd = 2 * u^2 * x1n / xd
24. y21 = y11 * u
25. y21 = y21 * c2
26. y22 = y21 * c3
27. gx2 = gx1 * tv1      # g(x2) = gx2 / gxd = 2 * u^2 * g(x1)
28. tv2 = y21^2
29. tv2 = tv2 * gxd
30. e2 = tv2 == gx2
31. y2 = CM0V(y22, y21, e2) # If g(x2) is square, this is its sqrt
32. tv2 = y1^2
33. tv2 = tv2 * gxd
34. e3 = tv2 == gx1
35. xn = CM0V(x2n, x1n, e3) # If e3, x = x1, else x = x2
36. y = CM0V(y2, y1, e3)   # If e3, y = y1, else y = y2
37. e4 = sgn0(y) == 1     # Fix sign of y
38. y = CM0V(y, -y, e3 XOR e4)
39. return (xn, xd, y, 1)

```

G.2.2. edwards25519

The following is a straight-line implementation of Elligator 2 for edwards25519 [RFC7748] as specified in Section 8.5. The subroutine `map_to_curve_elligator2_curve25519` is defined in Appendix G.2.1.

Note that the sign of the constant `c1` below is chosen as specified in Section 6.8.1, i.e., applying the rational map to the edwards25519 base point yields the curve25519 base point (see erratum [Err4730]).

```
map_to_curve_elligator2_edwards25519(u)
```

Input: u , an element of F .

Output: (x_n, x_d, y_n, y_d) such that $(x_n / x_d, y_n / y_d)$ is a point on edwards25519.

Constants:

1. $c_1 = \text{sqrt}(-486664)$ # $\text{sgn}_0(c_1)$ MUST equal 0

Steps:

```
1. (xMn, xMd, yMn, yMd) = map_to_curve_elligator2_curve25519(u)
2. xn = xMn * yMd
3. xn = xn * c1
4. xd = xMd * yMn    # xn / xd = c1 * xM / yM
5. yn = xMn - xMd
6. yd = xMn + xMd    # (n / d - 1) / (n / d + 1) = (n - d) / (n + d)
7. tv1 = xd * yd
8. e = tv1 == 0
9. xn = CMOV(xn, 0, e)
10. xd = CMOV(xd, 1, e)
11. yn = CMOV(yn, 1, e)
12. yd = CMOV(yd, 1, e)
13. return (xn, xd, yn, yd)
```

G.2.3. curve448 ($q = 3 \pmod{4}$, $K = 1$)

The following is a straight-line implementation of Elligator 2 for curve448 [RFC7748] as specified in Section 8.6.

This implementation can also be used for any Montgomery curve with $K = 1$ over $\text{GF}(q)$ where $q = 3 \pmod{4}$.

```
map_to_curve_elligator2_curve448(u)
```

Input: u , an element of F .

Output: (x_n, x_d, y_n, y_d) such that $(x_n / x_d, y_n / y_d)$ is a point on curve448.

Constants:

```
1. c1 = (q - 3) / 4      # Integer arithmetic
```

Steps:

```
1. tv1 = u^2
2. e1 = tv1 == 1
3. tv1 = CMOV(tv1, 0, e1) # If Z * u^2 == -1, set tv1 = 0
4. xd = 1 - tv1
5. x1n = -J
6. tv2 = xd^2
7. gxd = tv2 * xd      # gxd = xd^3
8. gx1 = -J * tv1     # x1n + J * xd
9. gx1 = gx1 * x1n    # x1n^2 + J * x1n * xd
10. gx1 = gx1 + tv2   # x1n^2 + J * x1n * xd + xd^2
11. gx1 = gx1 * x1n   # x1n^3 + J * x1n^2 * xd + x1n * xd^2
12. tv3 = gxd^2
13. tv2 = gx1 * gxd   # gx1 * gxd
14. tv3 = tv3 * tv2   # gx1 * gxd^3
15. y1 = tv3^c1      # (gx1 * gxd^3)^((p - 3) / 4)
16. y1 = y1 * tv2    # gx1 * gxd * (gx1 * gxd^3)^((p - 3) / 4)
17. x2n = -tv1 * x1n # x2 = x2n / xd = -1 * u^2 * x1n / xd
18. y2 = y1 * u
19. y2 = CMOV(y2, 0, e1)
20. tv2 = y1^2
21. tv2 = tv2 * gxd
22. e2 = tv2 == gx1
23. xn = CMOV(x2n, x1n, e2) # If e2, x = x1, else x = x2
24. y = CMOV(y2, y1, e2) # If e2, y = y1, else y = y2
25. e3 = sgn0(y) == 1 # Fix sign of y
26. y = CMOV(y, -y, e2 XOR e3)
27. return (xn, xd, y, 1)
```

G.2.4. edwards448

The following is a straight-line implementation of Elligator 2 for edwards448 [RFC7748] as specified in Section 8.6. The subroutine `map_to_curve_elligator2_curve448` is defined in Appendix G.2.3.

```
map_to_curve_elligator2_edwards448(u)
```

Input: u , an element of F .

Output: (x_n, x_d, y_n, y_d) such that $(x_n / x_d, y_n / y_d)$ is a point on edwards448.

Steps:

```

1.  $(x_n, x_d, y_n, y_d) = \text{map\_to\_curve\_elligator2\_curve448}(u)$ 
2.  $x_n2 = x_n^2$ 
3.  $x_d2 = x_d^2$ 
4.  $x_d4 = x_d2^2$ 
5.  $y_n2 = y_n^2$ 
6.  $y_d2 = y_d^2$ 
7.  $xEn = x_n2 - x_d2$ 
8.  $tv2 = xEn - x_d2$ 
9.  $xEn = xEn * x_d2$ 
10.  $xEn = xEn * y_d$ 
11.  $xEn = xEn * y_n$ 
12.  $xEn = xEn * 4$ 
13.  $tv2 = tv2 * x_n2$ 
14.  $tv2 = tv2 * y_d2$ 
15.  $tv3 = 4 * y_n2$ 
16.  $tv1 = tv3 + y_d2$ 
17.  $tv1 = tv1 * x_d4$ 
18.  $xEd = tv1 + tv2$ 
19.  $tv2 = tv2 * x_n$ 
20.  $tv4 = x_n * x_d4$ 
21.  $yEn = tv3 - y_d2$ 
22.  $yEn = yEn * tv4$ 
23.  $yEn = yEn - tv2$ 
24.  $tv1 = x_n2 + x_d2$ 
25.  $tv1 = tv1 * x_d2$ 
26.  $tv1 = tv1 * x_d$ 
27.  $tv1 = tv1 * y_n2$ 
28.  $tv1 = -2 * tv1$ 
29.  $yEd = tv2 + tv1$ 
30.  $tv4 = tv4 * y_d2$ 
31.  $yEd = yEd + tv4$ 
32.  $tv1 = xEd * yEd$ 
33.  $e = tv1 == 0$ 
34.  $xEn = \text{CMOV}(xEn, 0, e)$ 
35.  $xEd = \text{CMOV}(xEd, 1, e)$ 
36.  $yEn = \text{CMOV}(yEn, 1, e)$ 
37.  $yEd = \text{CMOV}(yEd, 1, e)$ 
38. return  $(xEn, xEd, yEn, yEd)$ 

```

G.2.5. Montgomery Curves with $q = 3 \pmod{4}$

The following is a straight-line implementation of Elligator 2 that applies to any Montgomery curve defined over $GF(q)$ where $q = 3 \pmod{4}$.

For curves where $K = 1$, the implementation given in [Appendix G.2.3](#) gives identical results with slightly reduced cost.

```
map_to_curve_elligator2_3mod4(u)
```

Input: u , an element of F .

Output: (x_n, x_d, y_n, y_d) such that $(x_n / x_d, y_n / y_d)$ is a point on the target curve.

Constants:

1. $c_1 = (q - 3) / 4$ # Integer arithmetic
2. $c_2 = K^2$

Steps:

1. $tv_1 = u^2$
2. $e_1 = tv_1 == 1$
3. $tv_1 = CMOV(tv_1, 0, e_1)$ # If $Z * u^2 == -1$, set $tv_1 = 0$
4. $x_d = 1 - tv_1$
5. $x_d = x_d * K$
6. $x_{1n} = -J$ # $x_1 = x_{1n} / x_d = -J / (K * (1 + 2 * u^2))$
7. $tv_2 = x_d^2$
8. $g_{xd} = tv_2 * x_d$
9. $g_{xd} = g_{xd} * c_2$ # $g_{xd} = x_d^3 * K^2$
10. $g_{x1} = x_{1n} * K$
11. $tv_3 = x_d * J$
12. $tv_3 = g_{x1} + tv_3$ # $x_{1n} * K + x_d * J$
13. $g_{x1} = g_{x1} * tv_3$ # $K^2 * x_{1n}^2 + J * K * x_{1n} * x_d$
14. $g_{x1} = g_{x1} + tv_2$ # $K^2 * x_{1n}^2 + J * K * x_{1n} * x_d + x_d^2$
15. $g_{x1} = g_{x1} * x_{1n}$ # $K^2 * x_{1n}^3 + J * K * x_{1n}^2 * x_d + x_{1n} * x_d^2$
16. $tv_3 = g_{xd}^2$
17. $tv_2 = g_{x1} * g_{xd}$ # $g_{x1} * g_{xd}$
18. $tv_3 = tv_3 * tv_2$ # $g_{x1} * g_{xd}^3$
19. $y_1 = tv_3^{c_1}$ # $(g_{x1} * g_{xd}^3)^{((q - 3) / 4)}$
20. $y_1 = y_1 * tv_2$ # $g_{x1} * g_{xd} * (g_{x1} * g_{xd}^3)^{((q - 3) / 4)}$
21. $x_{2n} = -tv_1 * x_{1n}$ # $x_2 = x_{2n} / x_d = -1 * u^2 * x_{1n} / x_d$
22. $y_2 = y_1 * u$
23. $y_2 = CMOV(y_2, 0, e_1)$
24. $tv_2 = y_1^2$
25. $tv_2 = tv_2 * g_{xd}$
26. $e_2 = tv_2 == g_{x1}$
27. $x_n = CMOV(x_{2n}, x_{1n}, e_2)$ # If e_2 , $x = x_1$, else $x = x_2$
28. $x_n = x_n * K$
29. $y = CMOV(y_2, y_1, e_2)$ # If e_2 , $y = y_1$, else $y = y_2$
30. $e_3 = \text{sgn}_0(y) == 1$ # Fix sign of y
31. $y = CMOV(y, -y, e_2 \text{ XOR } e_3)$
32. $y = y * K$
33. return $(x_n, x_d, y, 1)$

G.2.6. Montgomery Curves with $q = 5 \pmod{8}$

The following is a straight-line implementation of Elligator 2 that applies to any Montgomery curve defined over $GF(q)$ where $q = 5 \pmod{8}$.

For curves where $K = 1$, the implementation given in [Appendix G.2.1](#) gives identical results with slightly reduced cost.

```
map_to_curve_elligator2_5mod8(u)
```

Input: u , an element of F .

Output: (x_n, x_d, y_n, y_d) such that $(x_n / x_d, y_n / y_d)$ is a point on the target curve.

Constants:

1. $c_1 = (q + 3) / 8$ # Integer arithmetic
2. $c_2 = 2^{c_1}$
3. $c_3 = \text{sqrt}(-1)$
4. $c_4 = (q - 5) / 8$ # Integer arithmetic
5. $c_5 = K^2$

Steps:

1. $tv_1 = u^2$
2. $tv_1 = 2 * tv_1$
3. $x_d = tv_1 + 1$ # Nonzero: -1 is square (mod p), tv_1 is not
4. $x_d = x_d * K$
5. $x_{1n} = -J$ # $x_1 = x_{1n} / x_d = -J / (K * (1 + 2 * u^2))$
6. $tv_2 = x_d^2$
7. $g_{xd} = tv_2 * x_d$
8. $g_{xd} = g_{xd} * c_5$ # $g_{xd} = x_d^3 * K^2$
9. $g_{x1} = x_{1n} * K$
10. $tv_3 = x_d * J$
11. $tv_3 = g_{x1} + tv_3$ # $x_{1n} * K + x_d * J$
12. $g_{x1} = g_{x1} * tv_3$ # $K^2 * x_{1n}^2 + J * K * x_{1n} * x_d$
13. $g_{x1} = g_{x1} + tv_2$ # $K^2 * x_{1n}^2 + J * K * x_{1n} * x_d + x_d^2$
14. $g_{x1} = g_{x1} * x_{1n}$ # $K^2 * x_{1n}^3 + J * K * x_{1n}^2 * x_d + x_{1n} * x_d^2$
15. $tv_3 = g_{xd}^2$
16. $tv_2 = tv_3^2$ # g_{xd}^4
17. $tv_3 = tv_3 * g_{xd}$ # g_{xd}^3
18. $tv_3 = tv_3 * g_{x1}$ # $g_{x1} * g_{xd}^3$
19. $tv_2 = tv_2 * tv_3$ # $g_{x1} * g_{xd}^7$
20. $y_{11} = tv_2^{c_4}$ # $(g_{x1} * g_{xd}^7)^{((q - 5) / 8)}$
21. $y_{11} = y_{11} * tv_3$ # $g_{x1} * g_{xd}^3 * (g_{x1} * g_{xd}^7)^{((q - 5) / 8)}$
22. $y_{12} = y_{11} * c_3$
23. $tv_2 = y_{11}^2$
24. $tv_2 = tv_2 * g_{xd}$
25. $e_1 = tv_2 == g_{x1}$
26. $y_1 = \text{CMOV}(y_{12}, y_{11}, e_1)$ # If $g(x_1)$ is square, this is its sqrt
27. $x_{2n} = x_{1n} * tv_1$ # $x_2 = x_{2n} / x_d = 2 * u^2 * x_{1n} / x_d$
28. $y_{21} = y_{11} * u$
29. $y_{21} = y_{21} * c_2$
30. $y_{22} = y_{21} * c_3$
31. $g_{x2} = g_{x1} * tv_1$ # $g(x_2) = g_{x2} / g_{xd} = 2 * u^2 * g(x_1)$
32. $tv_2 = y_{21}^2$
33. $tv_2 = tv_2 * g_{xd}$
34. $e_2 = tv_2 == g_{x2}$
35. $y_2 = \text{CMOV}(y_{22}, y_{21}, e_2)$ # If $g(x_2)$ is square, this is its sqrt
36. $tv_2 = y_1^2$
37. $tv_2 = tv_2 * g_{xd}$
38. $e_3 = tv_2 == g_{x1}$
39. $x_n = \text{CMOV}(x_{2n}, x_{1n}, e_3)$ # If e_3 , $x = x_1$, else $x = x_2$
40. $x_n = x_n * K$
41. $y = \text{CMOV}(y_2, y_1, e_3)$ # If e_3 , $y = y_1$, else $y = y_2$
42. $e_4 = \text{sgn}_0(y) == 1$ # Fix sign of y
43. $y = \text{CMOV}(y, -y, e_3 \text{ XOR } e_4)$
44. $y = y * K$
45. return $(x_n, x_d, y, 1)$

G.3. Cofactor Clearing for BLS12-381 G2

The curve BLS12-381, whose parameters are defined in [Section 8.8.2](#), admits an efficiently computable endomorphism, ψ , that can be used to speed up cofactor clearing for G2 [[SBCDK09](#)] [[FKR11](#)] [[BP17](#)] (see also [Section 7](#)). This section implements the endomorphism ψ and a fast cofactor clearing method described by Budroni and Pintore [[BP17](#)].

The functions in this section operate on points whose coordinates are represented as ratios, i.e., (x_n, x_d, y_n, y_d) corresponds to the point $(x_n / x_d, y_n / y_d)$; see [Appendix G.1](#) for further discussion of projective coordinates. When points are represented in affine coordinates, one can simply ignore the denominators ($x_d == 1$ and $y_d == 1$).

The following function computes the Frobenius endomorphism for an element of $F = GF(p^2)$ with basis $(1, I)$, where $I^2 + 1 == 0$ in F . (This is the base field of the elliptic curve E defined in [Section 8.8.2](#).)

```
frobenius(x)
```

Input: x , an element of $GF(p^2)$.

Output: a , an element of $GF(p^2)$.

Notation: $x = x_0 + I * x_1$, where x_0 and x_1 are elements of $GF(p)$.

Steps:

1. $a = x_0 - I * x_1$
2. return a

The following function computes the endomorphism ψ for points on the elliptic curve E defined in [Section 8.8.2](#).

```
psi(xn, xd, yn, yd)
```

Input: P , a point $(x_n / x_d, y_n / y_d)$ on the curve E (see above).

Output: Q , a point on the same curve.

Constants:

1. $c_1 = 1 / (1 + I)^{((p - 1) / 3)}$ # in $GF(p^2)$
2. $c_2 = 1 / (1 + I)^{((p - 1) / 2)}$ # in $GF(p^2)$

Steps:

1. $q_{xn} = c_1 * \text{frobenius}(x_n)$
2. $q_{xd} = \text{frobenius}(x_d)$
3. $q_{yn} = c_2 * \text{frobenius}(y_n)$
4. $q_{yd} = \text{frobenius}(y_d)$
5. return $(q_{xn}, q_{xd}, q_{yn}, q_{yd})$

The following function efficiently computes $\psi(\psi(P))$.

```
psi2(xn, xd, yn, yd)
```

Input: P, a point (xn / xd, yn / yd) on the curve E (see above).
Output: Q, a point on the same curve.

Constants:

```
1. c1 = 1 / 2^((p - 1) / 3) # in GF(p^2)
```

Steps:

```
1. qxn = c1 * xn
2. qyn = -yn
3. return (qxn, xd, qyn, yd)
```

The following function maps any point on the elliptic curve E ([Section 8.8.2](#)) into the prime-order subgroup G2. This function returns a point equal to $h_{\text{eff}} * P$, where h_{eff} is the parameter given in [Section 8.8.2](#).

```
clear_cofactor_bls12381_g2(P)
```

Input: P, a point (xn / xd, yn / yd) on the curve E (see above).
Output: Q, a point in the subgroup G2 of BLS12-381.

Constants:

```
1. c1 = -15132376222941642752 # the BLS parameter for BLS12-381
# i.e., -0xd201000000010000
```

Notation: in this procedure, + and - represent elliptic curve point addition and subtraction, respectively, and * represents scalar multiplication.

Steps:

```
1. t1 = c1 * P
2. t2 = psi(P)
3. t3 = 2 * P
4. t3 = psi2(t3)
5. t3 = t3 - t2
6. t2 = t1 + t2
7. t2 = c1 * t2
8. t3 = t3 + t2
9. t3 = t3 - t1
10. Q = t3 - P
11. return Q
```

Appendix H. Scripts for Parameter Generation

This section gives Sage scripts [[SAGE](#)] used to generate parameters for the mappings of [Section 6](#).

H.1. Finding Z for the Shallue-van de Woestijne Map

The below function outputs an appropriate Z for the Shallue-van de Woestijne map ([Section 6.6.1](#)).


```

# Arguments:
# - F, a field object, e.g., F = GF(2^521 - 1)
# - A and B, the coefficients of the curve  $y^2 = x^3 + A * x + B$ 
def find_z_svdw(F, A, B, init_ctr=1):
    g = lambda x: F(x)^3 + F(A) * F(x) + F(B)
    h = lambda Z: -(F(3) * Z^2 + F(4) * A) / (F(4) * g(Z))
    # NOTE: if init_ctr=1 fails to find Z, try setting it to F.gen()
    ctr = init_ctr
    while True:
        for Z_cand in (F(ctr), F(-ctr)):
            # Criterion 1:
            # g(Z) != 0 in F.
            if g(Z_cand) == F(0):
                continue
            # Criterion 2:
            #  $-(3 * Z^2 + 4 * A) / (4 * g(Z)) != 0$  in F.
            if h(Z_cand) == F(0):
                continue
            # Criterion 3:
            #  $-(3 * Z^2 + 4 * A) / (4 * g(Z))$  is square in F.
            if not is_square(h(Z_cand)):
                continue
            # Criterion 4:
            # At least one of g(Z) and g(-Z / 2) is square in F.
            if is_square(g(Z_cand)) or is_square(g(-Z_cand / F(2))):
                return Z_cand
        ctr += 1

```

H.2. Finding Z for Simplified SWU

The below function outputs an appropriate Z for the Simplified SWU map ([Section 6.6.2](#)).

```

# Arguments:
# - F, a field object, e.g., F = GF(2^521 - 1)
# - A and B, the coefficients of the curve  $y^2 = x^3 + A * x + B$ 
def find_z_sswu(F, A, B):
    R.<xx> = F[] # Polynomial ring over F
    g = xx^3 + F(A) * xx + F(B) #  $y^2 = g(x) = x^3 + A * x + B$ 
    ctr = F.gen()
    while True:
        for Z_cand in (F(ctr), F(-ctr)):
            # Criterion 1: Z is non-square in F.
            if is_square(Z_cand):
                continue
            # Criterion 2: Z != -1 in F.
            if Z_cand == F(-1):
                continue
            # Criterion 3: g(x) - Z is irreducible over F.
            if not (g - Z_cand).is_irreducible():
                continue
            # Criterion 4: g(B / (Z * A)) is square in F.
            if is_square(g(B / (Z_cand * A))):
                return Z_cand
        ctr += 1

```

H.3. Finding Z for Elligator 2

The below function outputs an appropriate Z for the Elligator 2 map ([Section 6.7.1](#)).

```
# Argument:
# - F, a field object, e.g., F = GF(2^255 - 19)
def find_z_ell2(F):
    ctr = F.gen()
    while True:
        for Z_cand in (F(ctr), F(-ctr)):
            # Z must be a non-square in F.
            if is_square(Z_cand):
                continue
            return Z_cand
        ctr += 1
```

Appendix I. sqrt and is_square Functions

This section defines special-purpose sqrt functions for the three most common cases, $q = 3 \pmod{4}$, $q = 5 \pmod{8}$, and $q = 9 \pmod{16}$, plus a generic constant-time algorithm that works for any prime modulus.

In addition, it gives an optimized is_square method for $GF(p^2)$.

I.1. sqrt for $q = 3 \pmod{4}$

```
sqrt_3mod4(x)
```

Parameters:

- F, a finite field of characteristic p and order $q = p^m$.

Input: x, an element of F.

Output: z, an element of F such that $(z^2) == x$, if x is square in F.

Constants:

1. $c1 = (q + 1) / 4$ # Integer arithmetic

Procedure:

1. return x^{c1}

I.2. sqrt for $q = 5 \pmod{8}$

sqrt_5mod8(x)

Parameters:

- F, a finite field of characteristic p and order $q = p^m$.

Input: x, an element of F.

Output: z, an element of F such that $(z^2) == x$, if x is square in F.

Constants:

1. $c1 = \text{sqrt}(-1)$ in F, i.e., $(c1^2) == -1$ in F
2. $c2 = (q + 3) / 8$ # Integer arithmetic

Procedure:

1. $tv1 = x^{c2}$
2. $tv2 = tv1 * c1$
3. $e = (tv1^2) == x$
4. $z = \text{CMOV}(tv2, tv1, e)$
5. return z

I.3. sqrt for $q = 9 \pmod{16}$

sqrt_9mod16(x)

Parameters:

- F, a finite field of characteristic p and order $q = p^m$.

Input: x, an element of F.

Output: z, an element of F such that $(z^2) == x$, if x is square in F.

Constants:

1. $c1 = \text{sqrt}(-1)$ in F, i.e., $(c1^2) == -1$ in F
2. $c2 = \text{sqrt}(c1)$ in F, i.e., $(c2^2) == c1$ in F
3. $c3 = \text{sqrt}(-c1)$ in F, i.e., $(c3^2) == -c1$ in F
4. $c4 = (q + 7) / 16$ # Integer arithmetic

Procedure:

1. $tv1 = x^{c4}$
2. $tv2 = c1 * tv1$
3. $tv3 = c2 * tv1$
4. $tv4 = c3 * tv1$
5. $e1 = (tv2^2) == x$
6. $e2 = (tv3^2) == x$
7. $tv1 = \text{CMOV}(tv1, tv2, e1)$ # Select tv2 if $(tv2^2) == x$
8. $tv2 = \text{CMOV}(tv4, tv3, e2)$ # Select tv3 if $(tv3^2) == x$
9. $e3 = (tv2^2) == x$
10. $z = \text{CMOV}(tv1, tv2, e3)$ # Select the sqrt from tv1 and tv2
11. return z

I.4. Constant-Time Tonelli-Shanks Algorithm

This algorithm is a constant-time version of the classic Tonelli-Shanks algorithm ([C93], Algorithm 1.5.1) due to Sean Bowe, Jack Grigg, and Eirik Ogilvie-Wigley [jubjub-fq], adapted and optimized by Michael Scott.

This algorithm applies to $\text{GF}(p)$ for any p . Note, however, that the special-purpose algorithms given in the prior sections are faster, when they apply.

```

sqrt_ts_ct(x)

Parameters:
- F, a finite field of characteristic p and order q = p^m.

Input x, an element of F.
Output: z, an element of F such that z^2 == x, if x is square in F.

Constants:
1. c1, the largest integer such that 2^c1 divides q - 1.
2. c2 = (q - 1) / (2^c1)      # Integer arithmetic
3. c3 = (c2 - 1) / 2        # Integer arithmetic
4. c4, a non-square value in F
5. c5 = c4^c2 in F

Procedure:
1. z = x^c3
2. t = z * z
3. t = t * x
4. z = z * x
5. b = t
6. c = c5
7. for i in (c1, c1 - 1, ..., 2):
8.   for j in (1, 2, ..., i - 2):
9.     b = b * b
10.  e = b == 1
11.  zt = z * c
12.  z = CMOV(zt, z, e)
13.  c = c * c
14.  tt = t * c
15.  t = CMOV(tt, t, e)
16.  b = t
17. return z

```

I.5. is_square for $F = \text{GF}(p^2)$

The following `is_square` method applies to any field $F = \text{GF}(p^2)$ with basis $(1, I)$ represented as described in Section 2.1, i.e., an element $x = (x_1, x_2) = x_1 + x_2 * I$.

Other optimizations of this type are possible in other extension fields; see, for example, [AR13] for more information.

```
is_square(x)
```

Parameters:

- F, an extension field of characteristic p and order $q = p^2$ with basis (1, I).

Input: x, an element of F.

Output: True if x is square in F, and False otherwise.

Constants:

1. $c1 = (p - 1) / 2$ # Integer arithmetic

Procedure:

1. $tv1 = x_1^2$

2. $tv2 = I * x_2$

3. $tv2 = tv2^2$

4. $tv1 = tv1 - tv2$

5. $tv1 = tv1^{c1}$

6. $e1 = tv1 \neq -1$ # Note: -1 in F

7. return e1

Appendix J. Suite Test Vectors

This section gives test vectors for each suite defined in [Section 8](#). The test vectors in this section were generated using code that is available from [\[hash2curve-repo\]](#).

Each test vector in this section lists values computed by the appropriate encoding function, with variable names defined as in [Section 3](#). For example, for a suite whose encoding type is random oracle, the test vector gives the value for msg, u, Q0, Q1, and the output point P.

J.1. NIST P-256

J.1.1. P256_XMD:SHA-256_SSWU_RO_

```

suite = P256_XMD:SHA-256_SSWU_RO_
dst   = QUUX-V01-CS02-with-P256_XMD:SHA-256_SSWU_RO_

msg   =
P.x   = 2c15230b26dbc6fc9a37051158c95b79656e17a1a920b11394ca91
      c44247d3e4
P.y   = 8a7a74985cc5c776cdf4b1f19884970453912e9d31528c060be9a
      b5c43e8415
u[0]  = ad5342c66a6dd0ff080df1da0ea1c04b96e0330dd89406465eeba1
      1582515009
u[1]  = 8c0f1d43204bd6f6ea70ae8013070a1518b43873bcd850aafa0a9e
      220e2eea5a
Q0.x  = ab640a12220d3ff283510fff3f4b1953d09fad35795140b1c5d64f3
      13967934d5
Q0.y  = dccb558863804a881d4ffff3455716c836cef230e5209594ddd33d8
      5c565b19b1
Q1.x  = 51cce63c50d972a6e51c61334f0f4875c9ac1cd2d3238412f84e31
      da7d980ef5
Q1.y  = b45d1a36d00ad90e5ec7840a60a4de411917fbe7c82c3949a6e699
      e5a1b66aac

msg   = abc
P.x   = 0bb8b87485551aa43ed54f009230450b492fead5f1cc91658775da
      c4a3388a0f
P.y   = 5c41b3d0731a27a7b14bc0bf0ccded2d8751f83493404c84a88e71
      ffd424212e
u[0]  = afe47f2ea2b10465cc26ac403194dfb68b7f5ee865cda61e9f3e07
      a537220af1
u[1]  = 379a27833b0bfe6f7bdca08e1e83c760bf9a338ab335542704edcd
      69ce9e46e0
Q0.x  = 5219ad0ddef3cc49b714145e91b2f7de6ce0a7a7dc7406c7726c7e
      373c58cb48
Q0.y  = 7950144e52d30acbec7b624c203b1996c99617d0b61c2442354301
      b191d93ecf
Q1.x  = 019b7cb4efcfeaf39f738fe638e31d375ad6837f58a852d032ff60
      c69ee3875f
Q1.y  = 589a62d2b22357fed5449bc38065b760095ebe6aeac84b01156ee4
      252715446e

msg   = abcdef0123456789
P.x   = 65038ac8f2b1def042a5df0b33b1f4eca6bfff7cb0f9c6c15268118
      64e544ed80
P.y   = cad44d40a656e7aff4002a8de287abc8ae0482b5ae825822bb870d
      6df9b56ca3
u[0]  = 0fad9d125a9477d55cf9357105b0eb3a5c4259809bf87180aa01d6
      51f53d312c
u[1]  = b68597377392cd3419d8fcc7d7660948c8403b19ea78bbca4b133c
      9d2196c0fb
Q0.x  = a17bdf2965eb88074bc01157e644ed409dac97cfcf0c61c998ed0f
      a45e79e4a2
Q0.y  = 4f1bc80c70d411a3cc1d67aeae6e726f0f311639fee560c7f5a664
      554e3c9c2e
Q1.x  = 7da48bb67225c1a17d452c983798113f47e438e4202219dd0715f8
      419b274d66
Q1.y  = b765696b2913e36db3016c47edb99e24b1da30e761a8a3215dc0ec
      4d8f96e6f9

```


J.1.2. P256_XMD:SHA-256_SSWU_NU_


```
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x = 5c4bad52f81f39c8e8de1260e9a06d72b8b00a0829a8ea004a610b
      0691bea5d9
P.y = c801e7c0782af1f74f24fc385a8555da0582032a3ce038de637ccd
      cb16f7ef7b
u[0] = 0e1527840b9df2dfbef966678ff167140f2b27c4dccd884c25014d
      ce0e41dfa3
Q.x = 5c4bad52f81f39c8e8de1260e9a06d72b8b00a0829a8ea004a610b
      0691bea5d9
Q.y = c801e7c0782af1f74f24fc385a8555da0582032a3ce038de637ccd
      cb16f7ef7b
```

J.2. NIST P-384

J.2.1. P384_XMD:SHA-384_SSWU_RO_

```
suite = P384_XMD:SHA-384_SSWU_RO_  
dst   = QUUX-V01-CS02-with-P384_XMD:SHA-384_SSWU_RO_  
  
msg   =  
P.x   = eb9fe1b4f4e14e7140803c1d99d0a93cd823d2b024040f9c067a8e  
      ca1f5a2eeac9ad604973527a356f3fa3aeff0e4d83  
P.y   = 0c21708cff382b7f4643c07b105c2eaec2cead93a917d825601e63  
      c8f21f6abd9abc22c93c2bed6f235954b25048bb1a  
u[0]  = 25c8d7dc1acd4ee617766693f7f8829396065d1b447eedb155871f  
      effd9c6653279ac7e5c46edb7010a0e4ff64c9f3b4  
u[1]  = 59428be4ed69131df59a0c6a8e188d2d4ece3f1b2a3a02602962b4  
      7efa4d7905945b1e2cc80b36aa35c99451073521ac  
Q0.x  = e4717e29eef38d862bee4902a7d21b44efb58c464e3e1f0d03894d  
      94de310f8fffc6de86786dd3e15a1541b18d4eb2846  
Q0.y  = 6b95a6e639822312298a47526bb77d9cd7bcf76244c991c8cd7007  
      5e2ee6e8b9a135c4a37e3c0768c7ca871c0ceb53d4  
Q1.x  = 509527cfc0750eedc53147e6d5f78596c8a3b7360e0608e2fab056  
      3a1670d58d8ae107c9f04bcf90e89489ace5650efd  
Q1.y  = 33337b13cb35e173fdea4cb9e8cce915d836ff57803dbbeb7998aa  
      49d17df2ff09b67031773039d09fbd9305a1566bc4  
  
msg   = abc  
P.x   = e02fc1a5f44a7519419dd314e29863f30df55a514da2d655775a81  
      d413003c4d4e7fd59af0826dfaad4200ac6f60abe1  
P.y   = 01f638d04d98677d65bef99aef1a12a70a4cbb9270ec55248c0453  
      0d8bc1f8f90f8a6a859a7c1f1ddccedf8f96d675f6  
u[0]  = 53350214cb6bef0b51abb791b1c4209a2b4c16a0c67e1ab1401017  
      fad774cd3b3f9a8bcdf7f6229dd8dd5a075cb149a0  
u[1]  = c0473083898f63e03f26f14877a2407bd60c75ad491e7d26cbc6cc  
      5ce815654075ec6b6898c7a41d74ceaf720a10c02e  
Q0.x  = fc853b69437aee9a19d5acf96a4ee4c5e04cf7b53406dfaa2afbdd  
      7ad2351b7f554e4bbc6f5db4177d4d44f933a8f6ee  
Q0.y  = 7e042547e01834c9043b10f3a8221c4a879cb156f04f72bfccab0c  
      047a304e30f2aa8b2e260d34c4592c0c33dd0c6482  
Q1.x  = 57912293709b3556b43a2dfb137a315d256d573b82ded120ef8c78  
      2d607c05d930d958e50cb6dc1cc480b9afc38c45f1  
Q1.y  = de9387dab0eef0bda219c6f168a92645a84665c4f2137c14270fb4  
      24b7532ff84843c3da383ceea24c47fa343c227bb8  
  
msg   = abcdef0123456789  
P.x   = bdecc1c1d870624965f19505be50459d363c71a699a496ab672f9a  
      5d6b78676400926fbceee6fcd1780fe86e62b2aa89  
P.y   = 57cf1f99b5ee00f3c201139b3bfe4dd30a653193778d89a0acc5e  
      0f47e46e4e4b85a0595da29c9494c1814acafe183c  
u[0]  = aab7fb87238cf6b2ab56cdcca7e028959bb2ea599d34f68484139d  
      de85ec6548a6e48771d17956421bdb7790598ea52e  
u[1]  = 26e8d833552d7844d167833ca5a87c35bcfaa5a0d86023479fb28e  
      5cd6075c18b168bf1f5d2a0ea146d057971336d8d1  
Q0.x  = 0ceece45b73f89844671df962ad2932122e878ad2259e650626924  
      e4e7f132589341dec1480ebcbbbe3509d11fb570b7  
Q0.y  = fafd71a3115298f6be4ae5c6dfc96c400cfb55760f185b7b03f3fa  
      45f3f91eb65d27628b3c705cafd0466fafa54883ce  
Q1.x  = dea1be8d3f9be4cbf4fab9d71d549dde76875b5d9b876832313a08  
      3ec81e528cbc2a0a1d0596b3bcb0ba77866b129776  
Q1.y  = eb15fe71662214fb03b65541f40d3eb0f4cf5c3b559f647da138c9  
      f9b7484c48a08760e02c16f1992762cb7298fa52cf
```


J.2.2. P384_XMD:SHA-384_SSWU_NU_

```
suite = P384_XMD:SHA-384_SSWU_NU_
dst   = QUUX-V01-CS02-with-P384_XMD:SHA-384_SSWU_NU_

msg   =
P.x   = de5a893c83061b2d7ce6a0d8b049f0326f2ada4b966dc7e7292725
      6b033ef61058029a3bfb13c1c7eeced6641881ae20
P.y   = 63f46da6139785674da315c1947e06e9a0867f5608cf24724eb379
      3a1f5b3809ee28eb21a0c64be3be169afc6cdb38ca
u[0]  = bc7dc1b2cdc5d588a66de3276b0f24310d4aca4977efda7d6272e1
      be25187b001493d267dc53b56183c9e28282368e60
Q.x   = de5a893c83061b2d7ce6a0d8b049f0326f2ada4b966dc7e7292725
      6b033ef61058029a3bfb13c1c7eeced6641881ae20
Q.y   = 63f46da6139785674da315c1947e06e9a0867f5608cf24724eb379
      3a1f5b3809ee28eb21a0c64be3be169afc6cdb38ca

msg   = abc
P.x   = 1f08108b87e703c86c872ab3eb198a19f2b708237ac4be53d7929f
      b4bd5194583f40d052f32df66afe5249c9915d139b
P.y   = 1369dc8d5bf038032336b989994874a2270adadb67a7fcc32f0f88
      24bc5118613f0ac8de04a1041d90ff8a5ad555f96c
u[0]  = 9de6cf41e6e41c03e4a7784ac5c885b4d1e49d6de390b3cdd5a1ac
      5dd8c40afb3dfd7bb2686923bab644134483fc1926
Q.x   = 1f08108b87e703c86c872ab3eb198a19f2b708237ac4be53d7929f
      b4bd5194583f40d052f32df66afe5249c9915d139b
Q.y   = 1369dc8d5bf038032336b989994874a2270adadb67a7fcc32f0f88
      24bc5118613f0ac8de04a1041d90ff8a5ad555f96c

msg   = abcdef0123456789
P.x   = 4dac31ec8a82ee3c02ba2d7c9fa431f1e59ffe65bf977b948c59e1
      d813c2d7963c7be81aa6db39e78ff315a10115c0d0
P.y   = 845333cdb5702ad5c525e603f302904d6fc84879f0ef2ee2014a6b
      13edd39131bfd66f7bd7cdc2d9ccf778f0c8892c3f
u[0]  = 84e2d430a5e2543573e58e368af41821ca3ccc97baba7e9aab51a8
      4543d5a0298638a22ceee6090d9d642921112af5b7
Q.x   = 4dac31ec8a82ee3c02ba2d7c9fa431f1e59ffe65bf977b948c59e1
      d813c2d7963c7be81aa6db39e78ff315a10115c0d0
Q.y   = 845333cdb5702ad5c525e603f302904d6fc84879f0ef2ee2014a6b
      13edd39131bfd66f7bd7cdc2d9ccf778f0c8892c3f

msg   = q128_qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
      qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
      qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
P.x   = 13c1f8c52a492183f7c28e379b0475486718a7e3ac1dfef39283b9
      ce5fb02b73f70c6c1f3dfe0c286b03e2af1af12d1d
P.y   = 57e101887e73e40eab8963324ed16c177d55eb89f804ec9df06801
      579820420b5546b579008df2145fd770f584a1a54c
u[0]  = 504e4d5a529333b9205acaa283107bd1bffd5753898f7744161f7d
      d19ba57fbb6a64214a2e00ddd2613d76cd508ddb30
Q.x   = 13c1f8c52a492183f7c28e379b0475486718a7e3ac1dfef39283b9
      ce5fb02b73f70c6c1f3dfe0c286b03e2af1af12d1d
Q.y   = 57e101887e73e40eab8963324ed16c177d55eb89f804ec9df06801
      579820420b5546b579008df2145fd770f584a1a54c

msg   = a512_aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
```

```
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x = af129727a4207a8cb9e9dce656d88f79fce25edbcea350499d65e9
      bf1204537bdde73c7cefb752a6ed5ebcd44e183302
P.y = ce68a3d5e161b2e6a968e4ddaa9e51504ad1516ec170c7eef3ca6b
      5327943eca95d90b23b009ba45f58b72906f2a99e2
u[0] = 7b01ce9b8c5a60d9fbc202d6dde92822e46915d8c17e03fcb92ece
      1ed6074d01e149fc9236def40d673de903c1d4c166
Q.x = af129727a4207a8cb9e9dce656d88f79fce25edbcea350499d65e9
      bf1204537bdde73c7cefb752a6ed5ebcd44e183302
Q.y = ce68a3d5e161b2e6a968e4ddaa9e51504ad1516ec170c7eef3ca6b
      5327943eca95d90b23b009ba45f58b72906f2a99e2
```

J.3. NIST P-521

J.3.1. P521_XMD:SHA-512_SSWU_RO_

```

suite = P521_XMD:SHA-512_SSWU_RO_
dst   = QUUX-V01-CS02-with-P521_XMD:SHA-512_SSWU_RO_

msg   =
P.x   = 00fd767cebb2452030358d0e9cf907f525f50920c8f607889a6a35
      680727f64f4d66b161fafeb2654bea0d35086bec0a10b30b14adf
      3556ed9f7f1bc23cecc9c088
P.y   = 0169ba78d8d851e930680322596e39c78f4fe31b97e57629ef6460
      ddd68f8763fd7bd767a4e94a80d3d21a3c2ee98347e024fc73ee1c
      27166dc3fe5eeef782be411d
u[0]  = 01e5f09974e5724f25286763f00ce76238c7a6e03dc396600350ee
      2c4135fb17dc555be99a4a4bae0fd303d4f66d984ed7b6a3ba3860
      93752a855d26d559d69e7e9e
u[1]  = 00ae593b42ca2ef93ac488e9e09a5fe5a2f6fb330d18913734ff60
      2f2a761fcaaf5f596e790bcc572c9140ec03f6cccc38f767f1c197
      5a0b4d70b392d95a0c7278aa
Q0.x  = 00b70ae99b6339fffac19cb9bfde2098b84f75e50ac1e80d6acb95
      4e4534af5f0e9c4a5b8a9c10317b8e6421574bae2b133b4f2b8c6c
      e4b3063da1d91d34fa2b3a3c
Q0.y  = 007f368d98a4ddb381fb354de40e44b19e43bb11a1278759f4ea7
      b485e1b6db33e750507c071250e3e443c1aaed61f2c28541bb54b1
      b456843eda1eb15ec2a9b36e
Q1.x  = 01143d0e9cddcdac6a9aafe1bcf8d218c0afc45d4451239e821f5
      d2a56df92be942660b532b2aa59a9c635ae6b30e803c45a6ac8714
      32452e685d661cd41cf67214
Q1.y  = 00ff75515df265e996d702a5380defffab1a6d2bc232234c7bcffa
      433cd8aa791fbc8dcf667f08818bffa739ae25773b32073213cae9
      a0f2a917a0b1301a242dda0c

msg   = abc
P.x   = 002f89a1677b28054b50d15e1f81ed6669b5a2158211118ebdef8a
      6efc77f8ccaa528f698214e4340155abc1fa08f8f613ef14a04371
      7503d57e267d57155cf784a4
P.y   = 010e0be5dc8e753da8ce51091908b72396d3deed14ae166f66d8eb
      f0a4e7059ead169ea4bead0232e9b700dd380b316e9361cfdba55a
      08c73545563a80966ecbb86d
u[0]  = 003d00c37e95f19f358adeeaa47288ec39998039c3256e13c2a4c0
      0a7cb61a34c8969472960150a27276f2390eb5e53e47ab193351c2
      d2d9f164a85c6a5696d94fe8
u[1]  = 01f3cbd3df3893a45a2f1fecdac4d525eb16f345b03e2820d69bc5
      80f5cbe9cb89196fdf720ef933c4c0361fcfe29940fd0db0a5da6b
      afb0bee8876b589c41365f15
Q0.x  = 01b254e1c99c835836f0aceebba7d77750c48366ecb07fb658e4f5
      b76e229ae6ca5d271bb0006ffcc42324e15a6d3daae587f9049de2
      dbb0494378ffb60279406f56
Q0.y  = 01845f4af72fc2b1a5a2fe966f6a97298614288b456cfc385a425b
      686048b25c952fbb5674057e1eb055d04568c0679a8e2dda3158dc
      16ac598dbb1d006f5ad915b0
Q1.x  = 007f08e813c620e527c961b717ffc74aac7afccb9158cebc347d57
      15d5c2214f952c97e194f11d114d80d3481ed766ac0a3dba3eb73f
      6ff9ccb9304ad10bbd7b4a36
Q1.y  = 0022468f92041f9970a7cc025d71d5b647f822784d29ca7b3bc3b0
      829d6bb8581e745f8d0cc9dc6279d0450e779ac2275c4c3608064a
      d6779108a7828ebd9954caeb

msg   = abcdef0123456789
P.x   = 006e200e276a4a81760099677814d7f8794a4a5f3658442de63c18

```

```
d2244dcc957c645e94cb0754f95fcf103b2aeaf94411847c24187b
89fb7462ad3679066337cbc4
P.y = 001dd8dfa9775b60b1614f6f169089d8140d4b3e4012949b52f98d
b2deff3e1d97bf73a1fa4d437d1dcdcf39b6360cc518d8ebcc0f899
018206fded7617b654f6b168
u[0] = 00183ee1a9bbdc37181b09ec336bcaa34095f91ef14b66b1485c16
6720523dfb81d5c470d44afcb52a87b704dbc5c9bc9d0ef524dec2
9884a4795f55c1359945baf3
u[1] = 00504064fd137f06c81a7cf0f84aa7e92b6b3d56c2368f0a08f447
76aa8930480da1582d01d7f52df31dca35ee0a7876500ece3d8fe0
293cd285f790c9881c998d5e
Q0.x = 0021482e8622aac14da60e656043f79a6a110cbae5012268a62dd6
a152c41594549f373910ebed170ade892dd5a19f5d687fae7095a4
61d583f8c4295f7aaf8cd7da
Q0.y = 0177e2d8c6356b7de06e0b5712d8387d529b848748e54a8bc0ef5f
1475aa569f8f492fa85c3ad1c5edc51faf7911f11359bfa2a12d2e
f0bd73df9cb5abd1b101c8b1
Q1.x = 00abeafb16fdbb5eb95095678d5a65c1f293291dfd20a3751dbe05
d0a9bfe2d2eef19449fe59ec32cdd4a4adc3411177c0f2dff0159
438706159a1bbd0567d9b3d0
Q1.y = 007cc657f847db9db651d91c801741060d63dab4056d0a1d3524e2
eb0e819954d8f677aa353bd056244a88f00017e00c3ce8beedb43
82d83d74418bd48930c6c182

msg = q128_qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
P.x = 01b264a630bd6555be537b000b99a06761a9325c53322b65bdc41b
f196711f9708d58d34b3b90faf12640c27b91c70a507998e559406
48caa8e71098bf2bc8d24664
P.y = 01ea9f445bee198b3ee4c812dcf7b0f91e0881f0251aab272a1220
1fd89b1a95733fd2a699c162b639e9acdcc54fdc2f6536129b6beb
0432be01aa8da02df5e59aaa
u[0] = 0159871e222689aad7694dc4c3480a49807b1eedd9c8cb4ae1b219
d5ba51655ea5b38e2e4f56b36bf3e3da44a7b139849d28f598c816
fe1bc7ed15893b22f63363c3
u[1] = 004ef0cffd475152f3858c0a8ccbdf7902d8261da92744e98df9b7
fadb0a5502f29c5086e76e2cf498f47321434a40b1504911552ce4
4ad7356a04e08729ad9411f5
Q0.x = 0005eac7b0b81e38727efcab1e375f6779aea949c3e409b53a1d37
aa2acb87a7e6ad24aafb3c52f82f7f0e21b872e88c55e17b7fa
21ce08a94ea2121c42c2eb73
Q0.y = 00a173b6a53a7420dbd61d4a21a7c0a52de7a5c6ce05f31403bef7
47d16cc8604a039a73bdd6e114340e55dacd6bea8e217ffbadfb8c
292afa3e1b2afc839a6ce7bb
Q1.x = 01881e3c193a69e4d88d8180a6879b74782a0bc7e529233e9f84bf
7f17d2f319c36920ffba26f9e57a1e045cc7822c834c239593b6e1
42a694aa00c757b0db79e5e8
Q1.y = 01558b16d396d866e476e001f2dd0758927655450b84e12f154032
c7c2a6db837942cd9f44b814f79b4d729996ced61eec61d85c6751
39cbffe3fbf071d2c21cfecb

msg = a512_aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
```

```
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P . x   = 00c12bc3e28db07b6b4d2a2b1167ab9e26fc2fa85c7b0498a17b03
47edf52392856d7e28b8fa7a2dd004611159505835b687ecf1a764
857e27e9745848c436ef3925
P . y   = 01cd287df9a50c22a9231beb452346720bb163344a41c5f5a24e83
35b6ccc595fd436aea89737b1281aecb411eb835f0b939073fdd1d
d4d5a2492e91ef4a3c55bcbd
u[0]    = 0033d06d17bc3b9a3efc081a05d65805a14a3050a0dd4dfb488461
8eb5c73980a59c5a246b18f58ad022dd3630faa22889fbb8ba1593
466515e6ab4aeb7381c26334
u[1]    = 0092290ab99c3fea1a5b8fb2ca49f859994a04faee3301cefab312
d34227f6a2d0c3322cf76861c6a3683bdaa2dd2a6daa5d6906c663
e065338b2344d20e313f1114
Q0 . x  = 00041f6eb92af8777260718e4c22328a7d74203350c6c8f5794d99
d5789766698f459b83d5068276716f01429934e40af3d1111a2278
0b1e07e72238d2207e5386be
Q0 . y  = 001c712f0182813942b87cab8e72337db017126f52ed797dd23458
4ac9ae7e80dfe7abea11db02cf1855312eae1447dbaecc9d7e8c88
0a5e76a39f6258074e1bc2e0
Q1 . x  = 0125c0b69bcf55eab49280b14f707883405028e05c927cd7625d4e
04115bd0e0e6323b12f5d43d0d6d2eff16dbc244542f84ec05891
1260dc3bb6512ab5db285fbd
Q1 . y  = 008bddfb803b3f4c761458eb5f8a0aee3e1f7f68e9d7424405fa69
172919899317fb6ac1d6903a432d967d14e0f80af63e7035aaae0c
123e56862ce969456f99f102
```


J.3.2. P521_XMD:SHA-512_SSWU_NU_

```
suite = P521_XMD:SHA-512_SSWU_NU_  
dst   = QUUX-V01-CS02-with-P521_XMD:SHA-512_SSWU_NU_  
  
msg   =  
P.x   = 01ec604b4e1e3e4c7449b7a41e366e876655538acf51fd40d08b97  
      be066f7d020634e906b1b6942f9174b417027c953d75fb6ec64b8c  
      ee2a3672d4f1987d13974705  
P.y   = 00944fc439b4aad2463e5c9cfa0b0707af3c9a42e37c5a57bb4ecd  
      12fef9fb21508568aedcdd8d2490472df4bbafd79081c81e99f4da  
      3286eddf19be47e9c4cf0e91  
u[0]  = 01e4947fe62a4e47792cee2798912f672fff820b2556282d9843b4  
      b465940d7683a986f93ccb0e9a191fbc09a6e770a564490d2a4ae5  
      1b287ca39f69c3d910ba6a4f  
Q.x   = 01ec604b4e1e3e4c7449b7a41e366e876655538acf51fd40d08b97  
      be066f7d020634e906b1b6942f9174b417027c953d75fb6ec64b8c  
      ee2a3672d4f1987d13974705  
Q.y   = 00944fc439b4aad2463e5c9cfa0b0707af3c9a42e37c5a57bb4ecd  
      12fef9fb21508568aedcdd8d2490472df4bbafd79081c81e99f4da  
      3286eddf19be47e9c4cf0e91  
  
msg   = abc  
P.x   = 00c720ab56aa5a7a4c07a7732a0a4e1b909e32d063ae1b58db5f0e  
      b5e09f08a9884bff55a2bef4668f715788e692c18c1915cd034a6b  
      998311fcf46924ce66a2be9a  
P.y   = 003570e87f91a4f3c7a56be2cb2a078ffc153862a53d5e03e5dad5  
      bccc6c529b8bab0b7dbb157499e1949e4edab21cf5d10b782bc1e9  
      45e13d7421ad8121dbc72b1d  
u[0]  = 0019b85ef78596efc84783d42799e80d787591fe7432dee1d9fa2b  
      7651891321be732ddf653fa8fefaf34d86fb728db569d36b5b6ed39  
      83945854b2fc2dc6a75aa25b  
Q.x   = 00c720ab56aa5a7a4c07a7732a0a4e1b909e32d063ae1b58db5f0e  
      b5e09f08a9884bff55a2bef4668f715788e692c18c1915cd034a6b  
      998311fcf46924ce66a2be9a  
Q.y   = 003570e87f91a4f3c7a56be2cb2a078ffc153862a53d5e03e5dad5  
      bccc6c529b8bab0b7dbb157499e1949e4edab21cf5d10b782bc1e9  
      45e13d7421ad8121dbc72b1d  
  
msg   = abcdef0123456789  
P.x   = 00bcacf32a968ff7971b3bbd9ce8edfbee1309e2019d7ff373c3838  
      7a782b005dce6ceffccfed5c6511c8f7f312f343f3a891029c585  
      8f45ee0bf370aba25fc990cc  
P.y   = 00923517e767532d82cb8a0b59705eec2b7779ce05f9181c7d5d5e  
      25694ef8ebd4696343f0bc27006834d2517215ecf79482a84111f5  
      0c1bae25044fe1dd77744bbd  
u[0]  = 01dba0d7fa26a562ee8a9014ebc2cca4d66fd9de036176aca8fc11  
      ef254cd1bc208847ab7701dbca7af328b3f601b11a1737a899575a  
      5c14f4dca5aaca45e9935e07  
Q.x   = 00bcacf32a968ff7971b3bbd9ce8edfbee1309e2019d7ff373c3838  
      7a782b005dce6ceffccfed5c6511c8f7f312f343f3a891029c585  
      8f45ee0bf370aba25fc990cc  
Q.y   = 00923517e767532d82cb8a0b59705eec2b7779ce05f9181c7d5d5e  
      25694ef8ebd4696343f0bc27006834d2517215ecf79482a84111f5  
      0c1bae25044fe1dd77744bbd  
  
msg   = q128_qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq  
      qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq  
      qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq  
      qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
```


J.4. `curve25519`

J.4.1. `curve25519_XMD:SHA-512_ELL2_RO_`

```
suite = curve25519_XMD:SHA-512_ELL2_RO_
dst   = QUUX-V01-CS02-with-curve25519_XMD:SHA-512_ELL2_RO_

msg   =
P.x   = 2de3780abb67e861289f5749d16d3e217ffa722192d16bbd9d1bfb
      9d112b98c0
P.y   = 3b5dc2a498941a1033d176567d457845637554a2fe7a3507d21abd
      1c1bd6e878
u[0]  = 005fe8a7b8fef0a16c105e6cadf5a6740b3365e18692a9c05bfb4
      d97f645a6a
u[1]  = 1347edbec6a2b5d8c02e058819819bee177077c9d10a4ce165aab0
      fd0252261a
Q0.x  = 36b4df0c864c64707cbf6cf36e9ee2c09a6cb93b28313c169be295
      61bb904f98
Q0.y  = 6cd59d664fb58c66c892883cd0eb792e52055284dac3907dd756b4
      5d15c3983d
Q1.x  = 3fa114783a505c0b2b2fbee0102853c0b494e7757f2a089d0daae
      7ed9a0db2b
Q1.y  = 76c0fe7fec932aaafb8ee0fb42d9cbb32eb931158f469ff3050af15
      cfdbbef94

msg   = abc
P.x   = 2b4419f1f2d48f5872de692b0aca72cc7b0a60915dd70bde432e82
      6b6abc526d
P.y   = 1b8235f255a268f0a6fa8763e97eb3d22d149343d495da1160eff9
      703f2d07dd
u[0]  = 49bed021c7a3748f09fa8cdfcac044089f7829d3531066ac9e74e0
      994e05bc7d
u[1]  = 5c36525b663e63389d886105cee7ed712325d5a97e60e140aba7e2
      ce5ae851b6
Q0.x  = 16b3d86e056b7970fa00165f6f48d90b619ad618791661b7b5e1ec
      78be10eac1
Q0.y  = 4ab256422d84c5120b278cbdfc4e1facc5baadffeccecf8ee9bf39
      46106d50ca
Q1.x  = 7ec29ddb34539c40adfa98fcb39ec36368f47f30e8f888cc7e86f
      4d46e0c264
Q1.y  = 10d1abc1cae2d34c06e247f2141ba897657fb39f1080d54f09ce0a
      f128067c74

msg   = abcdef0123456789
P.x   = 68ca1ea5a6acf4e9956daa101709b1eee6c1bb0df1de3b90d46023
      82a104c036
P.y   = 2a375b656207123d10766e68b938b1812a4a6625ff83cb8d5e86f5
      8a4be08353
u[0]  = 6412b7485ba26d3d1b6c290a8e1435b2959f03721874939b21782d
      f17323d160
u[1]  = 24c7b46c1c6d9a21d32f5707be1380ab82db1054fde82865d5c9e3
      d968f287b2
Q0.x  = 71de3dadfe268872326c35ac512164850860567aea0e7325e6b91a
      98f86533ad
Q0.y  = 26a08b6e9a18084c56f2147bf515414b9b63f1522e1b6c5649f7d4
      b0324296ec
Q1.x  = 5704069021f61e41779e2ba6b932268316d6d2a6f064f997a22fef
      16d1eaeaca
Q1.y  = 50483c7540f64fb4497619c050f2c7fe55454ec0f0e79870bb4430
      2e34232210
```


J.4.2. `curve25519_XMD:SHA-512_ELL2_NU_`


```
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x = 5fd892c0958d1a75f54c3182a18d286efab784e774d1e017ba2fb2
52998b5dc1
P.y = 750af3c66101737423a4519ac792fb93337bd74ee751f19da4cf1e
94f4d6d0b8
u[0] = 1a68a1af9f663592291af987203393f707305c7bac9c8d63d6a729
bdc553dc19
Q.x = 3bcd651ee54d5f7b6013898aab251ee8ecc0688166fce6e9548d38
472f6bd196
Q.y = 1bb36ad9197299f111b4ef21271c41f4b7ecf5543db8bb5931307e
bdb2eaa465
```

J.5. edwards25519

J.5.1. edwards25519_XMD:SHA-512_ELL2_RO_

```

suite = edwards25519_XMD:SHA-512_ELL2_R0_
dst   = QUUX-V01-CS02-with-edwards25519_XMD:SHA-512_ELL2_R0_

msg   =
P.x   = 3c3da6925a3c3c268448dcabb47ccde5439559d9599646a8260e47
      b1e4822fc6
P.y   = 09a6c8561a0b22bef63124c588ce4c62ea83a3c899763af26d7953
      02e115dc21
u[0]  = 03fef4813c8cb5f98c6eef88fae174e6e7d5380de2b007799ac7ee
      712d203f3a
u[1]  = 780bddd137290c8f589dc687795aafae35f6b674668d92bf92ae7
      93e6a60c75
Q0.x  = 6549118f65bb617b9e8b438decedc73c496eaed496806d3b2eb9ee
      60b88e09a7
Q0.y  = 7315bcc8cf47ed68048d22bad602c6680b3382a08c7c5d3f439a97
      3fb4cf9feb
Q1.x  = 31dcfc5c58aa1bee6e760bf78cbe71c2bead8cebb2e397ece0f37a
      3da19c9ed2
Q1.y  = 7876d81474828d8a5928b50c82420b2bd0898d819e9550c5c82c39
      fc9bafa196

msg   = abc
P.x   = 608040b42285cc0d72cbb3985c6b04c935370c7361f4b7fbdb1ae7
      f8c1a8ecad
P.y   = 1a8395b88338f22e435bbd301183e7f20a5f9de643f11882fb237f
      88268a5531
u[0]  = 5081955c4141e4e7d02ec0e36becffaa1934df4d7a270f70679c78
      f9bd57c227
u[1]  = 005bdc17a9b378b6272573a31b04361f21c371b256252ae5463119
      aa0b925b76
Q0.x  = 5c1525bd5d4b4e034512949d187c39d48e8cd84242aa4758956e4a
      dc7d445573
Q0.y  = 2bf426cf7122d1a90abc7f2d108befc2ef415ce8c2d09695a74072
      40faa01f29
Q1.x  = 37b03bba828860c6b459ddad476c83e0f9285787a269df2156219b
      7e5c86210c
Q1.y  = 285ebf5412f84d0ad7bb4e136729a9ffd2195d5b8e73c0dc85110c
      e06958f432

msg   = abcdef0123456789
P.x   = 6d7fabf47a2dc03fe7d47f7dddd21082c5fb8f86743cd020f3fb14
      7d57161472
P.y   = 53060a3d140e7fbcda641ed3cf42c88a75411e648a1add71217f70
      ea8ec561a6
u[0]  = 285ebaa3be701b79871bcb6e225ecc9b0b32dff2d60424b4c50642
      636a78d5b3
u[1]  = 2e253e6a0ef658fedb8e4bd6a62d1544fd6547922acb3598ec6b36
      9760b81b31
Q0.x  = 3ac463dd7fddb773b069c5b2b01c0f6b340638f54ee3bd92d452fc
      ec3015b52d
Q0.y  = 7b03ba1e8db9ec0b390d5c90168a6a0b7107156c994c674b61fe69
      6cbeb46baf
Q1.x  = 0757e7e904f5e86d2d2f4acf7e01c63827fde2d363985aa7432106
      f1b3a444ec
Q1.y  = 50026c96930a24961e9d86aa91ea1465398ff8e42015e2ec1fa397
      d416f6a1c0

```


J.5.2. edwards25519_XMD:SHA-512_ELL2_NU_


```
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x = 6e5e1f37e99345887fc12111575fc1c3e36df4b289b8759d23af14
      d774b66bfff
P.y = 2c90c3d39eb18ff291d33441b35f3262cdd307162cc97c31bfcc7a
      4245891a37
u[0] = 3cb0178a8137cefa5b79a3a57c858d7eeea787b2781be4a362a2f
      0750d24fa0
Q.x = 3e6368cff6e88a58e250c54bd27d2c989ae9b3acb6067f2651ad28
      2ab8c21cd9
Q.y = 38fb39f1566ca118ae6c7af42810c0bb9767ae5960abb5a8ca7925
      30fb9447d
```

J.6. curve448

J.6.1. curve448_XOF:SHAKE256_ELL2_RO_


```
a55ac204e153f77285d780c4acee6c96abe3577a0c0b00be6e790c
f194
P.y = 935247a64bf78c107069943c7e3ecc52acb27ce4a3230407c83573
41685ea2152e8c3da93f8cd77da1bddd5bb759c6e7ae7d516dced4
2850
u[0] = 6aab71a38391639f27e49eae8b1cb6b7172a1f478190ece293957e
7cdb2391e7cc1c4261970d9c1bbf9c3915438f74fbd7eb5cd4d4d1
7ace
u[1] = c80b8380ca47a3bcbf76caa75cef0e09f3d270d5ee8f676cde11ae
df41aaca6741bd81a86232bd336ccb42efad39f06542bc06a67b65
909e
Q0.x = 946d91bd50c90ef70743e0dd194bddd68bb630f4e67e5b93e15a9b
94e62cb85134467993501759525c1f4fdbf06f10ddaf817847d735
e062
Q0.y = 185cf511262ec1e9b3c3cbdc015ab93df4e71cbe87766917d81c9f
3419d480407c1462385122c84982d4dae60c3ae4acce0089e37ad6
5934
Q1.x = 01778f4797b717cd6f83c193b2dfb92a1606a36ede941b0f6ab0ac
71ad0eac756d17604bf054398887da907e41065d3595f178ae802f
2087
Q1.y = b4ca727d0bda895e0eee7eb3cbc28710fa2e90a73b568cae26bd7c
2e73b70a9fa0affe1096f0810198890ed65d8935886b6e60dc4c56
9dc6

msg = q128_qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
P.x = 5bd67c4f88adf6beb10f7e0d0054659776a55c97b809ec8b310172
9e104fd0f684e103792f267fd87cc4afc25a073956ef4f268fb028
24d5
P.y = da1f5cb16a352719e4cb064cf47ba72aeba7752d03e8ca2c56229f
419b4ef378785a5af1a53dd7ab4d467c1f92f7b139b3752faf29c9
6432
u[0] = cb5c27e51f9c18ee8ffdb6be230f4eb4f2c2481963b2293484f08d
a2241c1ff59f80978e6defe9d70e34abba2fcbe12dc3a1eb2c5d3d
2e4a
u[1] = c895e8afeccec5466e126fa70fc4aa784b8009063afb10e3ee06a9b
22318256aa8693b0c85b955cf2d6540b8ed71e729af1b8d5ca3b11
6cd7
Q0.x = c2d275826d6ad55e41a22318f6b6240f1f862a2e231120ff41eadb
ec319756032e8cef2a7ac6c10214fa0608c17fc61ec2694a8a2b
358b
Q0.y = 93d2e092762b135509840e609d413200df800d99da91d8b8284066
6cac30e7a3520adbaa4b089bfdc86132e42729f651d022f4782502
f12c
Q1.x = 3c0880ece7244036e9a45944a85599f9809d772f770cc237ac41b2
1aa71615e4f3bb08f64fca618896e4f6cf5bd92e16b89d2cf6e195
6bfb
Q1.y = 45cce4beb96505cac5976b3d2673641e9bcd18d3462bbb453d293e
5282740a6389cfeae610adc7bd425c728541ceec83fcc999164af4
3fb5

msg = a512_aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
```

```
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x = ea441c10b3636ecedd5c0dfcae96384cc40de8390a0ab648765b45
08da12c586d55dc981275776507ebca0e4d1bcaa302bb69dcfa31b
3451
P.y = fee0192d49bcc0c28d954763c2cbe739b9265c4bebe3883803c649
71220cfda60b9ac99ad986cd908c0534b260b5cfca46f6c2b0f3f2
1bda
u[0] = 8cba93a007bb2c801b1769e026b1fa1640b14a34cf3029db3c7fd6
392745d6fec0f7870b5071d6da4402cedbbde28ae4e50ab30e1049
a238
u[1] = 4223746145069e4b8a981acc3404259d1a2c3ecfed5d864798a89d
45f81a2c59e2d40eb1d5f0fe11478cbb2bb30246dd388cb932ad7b
b330
Q0.x = 4321ab02a9849128691e9b80a5c5576793a218de14885fddccb91f
17ceb1646ea00a28b69ad211e1f14f17739612dbde3782319bdf00
9689
Q0.y = 1b8a7b539519eec0ea9f7a46a43822e16cba39a439733d6847ac44
a806b8adb3e1a75ea48a1228b8937ba85c6cb6ee01046e10cad895
3b1e
Q1.x = 126d744da6a14fddec0f78a9cee4571c1320ac7645b600187812e4
d7021f98fc4703732c54daec787206e1f34d9dbbf4b292c68160b8
bfbf
Q1.y = 136eebe6020f2389d448923899a1a38a4c8ad74254e0686e91c4f9
3c1f8f8e1bd619ffb7c1281467882a9c957d22d50f65c5b72b2aee
11af
```

J.6.2. curve448_XOF:SHAKE256_ELL2_NU_

```
suite = curve448_XOF:SHAKE256_ELL2_NU_
dst   = QUUX-V01-CS02-with-curve448_XOF:SHAKE256_ELL2_NU_

msg   =
P.x   = b65e8dbb279fd656f926f68d463b13ca7a982b32f5da9c7cc58afc
      f6199e4729863fb75ca9ae3c95c6887d95a5102637a1c5c40ff0aa
      fadc
P.y   = ea1ea211cf29eca11c057fe8248181591a19f6ac51d45843a65d4b
      b8b71bc83a64c771ed7686218a278ef1c5d620f3d26b5316218864
      5453
u[0]  = 242c70f74eac8184116c71630d284cf8a742fc463e710545847ff6
      4d8e9161cb9f599728a18a32dbd8b67c3bec5d64c9b1d2f2cde7b5
      888d
Q.x   = e6304424de5af3f556d3e645600530c53ad949891c3e60ba041dd5
      f68a93901befff8440164477d348c13d28e27bfc360c44c80b4c7d
      4cea
Q.y   = 4160a8f2043a347185406a6a7e50973b98b82edbdfa3209b0e1c90
      118e10eeb45045b0990d4b2b0708a30eca17df40ad53c9100f20c1
      0b44

msg   = abc
P.x   = 51aceca4fa95854bbaba58d8a5e17a86c07acadef32e1188cafda2
      6232131800002cc2f27c7aec454e5e0c615bddfffb7df6a5f7f0f14
      793f
P.y   = c590c9246eb28b08dee816d608ef233ea5d76e305dc458774a1e1b
      d880387e6734219e2018e4aa50a49486dce0ba8740065da37e6cf5
      212c
u[0]  = ef6dcb75b696d325fb36d66b104700df1480c4c17ea9190d447eee
      1e7e4c9b7f36bbfb8ba7ba7c4cb6b07fed16531c1ac7a26a3618b4
      0b34
Q.x   = de0dc93df9ce7953452f20e270699c1e7dacd5d571c226d77f53b7
      e3053d16f8a81b1601efb362054e973c8e733b663af93f00cb81ba
      f130
Q.y   = 8c5bdec6fa6690905f6eff966b0f98f5a8161493bd04976684d4ec
      1f4512fa8743d86860b2ff2c5d67e9c145fd906f2cb89ff812c6b9
      883f

msg   = abcdef0123456789
P.x   = c6d65987f146b8d0cb5d2c44e1872ac3af1f458f6a8bd8c232ffe8
      b9d09496229a5a27f350eb7d97305bcc4e0f38328718352e8e3129
      ed71
P.y   = 4d2f901bf333fdc4135b954f20d59207e9f6a4ecf88ce5af11c892
      b44f79766ec4ecc9f60d669b95ca8940f39b1b7044140ac2040c1b
      f659
u[0]  = 3012ba5d9b3bb648e4613833a26ecaeadb3e8c8bba07fc90ac3da0
      375769289c44d3dc87474b23df7f45f9a4030892cda689e343aeae
      a6ad
Q.x   = dc29532761f03c24d57f530da4c24acc4c676d185becaa89fcc083
      266541fb7f10ecec91dac64a34cd988274633ae25c4d784aee52de
      47a8
Q.y   = a5f6da11259c69f2e07fce6a7b6afec4c25bd2df83426765f9c070
      4111da24c6a0550d5c7aac7d648d55f7640d50be99c926195e852a
      daac

msg   = q128_qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
      qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
      qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
      qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
```


J.7. **edwards448**

J.7.1. edwards448_XOF:SHAKE256_ELL2_RO_

```

suite = edwards448_XOF:SHAKE256_ELL2_RO_
dst   = QUUX-V01-CS02-with-edwards448_XOF:SHAKE256_ELL2_RO_

msg   =
P.x   = 73036d4a88949c032f01507005c133884e2f0d81f9a950826245dd
      a9e844fc78186c39daaa7147ead3e462cff60e9c6340b58134480b
      4d17
P.y   = 94c1d61b43728e5d784ef4fcb1f38e1075f3aef5e99866911de5a2
      34f1aafdc26b554344742e6ba0420b71b298671bb2b773661863
      4610
u[0]  = 0847c5ebf957d3370b1f98fde499fb3e659996d9fc9b5707176ade
      785ba72cd84b8a5597c12b1024be5f510fa5ba99642c4cec7f3f69
      d3e7
u[1]  = f8cbd8a7ae8c8deed071f3ac4b93e7cfc8f1eac1645d699fd6d38
      81cb295a5d3006d9449ed7cad412a77a1fe61e84a9e41d59ef384d
      6f9a
Q0.x  = c08177330869db17fb81a5e6e53b36d29086d806269760f2e4caba
      a4015f5dbadb7ca2ba594d96a89d0ca4f0944489e1ef393d53db85
      096f
Q0.y  = 02e894598c050eeb7195f5791f1a5f65da3776b7534be37640bcbf
      95d4b915bd22333c50387583507169708fbd7bea0d7aa385dcc614
      be9c
Q1.x  = 770877fd3b6c5503398157b68a9d3609f585f40e1ebebdd69bb0e4
      d3d9aa811995ce75333fdadfa50db886a35959cc59cffd5c9710da
      ca25
Q1.y  = b27fef77aa6231fbbc27538fa90eaca8abd03eb1e62fdae4ec5e82
      8117c3b8b3ff8c34d0a6e6d79fff16d339b94ae8ede33331d5b464
      c792

msg   = abc
P.x   = 4e0158acacffa545adb818a6ed8e0b870e6abc24dfc1dc45cf9a05
      2e98469275d9ff0c168d6a5ac7ec05b742412ee090581f12aa398f
      9f8c
P.y   = 894d3fa437b2d2e28cdc3bfaade035430f350ec5239b6b406b5501
      da6f6d6210ff26719cad83b63e97ab26a12df6dec851d6bf38e294
      af9a
u[0]  = 04d975cd938ab49be3e81703d6a57cca84ed80d2ff6d4756d3f229
      47fb5b70ab0231f0087cbfb4b7cae73b41b0c9396b356a4831d9a1
      4322
u[1]  = 2547ca887ac3db7b5fad3a098aa476e90078afe1358af6c63d677d
      6edfd2100bc004e0f5db94dd2560fc5b308e223241d00488c9ca6b
      0ef2
Q0.x  = 7544612a97f4419c94ab0f621a1ee8ccf46c6657b8e0778ec9718b
      f4b41bc774487ad87d9b1e617aa49d3a4dd35a3cf57cd390ebf042
      9952
Q0.y  = d3ab703e60267d796b485bb58a28f934bd0133a6d1bbdfeda5277f
      a293310be262d7f653a5adffa608c37ed45c0e6008e54a16e1a342
      e4df
Q1.x  = 6262f18d064bc131ade1b8bbcf1cbdf984f4f88153fcc9f94c888a
      f35d5e41aae84c12f169a55d8abf06e6de6c5b23079e587a58cf73
      303e
Q1.y  = 6d57589e901abe7d947c93ab02c307ad9093ed9a83eb0b6e829fb7
      318d590381ca25f3cc628a36a924a9ddf94edf3b338ea7
      7403

msg   = abcdef0123456789
P.x   = 2c25b4503fad94b27391933b557abdecc601c13ed51c5de683894

```

```
84f93dbd6c22e5f962d9babf7a39f39f994312f8ca23344847e1fb
P.y = f176
      = d5e6f5350f430e53a110f5ac7fcc82a96cb865aeca982029522d32
u[0] = 601e41c042a9dfbdfbefa2b0bdc3bc58cca8a7cd546803083d3a
      = 8548
      = 10659ce25588db4e4be6f7c791a79eb21a7f24aaaca76a6ca3b83b
u[1] = 80aaf95aa328fe7d569a1ac99f9cd216edf3915d72632f1a8b990e
      = 250c
      = 9243e5b6c480683fd533e81f4a778349a309ce00bd163a29eb9fa8
Q0.x = dbc8f549242bef33e030db21cffacd408d2c4264b93e476c6a8590
      = e7aa
      = 1457b60c12e00e47ceb3ce64b57e7c3c61636475443d704a8e2b2a
Q0.y = b0a5ac7e4b3909435416784e16e19929c653b1bdc9478a8e5331c
      = a9ae
      = 935d9f75f7a0babbc39c0a1c3b412518ed8a24bc2c4886722fb4b7
Q1.x = d4a747af98e4e2528c75221e2dff3424abb436e10539a74caafa
      = 3ea3
      = b44d9e34211b4028f24117e856585ed81448f3c8b934987a1c5939
Q1.y = c86048737a08d85934fec6b3c2ef9f09cbd365cf22744f2e4ce697
      = 62a4
      = dc996c1736f4319868f897d9a27c45b02dd3bc6b7ca356a039606e
msg   = 5406e131a0bbe8238208b327b00853e8af84b58b13443e70542556
      = 3323

msg   = q128_qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
      = qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
P.x   = a1861a9464ae31249a0e60bf38791f3663049a3f5378998499a832
      = 92e159a2fecff838eb9bc6939e5c6ae76eb074ad4aae39b55b72ca
P.y   = 0b9a
      = 580a2798c5b904f8adfec5bd29fb49b4633cd9f8c2935eb4a0f12e
u[0] = 5dfa0285680880296bb729c6405337525fb5ed3dff930c137314f6
      = 0401
      = c80390020e578f009ead417029eff6cd0926110922db63ab98395e
u[1] = 3bdfdd5d8a65b1a2b8d495dc8c5e59b7f3518731f7dfc0f93ace5d
      = ee4b
      = 1c4dc6653a445bbef2add81d8e90a6c8591a788deb91d0d3f1519a
Q0.x = 2e4a460313041b77c1b0817f2e80b388e5c3e49f37d787dc1f85e4
      = 324a
      = 9d355251e245e4b13ed4ea3e5a3c55bf9b7211f1704771f2e1d8f1
Q0.y = a65610c468b1cf70c6c2ce30dcaad54ad9e5439471ec554b862ec8
      = 875a
      = 6689ba36a242af69ac2aad955d15e982d9b04f5d77f7609ebf742
Q1.x = 9587feb7e5ce27490b9c72114509f89565122074e46a614d7fd7c8
      = 00bd
      = c4b3d3ad4d2d62739a62989532992c1081e9474a201085b4616da5
Q1.y = 706cab824693b9fb428a201bcd1639a4588cc43b9eb841dbca7421
      = 9b1f
      = 265286f5dee8f3d894b5649da8565b58e96b4cfd44b462a2883ea6
msg   = 4dbcda21a00706ea3fea53fc2d769084b0b74589e91d0384d71189
      = 09fb

msg   = a512_aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      = aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      = aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      = aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      = aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      = aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      = aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      = aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
```

```
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x = 987c5ac19dd4b47835466a50b2d9feba7c8491b8885a04edf577e1
5a9f2c98b203ec2cd3e5390b3d20bba0fa6fc3eecefb5029a31723
4401
P.y = 5e273fcfff6b007bb6771e90509275a71ff1480c459ded26fc7b10
664db0a68aaa98bc7ecb07e49cf05b80ae5ac653fbdd14276bbd35
ccbc
u[0] = 163c79ab0210a4b5e4f44fb19437ea965bf5431ab233ef16606f0b
03c5f16a3feb7d46a5a675ce8f606e9c2bf74ee5336c54a1e54919
f13f
u[1] = f99666bde4995c4088333d6c2734687e815f80a99c6da02c47df4b
51f6c9d9ed466b4fecf7d9884990a8e0d0be6907fa437e0b1a27f4
9265
Q0.x = d1a5eba4a332514b69760948af09ceaeddbbb9fd4cb1f19b78349c
2ee4cf9ee86dbc9064659a4a0566fe9c34d90aec86f0801edc131
ad9b
Q0.y = 5d0a75a3014c3269c33b1b5da80706a4f097893461df286353484d
8031cd607c98edc2a846c77a841f057c7251eb45077853c7b20595
7e52
Q1.x = 69583b00dc6b2aced6ffa44630cc8c8cd0dd0649f57588dd0fb1da
ad2ce132e281d01e3f25ccd3f405be759975c6484268bfe8f5e5f2
3c30
Q1.y = 8418484035f60bdccf48cb488634c2dfb40272123435f7e654fb6f
254c6c42e7e38f1fa79a637a168a28de6c275232b704f9ded0ff76
dd94
```

J.7.2. edwards448_XOF:SHAKE256_ELL2_NU_

J.8. **secp256k1**

J.8.1. secp256k1_XMD:SHA-256_SSWU_RO_

```
suite = secp256k1_XMD:SHA-256_SSWU_RO_
dst   = QUUX-V01-CS02-with-secp256k1_XMD:SHA-256_SSWU_RO_

msg   =
P.x   = c1cae290e291aee617ebaf1be6d73861479c48b841eaba9b7b585
      2ddfcb1346
P.y   = 64fa678e07ae116126f08b022a94af6de15985c996c3a91b64c406
      a960e51067
u[0]  = 6b0f9910dd2ba71c78f2ee9f04d73b5f4c5f7fc773a701abea1e57
      3cab002fb3
u[1]  = 1ae6c212e08fe1a5937f6202f929a2cc8ef4ee5b9782db68b0d579
      9fd8f09e16
Q0.x  = 74519ef88b32b425a095e4ebcc84d81b64e9e2c2675340a720bb1a
      1857b99f1e
Q0.y  = c174fa322ab7c192e11748beed45b508e9fdb1ce046dee9c2cd3a2
      a86b410936
Q1.x  = 44548adb1b399263ded3510554d28b4bead34b8cf9a37b4bd0bd2b
      a4db87ae63
Q1.y  = 96eb8e2faf05e368efe5957c6167001760233e6dd2487516b46ae7
      25c4cce0c6

msg   = abc
P.x   = 3377e01eab42db296b512293120c6cee72b6ecf9f9205760bd9ff1
      1fb3cb2c4b
P.y   = 7f95890f33efebd1044d382a01b1bee0900fb6116f94688d487c6c
      7b9c8371f6
u[0]  = 128aab5d3679a1f7601e3bdf94ced1f43e491f544767e18a4873f3
      97b08a2b61
u[1]  = 5897b65da3b595a813d0fdcc75c895dc531be76a03518b044daaa0
      f2e4689e00
Q0.x  = 07dd9432d426845fb19857d1b3a91722436604ccbbbadaad8523b8f
      c38a5322d7
Q0.y  = 604588ef5138cffe3277bbd590b8550bcbe0e523bbaf1bed4014a4
      67122eb33f
Q1.x  = e9ef9794d15d4e77dde751e06c182782046b8dac05f8491eb88764
      fc65321f78
Q1.y  = cb07ce53670d5314bf236ee2c871455c562dd76314aa41f012919f
      e8e7f717b3

msg   = abcdef0123456789
P.x   = bac54083f293f1fe08e4a70137260aa90783a5cb84d3f35848b324
      d0674b0e3a
P.y   = 4436476085d4c3c4508b60fcf4389c40176adce756b398bdee27bc
      a19758d828
u[0]  = ea67a7c02f2cd5d8b87715c169d055a22520f74daeb080e6180958
      380e2f98b9
u[1]  = 7434d0d1a500d38380d1f9615c021857ac8d546925f5f2355319d8
      23a478da18
Q0.x  = 576d43ab0260275adf11af990d130a5752704f7947862876172080
      8862544b5d
Q0.y  = 643c4a7fb68ae6cff55edd66b809087434bbaff0c07f3f9ec4d49b
      b3c16623c3
Q1.x  = f89d6d261a5e00fe5cf45e827b507643e67c2a947a20fd9ad71039
      f8b0e29ff8
Q1.y  = b33855e0cc34a9176ead91c6c3acb1aacb1ce936d563bc1cee1dcf
      fc806caf57
```


J.8.2. secp256k1_XMD:SHA-256_SSWU_NU_


```
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x = 17d22b867658977b5002dbe8d0ee70a8cfddec3eec50fb93f36136
      070fd9fa6c
P.y = e9178ff02f4dab73480f8dd590328aea99856a7b6cc8e5a6cdf289
      ecc2a51718
u[0] = a9ffbbee1d6e41ac33c248fb3364612ff591b502386c1bf6ac4aaf
      1ea51f8c3b
Q.x = 17d22b867658977b5002dbe8d0ee70a8cfddec3eec50fb93f36136
      070fd9fa6c
Q.y = e9178ff02f4dab73480f8dd590328aea99856a7b6cc8e5a6cdf289
      ecc2a51718
```


J.9. BLS12-381 G1

J.9.1. BLS12381G1_XMD:SHA-256_SSWU_RO_

```

suite = BLS12381G1_XMD:SHA-256_SSWU_RO_
dst   = QUUX-V01-CS02-with-BLS12381G1_XMD:SHA-256_SSWU_RO_

msg   =
P.x   = 052926add2207b76ca4fa57a8734416c8dc95e24501772c8142787
      00eed6d1e4e8cf62d9c09db0fac349612b759e79a1
P.y   = 08ba738453bfed09cb546dbb0783dbb3a5f1f566ed67bb6be0e8c6
      7e2e81a4cc68ee29813bb7994998f3eae0c9c6a265
u[0]  = 0ba14bd907ad64a016293ee7c2d276b8eae71f25a4b941eece7b0d
      89f17f75cb3ae5438a614fb61d6835ad59f29c564f
u[1]  = 019b9bd7979f12657976de2884c7cce192b82c177c80e0ec604436
      a7f538d231552f0d96d9f7babe5fa3b19b3ff25ac9
Q0.x  = 11a3cce7e1d90975990066b2f2643b9540fa40d6137780df4e753a
      8054d07580db3b7f1f03396333d4a359d1fe3766fe
Q0.y  = 0eeaf6d794e479e270da10fdaf768db4c96b650a74518fc67b04b0
      3927754bac66f3ac720404f339ecdcc028afa091b7
Q1.x  = 160003aaf1632b13396dbad518effa00fff532f604de1a7fc2082f
      f4cb0afa2d63b2c32da1bef2bf6c5ca62dc6b72f9c
Q1.y  = 0d8bb2d14e20cf9f6036152ed386d79189415b6d015a20133acb4e
      019139b94e9c146aaad5817f866c95d609a361735e

msg   = abc
P.x   = 03567bc5ef9c690c2ab2ecdf6a96ef1c139cc0b2f284dca0a9a794
      3388a49a3aee664ba5379a7655d3c68900be2f6903
P.y   = 0b9c15f3fe6e5cf4211f346271d7b01c8f3b28be689c8429c85b67
      af215533311f0b8dfaaa154fa6b88176c229f2885d
u[0]  = 0d921c33f2bad966478a03ca35d05719bdf92d347557ea166e5bba
      579eea9b83e9afa5c088573c2281410369fbd32951
u[1]  = 003574a00b109ada2f26a37a91f9d1e740dff8d69ec0c35e1e9f4
      652c7dba61123e9dd2e76c655d956e2b3462611139
Q0.x  = 125435adce8e1cbd1c803e7123f45392dc6e326d292499c2c45c58
      65985fd74fe8f042ecdeec5ecac80680d04317d80
Q0.y  = 0e8828948c989126595ee30e4f7c931cbd6f4570735624fd25aef2
      fa41d3f79cfb4b4ee7b7e55a8ce013af2a5ba20bf2
Q1.x  = 11def93719829ecda3b46aa8c31fc3ac9c34b428982b898369608e
      4f042babee6c77ab9218aad5c87ba785481eff8ae4
Q1.y  = 0007c9cef122ccf2efd233d6eb9bfc680aa276652b0661f4f820a6
      53cec1db7ff69899f8e52b8e92b025a12c822a6ce6

msg   = abcdef0123456789
P.x   = 11e0b079dea29a68f0383ee94fed1b940995272407e3bb916bbf26
      8c263ddd57a6a27200a784cbc248e84f357ce82d98
P.y   = 03a87ae2caf14e8ee52e51fa2ed8ee fe80f02457004ba4d486d6aa
      1f517c0889501dc7413753f9599b099ebcbbd2d709
u[0]  = 062d1865eb80ebfa73dcfc45db1ad4266b9f3a93219976a3790ab8
      d52d3e5f1e62f3b01795e36834b17b70e7b76246d4
u[1]  = 0cdc3e2f271f29c4ff75020857ce6c5d36008c9b48385ea2f2bf6f
      96f428a3deb798aa033cd482d1cdc8b30178b08e3a
Q0.x  = 08834484878c217682f6d09a4b51444802fdb3d7f2df9903a0dda
      db92130ebbf807fffa0eabf257d7b48272410afff
Q0.y  = 0b318f7ecf77f45a0f038e62d7098221d2dbbca2a394164e2e3fe9
      53dc714ac2cde412d8f2d7f0c03b259e6795a2508e
Q1.x  = 158418ed6b27e2549f05531a8281b5822b31c3bf3144277fbb977f
      8d6e2694fedceb7011b3c2b192f23e2a44b2bd106e
Q1.y  = 1879074f344471fac5f839e2b4920789643c075792bec5af4282c7
      3f7941cda5aa77b00085eb10e206171b9787c4169f

```


J.9.2. BLS12381G1_XMD:SHA-256_SSWU_NU_


```
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x = 0e7a16a975904f131682edbb03d9560d3e48214c9986bd50417a77
108d13dc957500edf96462a3d01e62dc6cd468ef11
P.y = 0ae89e677711d05c30a48d6d75e76ca9fb70fe06c6dd6ff988683d
89ccde29ac7d46c53bb97a59b1901abf1db66052db
u[0] = 0dd824886d2123a96447f6c56e3a3fa992fbfefdba17b6673f9f63
0ff19e4d326529db37e1c1be43f905bf9202e0278d
Q.x = 1775d400a1bacc1c39c355da7e96d2d1c97baa9430c4a3476881f8
521c09a01f921f592607961efc99c4cd46bd78ca19
Q.y = 1109b5d59f65964315de65a7a143e86eabc053104ed289cf480949
317a5685fad7254ff8e7fe6d24d3104e5d55ad6370
```

J.10. BLS12-381 G2

J.10.1. BLS12381G2_XMD:SHA-256_SSWU_RO_

```

suite = BLS12381G2_XMD:SHA-256_SSWU_RO_
dst   = QUUX-V01-CS02-with-BLS12381G2_XMD:SHA-256_SSWU_RO_

msg   =
P.x   = 0141ebfbdc40eb85b87142e130ab689c673cf60f1a3e98d693352
      66f30d9b8d4ac44c1038e9dcdd5393faf5c41fb78a
+ I * 05cb8437535e20ecffae7752baddf98034139c38452458baefab
      379ba13dff5bf5dd71b72418717047f5b0f37da03d
P.y   = 0503921d7f6a12805e72940b963c0cf3471c7b2a524950ca195d11
      062ee75ec076daf2d4bc358c4b190c0c98064fdd92
+ I * 12424ac32561493f3fe3c260708a12b7c620e7be00099a974e259d
      dc7d1f6395c3c811cdd19f1e8dbf3e9ecfcdcbab8d6
u[0]  = 03dbc2cce174e91ba93cbb08f26b917f98194a2ea08d1cce75b2b9
      cc9f21689d80bd79b594a613d0a68eb807dfdc1cf8
+ I * 05a2acec64114845711a54199ea339abd125ba38253b70a92c876d
      f10598bd1986b739cad67961eb94f7076511b3b39a
u[1]  = 02f99798e8a5acdeed60d7e18e9120521ba1f47ec090984662846b
      c825de191b5b7641148c0dbc237726a334473eee94
+ I * 145a81e418d4010cc027a68f14391b30074e89e60ee7a22f87217b
      2f6eb0c4b94c9115b436e6fa4607e95a98de30a435
Q0.x  = 019ad3fc9c72425a998d7ab1ea0e646a1f6093444fc6965f1cad5a
      3195a7b1e099c050d57f45e3fa191cc6d75ed7458c
+ I * 171c88b0b0efb5eb2b88913a9e74fe111a4f68867b59db252ce586
      8af4d1254bfab77ebde5d61cd1a86fb2fe4a5a1c1d
Q0.y  = 0ba10604e62bdd9eeeb4156652066167b72c8d743b050fb4c1016c
      31b505129374f76e03fa127d6a156213576910fef3
+ I * 0eb22c7a543d3d376e9716a49b72e79a89c9bfe9feee8533ed931c
      bb5373dde1fbcd7411d8052e02693654f71e15410a
Q1.x  = 113d2b9cd4bd98aee53470b27abc658d91b47a78a51584f3d4b950
      677cfb8a3e99c24222c406128c91296ef6b45608be
+ I * 13855912321c5cb793e9d1e88f6f8d342d49c0b0dbac613ee9e17e
      3c0b3c97dfbb5a49cc3fb45102fdbaf65e0efe2632
Q1.y  = 0fd3def0b7574a1d801be44fde617162aa2e89da47f464317d9bb5
      abc3a7071763ce74180883ad7ad9a723a9afafcdca
+ I * 056f617902b3c0d0f78a9a8cbda43a26b65f602f8786540b9469b0
      60db7b38417915b413ca65f875c130bebf5a59790c

msg   = abc
P.x   = 02c2d18e033b960562aae3cab37a27ce00d80ccd5ba4b7fe0e7a21
      0245129dbec7780ccc7954725f4168aff2787776e6
+ I * 139cddbccc5e91b9623efd38c49f81a6f83f175e80b06fc374de9
      eb4b41dfe4ca3a230ed250fbc3a2acf73a41177fd8
P.y   = 1787327b68159716a37440985269cf584bcb1e621d3a7202be6ea0
      5c4cfe244aeb197642555a0645fb87bf7466b2ba48
+ I * 00aa65dae3c8d732d10ecd2c50f8a1baf3001578f71c694e03866e
      9f3d49ac1e1ce70dd94a733534f106d4cec0eddd16
u[0]  = 15f7c0aa8f6b296ab5ff9c2c7581ade64f4ee6f1bf18f55179ff44
      a2cf355fa53dd2a2158c5ecb17d7c52f63e7195771
+ I * 01c8067bf4c0ba709aa8b9abc3d1cef589a4758e09ef53732d670f
      d8739a7274e111ba2fcaa71b3d33df2a3a0c8529dd
u[1]  = 187111d5e088b6b9acfdfad078c4dacf72dcd17ca17c82be35e79f
      8c372a693f60a033b461d81b025864a0ad051a06e4
+ I * 08b852331c96ed983e497ebc6dee9b75e373d923b729194af8e72a
      051ea586f3538a6ebb1e80881a082fa2b24df9f566
Q0.x  = 12b2e525281b5f4d2276954e84ac4f42cf4e13b6ac4228624e1776
      0faf94ce5706d53f0ca1952f1c5ef75239aeed55ad
+ I * 05d8a724db78e570e34100c0bc4a5fa84ad5839359b40398151f37

```



```
+ I * 12c3c913ba4ed03c24f0721a81a6be7430f2971ffca8fd1729aafe
496bb725807531b44b34b59b3ae5495e5a2dcdbd5c8
Q1.x   = 16ec57b7fe04c71dfe34fb5ad84dbce5a2dbbd6ee085f1d8cd17f4
5e8868976fc3c51ad9eeda682c7869024d24579bfd
+ I * 13103f7aace1ae1420d208a537f7d3a9679c287208026e4e3439ab
8cd534c12856284d95e27f5e1f33eec2ce656533b0
Q1.y   = 0958b2c4c2c10fcecf5a6c59b9e92c4a67b0fae3e2e0f1b6b5edad9
c940b8f3524ba9ebbc3f2ceb3cfe377655b3163bd7
+ I * 0ccb594ed8bd14ca64ed9cb4e0aba221be540f25dd0d6ba15a4a4b
e5d67bcf35df7853b2d8dad3ba245f1ea3697f66aa
```

J.10.2. BLS12381G2_XMD:SHA-256_SSWU_NU_

```

suite = BLS12381G2_XMD:SHA-256_SSWU_NU_
dst   = QUUX-V01-CS02-with-BLS12381G2_XMD:SHA-256_SSWU_NU_

msg   =
P.x   = 00e7f4568a82b4b7dc1f14c6aaa055edf51502319c723c4dc2688c
      7fe5944c213f510328082396515734b6612c4e7bb7
+ I * 126b855e9e69b1f691f816e48ac6977664d24d99f8724868a18418
      6469ddfd4617367e94527d4b74fc86413483afb35b
P.y   = 0caead0fd7b6176c01436833c79d305c78be307da5f6af6c133c47
      311def6ff1e0babf57a0fb5539fce7ee12407b0a42
+ I * 1498aadcf7ae2b345243e281ae076df6de84455d766ab6fcdad71
      fab60abb2e8b980a440043cd305db09d283c895e3d
u[0]  = 07355d25caf6e7f2f0cb2812ca0e513bd026ed09dda65b177500fa
      31714e09ea0ded3a078b526bed3307f804d4b93b04
+ I * 02829ce3c021339ccb5caf3e187f6370e1e2a311dec9b753631170
      63ab2015603fff52c3d3b98f19c2f65575e99e8b78c
Q.x   = 18ed3794ad43c781816c523776188deafba67ab773189b8f18c49b
      c7aa841cd81525171f7a5203b2a340579192403bef
+ I * 0727d90785d179e7b5732c8a34b660335fed03b913710b60903cf4
      954b651ed3466dc3728e21855ae822d4a0f1d06587
Q.y   = 00764a5cf6c5f61c52c838523460eb2168b5a5b43705e19cb612e0
      06f29b717897facfd15dd1c8874c915f6d53d0342d
+ I * 19290bb9797c12c1d275817aa2605ebe42275b66860f0e4d04487e
      bc2e47c50b36edd86c685a60c20a2bd584a82b011a

msg   = abc
P.x   = 108ed59fd9fae381abfd1d6bce2fd2fa220990f0f837fa30e0f279
      14ed6e1454db0d1ee957b219f61da6ff8be0d6441f
+ I * 0296238ea82c6d4adb3c838ee3cb2346049c90b96d602d7bb1b469
      b905c9228be25c627bffee872def773d5b2a2eb57d
P.y   = 033f90f6057aadacae7963b0a0b379dd46750c1c94a6357c99b65f
      63b79e321ff50fe3053330911c56b6ceea08fee656
+ I * 153606c417e59fb331b7ae6bce4fbf7c5190c33ce9402b5ebe2b70
      e44fca614f3f1382a3625ed5493843d0b0a652fc3f
u[0]  = 138879a9559e24cecee8697b8b4ad32cced053138ab913b9987277
      2dc753a2967ed50aabc907937aefb2439ba06cc50c
+ I * 0a1ae7999ea9bab1dcc9ef8887a6cb6e8f1e22566015428d220b7e
      ec90ffa70ad1f624018a9ad11e78d588bd3617f9f2
Q.x   = 0f40e1d5025ecf0d850aa0bb7bbeceab21a3d4e85e6bee857805b
      09693051f5b25428c6be343edba5f14317fcc30143
+ I * 02e0d261f2b9fee88b82804ec83db330caa75fbb12719cfa71ccce
      1c532dc4e1e79b0a6a281ed8d3817524286c8bc04c
Q.y   = 0cf4a4adc5c66da0bca4caddc6a57ecd97c8252d7526a8ff478e0d
      fed816c4d321b5c3039c6683ae9b1e6a3a38c9c0ae
+ I * 11cad1646bb3768c04be2ab2bbe1f80263b7ff6f8f9488f5bc3b68
      50e5a3e97e20acc583613c69cf3d2bfe8489744ebb

msg   = abcdef0123456789
P.x   = 038af300ef34c7759a6caaa4e69363cafeed218a1f207e93b2c70d
      91a1263d375d6730bd6b6509dcac3ba5b567e85bf3
+ I * 0da75be60fb6aa0e9e3143e40c42796edf15685cafe0279afd2a67
      c3dff1c82341f17effd402e4f1af240ea90f4b659b
P.y   = 19b148cbdf163cf0894f29660d2e7bfb2b68e37d54cc83fd4e6e62
      c020eaa48709302ef8e746736c0e19342cc1ce3df4
+ I * 0492f4fed741b073e5a82580f7c663f9b79e036b70ab3e51162359
      cec4e77c78086fe879b65ca7a47d34374c8315ac5e
u[0]  = 18c16fe362b7dbdfa102e42bdfd3e2f4e6191d479437a59db4eb71

```



```
Q.x      = 19592c812d5a50c5601062faba14c7d670711745311c879de1235a
          0a11c75aab61327bf2d1725db07ec4d6996a682886
+ I *    0eef4fa41ddc17ed47baf447a2c498548f3c72a02381313d13bef9
          16e240b61ce125539090d62d9fbb14a900bf1b8e90
Q.y      = 1260d6e0987eae96af9ebe551e08de22b37791d53f4db9e0d59da7
          36e66699735793e853e26362531fe4adf99c1883e3
+ I *    0dbace5df0a4ac4ac2f45d8fdf8aee45484576fdd6efc4f98ab9b9
          f4112309e628255e183022d98ea5ed6e47ca00306c
```

Appendix K. Expand Test Vectors

This section gives test vectors for `expand_message` variants specified in [Section 5.3](#). The test vectors in this section were generated using code that is available from [\[hash2curve-repo\]](#).

Each test vector in this section lists the `expand_message` name, hash function, and DST, along with a series of tuples of the function inputs (`msg` and `len_in_bytes`), output (`uniform_bytes`), and intermediate values (`dst_prime` and `msg_prime`). DST and `msg` are represented as ASCII strings. Intermediate and output values are represented as byte strings in hexadecimal.

K.1. `expand_message_xmd(SHA-256)`

K.2. `expand_message_xmd(SHA-256)` (Long DST)

K.3. expand_message_xmd(SHA-512)

K.4. `expand_message_xof(SHAKE128)`

K.5. `expand_message_xof`(SHAKE128) (Long DST)

K.6. `expand_message_xof(SHAKE256)`

Sam Scott

Cornell Tech
2 West Loop Rd
New York, New York 10044
United States of America
Email: sam.scott@cornell.edu

Nick Sullivan

Cloudflare, Inc.
101 Townsend St
San Francisco,
United States of America
Email: nick@cloudflare.com

Riad S. Wahby

Stanford University
Email: rsw@cs.stanford.edu

Christopher A. Wood

Cloudflare, Inc.
101 Townsend St
San Francisco, CA 94107
United States of America
Email: caw@heapingbits.net